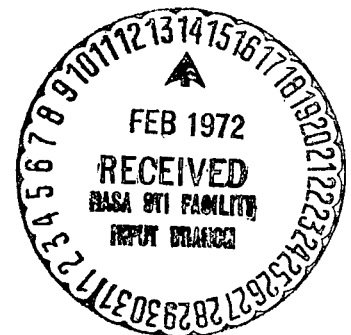
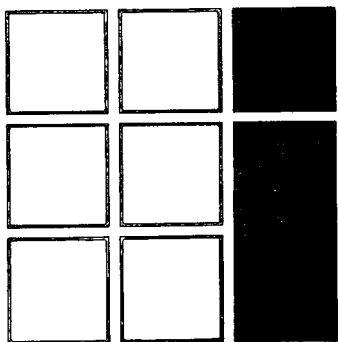


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Final Report

CONTINUED ADVANCEMENT OF THE  
PROGRAMMING LANGUAGE HAL TO AN  
OPERATIONAL STATUS

NAS 9-11944

December 30, 1971

N O T I C E

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## FOREWORD

This document represents the final report of a contract for the continued advancement of the programming language HAL to an operational status. The effort was sponsored by the National Aeronautics and Space Administration's Manned Spacecraft Center in Houston, Texas under Contract NAS-9-11944 . It was performed by Intermetrics, Inc., Cambridge, Mass. under the technical direction of Mr. Daniel J. Lickly. The Technical Monitor for NASA/MSC was Mr. John Garman, FS/6.

The publication of this report does not constitute approval by the National Aeronautics and Space Administration of the findings or the conclusions contained therein. It is published only for the exchange and stimulation of ideas.

## 1.0 TASK SUMMARY

The objectives and time duration of this contract were limited. Essentially, over the summer months of 1971 HAL was to be installed on the 360/75 at MSC, on-site support provided, users trained by class sessions, the compiler updated as necessary and experience gained with the language. The work was divided into two broad areas: maintenance and training, and advanced HAL development.

### 1.1 Task I: Maintenance and Training

Under this task, Intermetrics established a support systems programmer (Mr. Ronald Kole) at MSC within the Flight Software Section. Mr. Kole succeeded in solving some of the formidable problems associated with running at the RTCC under RTOS and transferred HAL from Intermetrics' 360/65 to MSC's 360/75. In addition, from time to time, necessary modifications originating in Cambridge and in Houston were incorporated at both sites and specifically, a general compiler-version update procedure was developed and implemented at MSC. The most significant modification was a redesign of the storage allocation algorithms, described in detail in Section 4.2 of this report.

Also as part of this task activity, 36 hours of training classes were conducted at MSC and the MIT Draper Laboratory. A total of approximately 50 people attended, including both government and industrial personnel.

Although no modifications were made which affected the HAL Specification or Guide documents, a new complete description of HALMAT, the intermediate code, was issued.

### 1.2 Task II: Advanced HAL Development

The objective of this development task was to increase the transferability of HAL to another host computer. The approach taken was to demonstrate that the compiler itself could be written in HAL. If this were accomplished then the entire compiler could be compiled on the 360 into FORTRAN and then the FORTRAN moved to almost any other large computer with only minor modifications. Toward this end a portion of the HAL compiler was coded in HAL and demonstrated to work on the 360/75 at the MIT Draper Laboratory. The portion selected exercised the bit and character handling features of HAL and indicates the feasibility of the approach.

## 2.0 HAL INSTALLATION AND ON-SITE SUPPORT

### 2.1 Installation

The HAL compiler and associated software were developed at the Harvard Computing Center in Cambridge, Massachusetts using a 360/65 running under OS/360 MVT Release 18. Transferability problems were encountered in establishing a usable HAL system at MSC. These problems may be broken down into two categories: (1) logistics and (2) internal software compatibility. Before the second set could even be recognized, the first had to be solved.

#### 2.1.1 Logistics Problems

The logistics problems are the ones associated with the differences in operating procedures and at the systems' two installations. The preparation and submission of jobs at Harvard were done almost entirely through the Conversational Remote Batch Entry System (CRBE). The Harvard system was equipped with four IBM 1403 printers, each having a full PL/1 character set. The entire HAL system was maintained on a disk pack that was mounted by the operator when needed.

In contrast, the RTCC at MSC uses IBM 360/75 computers running a modified OS MVT which they call RTOS (Real Time Operating System). The RTOS version available during the June-September 1971 installation period corresponded roughly to OS/360 Release 18. The RTCC does not support the CRBE system, so all input submission was done via punched cards. This in itself was something of a problem since some pieces of the compiler and even some updates to the compiler were quite long. Also, the handling of cards and the chance of error in mixing up cards was considered less desirable than an on-line editing/submitting system like CRBE.

2.1.1.1 RTCC Limited Disk Space. The RTCC also has very limited available disk space. In fact, it is impossible to have a permanently saved disk on the system. Therefore, a reasonable way had to be found to give users access to the HAL compiler. The method settled on was to use a dump/restore tape. When a user wants to run something that is located on the HAL disk (known as HAL001 at the RTCC) he tells the operator of the requirement for this special disk. The operator must "restore" the disk from a tape before the user's actual job can run. This means that he must run a background utility to transfer the contents of a special tape to some existing disk on the

system so that any user programs requesting HAL001 will find the disk mounted. After the user's run finishes, the HAL001 disk is scratched to make room for perhaps some other user's restored disk.

This system works, but it puts all jobs that require HAL001 in a special class that is usually run only during the early morning hours. This obviously precludes any same-day turnaround as can sometimes be obtained with "non-special" device requirements. There is, however, an advantage that is gained indirectly through use of the dump/restore tape. Since it is not any one physical disk that is mounted to satisfy requests for HAL001, it is possible to have serial versions of the disk available as separate dump/restore tapes. The user need only specify which dump/restore tape he wants used to create HAL001. Since one HAL001 is the same to the system as any other, all of his Job Control Language will work no matter which dump/restore tape is used.

The more troublesome aspect of the dump/restore system is the creation of an updated HAL001 and getting that new version put onto a D/R tape. The method of updating is to submit the job that makes the desired changes, requesting that HAL001 be restored in the normal manner, but also requesting that the disk to which the restore will take place be completely erased first. The job is then run as a regular batch job. After the run is completed the operator is requested to "dump" HAL001 to some specified tape. This tape then becomes the updated restore tape.

The problem in this system can occur in many ways. The disk to be restored must be erased first. This is to assure that after the update run, HAL001 contains only the HAL001 files. The dump program that does the disk to tape transfer will copy anything that it finds on the disk, no matter where it originated. So failure to erase the disk, while not really causing any errors in the resulting D/R tape, makes the tape file very long. This also means that subsequent restores from the tape will take a much longer time. All of this degrades performance of the overall system. It is also possible, if an unerased disk is cluttered enough, to overflow one D/R tape which causes even more complications.

Another source of error is the possibility that the operator will forget to dump the disk at all. Since this is an operator controlled utility, the user gets no indication on his output if the disk was actually dumped. The only real way to tell is to submit another job the next day to see if the updates were indeed saved. There is no way of telling if the restore tape file is too big because of a non-erased disk, other than watching the restore take place and guessing whether the tape moves a reasonable distance. Once an oversized tape file has been created it is very difficult to get rid of the unwanted "garbage". Thus, the



dump/restore tape method of supplying users with a HAL system works fairly well from the user's viewpoint (although turnaround is probably adversely affected) but is error prone from the system maintenance point of view.

The real solution to the D/R problem is to eventually catalogue all of the HAL system components on the system library as is done with the other language translators. This could be done now, but such a method is even more difficult to update and maintain and should not be used for anything short of a non-changing, well established version of the system.

2.1.1.2 RTCC Limited Print Facility. One other source of frustration at the RTCC can be identified. The RTCC is basically designed to run space missions and not necessarily to support batch users. The main outputs of mission programs are real-time displays and telemetry. These require little hard copy I/O. The main output of the batch user is printout. Yet the RTCC has only one printer per machine with two others shared between all machines. The 360/65 at Harvard manages to keep four printers busy by itself. So it would seem that the faster model 75 would generate a print backlog with even the maximum of three printers attached. This did indeed seem to be the case at MSC. This print backlog, of course, results in reduced turnaround. The HAL compiler proper output requires a full PL/1 character set. At the RTCC there is only one printer that has the UCS (Universal Character Set) feature that can support a print set other than the standard limited FORTRAN 48 character one. All of the output that requires this special character set is spooled on tapes until such time as there is enough to make the mounting of the special print train worthwhile. That point is reached two or three times a day. So HAL print delays are generally even longer than those caused by the regular print backlog. A small point further slowing down the printing is that the print train that was used for HAL was a special one, known as the Philco train. Some characters used by HAL only appear once on this train. This delay in waiting for the single character to move to the correct print position makes a large reduction in the speed of the printers (a visual estimate would be 20 to 40% slower).

The necessity of mounting the special character set created another potential error situation. If for some reason the class of output was mis-written on the spool tape, or misread, the output might (and has on occasion) be printed on the Fortran print train. It is, of course, then unreadable. The net result of these operation dependent procedures was a 1 run per day situation at best. Some preliminary users experienced even longer delays.

## 2.1.2 Software Incompatibility

The second class of problems relating to implementing HAL at MSC concern the operating system peculiarities that are present in any system. The most significant difference between RTOS at MSC and OS/360 was the handling of core allocation. One part of the HAL compiler called the Submonitor was significantly affected by the RTOS scheme.

The Submonitor is an assembler language program whose purpose essentially is to provide an I/O interface with the operating system. It also has the task of obtaining a block of core storage into which the actual compiler machine code can be placed for execution. In the original version of the Submonitor at Harvard, the necessary core was requested by means of a GETMAIN macro instruction. The operating system then gave the Submonitor the requested core out of the remaining part of the task's region as specified in the REGION parameter. Thus, it was possible to use the region parameter to determine how much core was made available to run the compiler.

2.1.2.1 OS/360 Core Allocation. The standard OS scheduling algorithms keep jobs in a waiting queue until such time as there is available in the machine, enough core in a contiguous block to satisfy the job's REGION request. Then the job is started and is free to do whatever it wants with its core. This block of contiguous core is reserved for the one task, even if the task uses only a small part of it. Thus, if an adequate REGION parameter is supplied, the Submonitor is guaranteed the availability of the core it needs to run the compiler. The Submonitor, in fact, makes maximum use of the entire available core by using a form of the GETMAIN Macro instruction that gives the operating system a minimum and maximum value of acceptable core regions. The maximum is set very high so that the operating system, in trying to come as close as possible to the maximum requested allocates all of the core remaining in the Submonitor's region.

2.1.2.2 RTOS Core Allocation. RTOS, however, does not make a job wait until the requested REGION is available in a contiguous block. It waits only until the sum of all free core in the system meets or exceeds the requested REGION. In fact, the decision of which jobs to run at any one time is based on a 125% allocation of core on the assumption that not all jobs will require their entire region at the same time. So in the case of the HAL compiler which requires about 4K bytes of core,

the operating system will load and execute the Submonitor as soon as there is 4K contiguous core available and enough other free core to satisfy the rest of the specified REGION parameter (this is usually 300K for HAL). The Submonitor then immediately issues a GETMAIN for the remaining core. This is where the operating system differences cause problems.

Under RTOS, a GETMAIN of the original form used at Harvard causes the operating system to put the Submonitor task into a wait state until the maximum of the minimum/maximum pair is available. For a large maximum this never happens and the job must eventually be cancelled by the operator.

2.1.2.3 Some Solutions. The first attempt to fix this was made by allowing the user to specify in his JCL the minimum and maximum values to be used in the GETMAIN request instead of the Submonitor's default values. This approach will apparently work for smaller programs, but for larger programs, there seems to be some system prescribed limit to the amount of core that can be obtained in this way. The Submonitor was never able to obtain more than 262, 144 (or  $2^{18}$ ) bytes of core. This is not enough to run the compiler and another method was needed.

The second method was to change the type of GETMAIN to one that requested a single specific size piece of core. The user was given the ability to specify this number through a keyword in his JCL. This system partially solved the problem. It was possible to obtain the correct amount of core this way, but another problem persisted. Even though RTOS could guarantee the existence of 300K bytes of free core, it could not guarantee how this core was divided up as could the regular OS. It was possible for the Submonitor to go into a wait state while the operating system tried to supply its contiguous core requirement. Under the right circumstances, this might take an hour or more. Whenever a task goes into the wait state, RTOS monitors how long it stays there. After some length of time RTOS begins sending the operator messages informing him of the lack of progress of the task. The operator makes the decision on the length of time the task is allowed to wait. After he gets tired of seeing the periodic messages, he usually cancels the job. The problem with this method is that the operator is never told why the job is waiting; it may be waiting because of some programmer error. He really has no chance to evaluate the situation.

The next step taken toward insuring a successful run was to take the wait for core out of the running Submonitor and put it in the pre-execution allocation. This simply meant doing away with the GETMAIN and giving the Submonitor a built-in

storage area big enough to run HAL. Under this system, the Submonitor became about 300K bytes long. This forced the operating system to find the 300K in a contiguous region before it could even load the Submonitor. As before, it was still possible for the available core to total more than 300K and yet not have 300K contiguous, but now the operator got a message saying "JOBXXX AWAITING A REGION". In this way the operator was informed of the real cause for the delay and was more willing to let the job wait for the core to become available.

Although the probability of getting a job to run has been increased, it is still possible for the job to be cancelled when it was the only user job in the system and was still unable to get its core. This happens when there are "background" utilities running. Although they are termed "background", they still compete for core like any other job on the system and it is possible for them to tie up core in such a way that a HAL job will not run even though it appears to be all alone on the system.

One small drawback to this final state of the Submonitor is that it is no longer possible to use the REGION parameter, or any user keywords to limit the size of the available core. This is not considered much of a handicap since the size of the HAL compiler is quite stable and is expected to remain so.

There is an alternate approach to solving the core lockout problem. This would involve a redesign of the HAL compiler structure to give it a scatter loadable attribute. If this were done, the required core would not need to be contiguous; several smaller contiguous areas would be requested. The probability of finding these smaller areas would be greater than the present system. There would, however, still be a finite chance that even these smaller regions would not all be available. The situation is such that the more the core requirement is split up, the more chance there is that the resulting smaller pieces will be found. More pieces of code, however, require much more overhead to maintain. Also, the redesign of the HAL code to allow such a split would be a difficult job that would not further the goal of producing a better overall compiler. The frequency of run failures under the present system is very low and sporadic. For a large, ground-based, batch-oriented system like the RTCC, additional time spent on refinement of HAL running procedures, would be of little value compared to the same time spent on refinement of the actual compiler code.

This potential lockout problem is not peculiar to HAL. It can happen to any job on the system whose core requirements are of the same size and nature as HAL's.

## 2.2 Program Changes and Maintenance Procedures

During the June-September 1971 time period many changes were made to both the HAL Submonitor and the compiler itself. Most of the changes made to the Submonitor are detailed in 2.1 above. These changes were basic in nature and took a large part of the summer to research and implement properly at MSC. In addition, new compiler versions were developed at Intermetrics and were sent to Houston on magnetic tape where copies of all of the files on the tape were put on MSC-owned tapes. The required files were then transferred from tape to the HAL001 disk and the disk then dumped to a dump/restore tape as described above.

### 2.2.1 Updates

A three tape dump/restore system was established to maintain the integrity of the system. One tape, called the system tape was the only one available to users. It always contained the most recent released version of the compiler. Thus, users only had to have this tape number to run HAL compilation. The two other tapes were development tapes. They were used in an alternating manner to build and checkout a new release. The alternation was necessary to provide a backup in case of some failure to make a good update. To make an update, the newest version development tape was used to restore HAL001. The update was made to the disk and then HAL001 was dumped to the alternate development tape. Even if the dump was not done, or if the update was unsuccessful, the original restore tape was still intact.

Once a version was considered ready for release, the development tape on which it resided was simply copied onto the system tape. Users specifying the system tape number as the HAL001 dump/restore automatically got the new release.

In addition to updates originating in Cambridge, some changes were made to the compiler at MSC. Small changes were communicated to Cambridge directly by long distance through the CRBE system at Harvard. In the case of larger updates, tapes were exchanged.

### 2.2.2 Summary of Changes Made At MSC

- a) Research and implement the changes to the Submonitor to allow a more reliable core allocation.

- b) Fix numerous small bugs found during checkout both in the HAL code itself and in the HAL run time library.
- c) Partially implement and lay the ground work for a more complete listing generator as detailed in the HAL Guide. This involved providing additional functions in the Submonitor to allow the HAL compiler to set a maximum number of lines per page of listing and to dynamically request the line number of the current line on the page. These new functions helped to lay the framework for the ability to control completely the layout of the HAL listing. The listing was changed to the extent that the statement and line numbers were made available and the format of the format of the printed source code changed to increase readability.

## 3.0 HAL COURSES

### 3.1 General Description

Intermetrics personnel prepared and conducted three HAL language courses during the contract period. The material was designed for two types of audiences: 1) those seeking a broad "brush" overview of HAL, 2) those intending an indepth exposure to HAL. Two 15-hour sessions (2 1/2 days each) were given at the Manned Spacecraft Center in Houston. The first, primarily for NASA personnel and the second, for industrial contractors and other government agencies with an interest in higher order languages. For each session, the first three hours were devoted to the HAL overview; however, the overview itself was considered an integral part of the longer course.

A special third session was also conducted at the MIT Draper Laboratory, for Laboratory personnel and local industrial contractors. Because of the familiarity of these personnel with MIT's MAC language and certain similarities between HAL and MAC, an effective 1-day, 6-hour course was held. The course consisted of the overview, with elaborations and discussions, followed by a rapid presentation of the salient features of HAL.

In general, the participating students at MSC and MIT were highly motivated to learn HAL and always attempted the place HAL in perspective with respect to Shuttle applications. As a result, many provocative questions were asked and in some circumstances material discussed in class was fed back into the HAL design.

### 3.2 Course Preparation

The HAL courses were prepared with two objectives in mind: an overview, and a detailed study. For the overview, a balanced presentation of most of the important features and rationale incorporated into the HAL Specification Document (MSC-#01846) was designed. The purpose here was to illustrate how HAL satisfied, for the most part, the requirements imposed on a programming language for the Shuttle. Toward this end readability, vector-matrix arithmetic, data management, systems programming, real-time control and software reliability were emphasized.

The material was presented in vu-graph form and included numerous "Shuttle-like" application examples and commentary which included Intermetrics' experience with Apollo software

development. Particular attention was paid during the overview to indicate which HAL features would not be included in the first implementation for the IBM 360/75 at MSC. This was especially true for the descriptions of real-time control, controlled data sharing, and the error recovery features.

The longer HAL course was designed as an in-depth study of HAL and the overview served as an excellent orientation. This part of the course was based closely on the HAL Guide (MSC #01848) and the material was a combination of vu-graphs, references, to the Guide text and blackboard work. Only those features actually intended for implementation on the first 360/75 version were covered. (This specifically excluded real-time control, etc.)

Levels of increasing detail were presented, first with a set of vu-graphs covering all of the language features of HAL; i.e., operations, declarations, indexing, control, etc. followed by a careful tour through selected portions of the Guide. The Guide work illustrated usage, described many examples and motivated class discussions (and, in fact, contributed to subsequent corrections to the Guide). The technique of repeating subject material in levels of increasing detail; i.e. from overview to construct description to Guide with examples, proved to be an effective method of rapid assimilation and study.

In addition to text material and lecture, each student was provided with a HAL problem set as a homework exercise. Unfortunately, few found the time to actually address these problems out of class. However, during the last class session, prepared problem answers were distributed and each problem was carefully "talked-through". Actual runs on the 360/75 by the students were contemplated during the course preparation, but 360 turn-around time within the RTCC facility was not consistent with the 2 1/2 day course duration.

An outline of the HAL course material is presented in the next section and the vu-graphs for both the overview and the longer course are collected in Appendix .

### 3.3 Course Outline

#### 3.3.1 Overview (vu-graph material)

1. Higher order language motivation and capabilities.



2. Salient features of HAL
3. Data types
4. Program organization and structure
5. HAL Statements
6. Specific Examples
7. Real-time control, including data sharing and error recovery
8. Summary

### 3.3.2 Longer Course (Vu-graph Material)

1. Data Operations
2. Data Declarations
3. Indexing: partitions and use of subscripts
4. Control and branching mechanisms
5. Name scope rules

### 3.3.3 Longer Course (Guide Material)

1. Two-dimensional input-output format
2. HAL<sub>M</sub> (HAL Mathematical Subset)
  - a. Data and declarations
  - b. Arithmetic expressions
  - c. Assignment statements
  - d. User-defined functions (SCALAR, VECTOR, MATRIX)
  - e. IF Statements
  - f. Illustrative problems - I
  - g. Subscripts
  - h. DO Statements
  - i. Illustrative problems - II
  - j. Subroutines; i.e., HAL PROCEDURES

- k. Illustrative problems - III
  - l. Name scope
  - m. I/O Facilities
  - n. Illustrative problems - IV
- 3. Integer and Bit String Data
- 4. Structures
- 5. Bit and Character String Manipulations
- 6. Subscript facilities: complete
- 7. Implicit conversion of mixed data types
- 8. User-defined functions: complete
- 9. Array processing
- 10. Shaping functions: complete
- 11. REPLACE and DEFAULT Statements
- 12. "Talk-through" of problem set

## 4.0 NECESSARY MODIFICATIONS AND ADDITIONS

Redesign work was undertaken to increase the scope and capabilities of the HAL compiler and to promote its transferability to other computers. The first step was to redesign the variable storage philosophy and mechanization. Extensive design sessions were conducted to develop a suitable memory storage allocation system that would support the most general future goals of HAL, especially transferability (see 4.1 for more detail). This included the techniques necessary to support the calling of separately compiled HAL programs and the sharing of their data through a COMPOOL. This capability, in some form, is vital to the production of a multipass compiler.

During this time items were also dealt with that were either incomplete or had been newly defined. Thus, certain "holes" in HAL's capabilities were filled in. In addition, a number of shortcomings which had been uncovered were remedied.

### 4.1 Storage Allocation Problem

Certain storage allocation problems encountered during the implementation of some of the more advanced features of HAL in Phase II of the compiler (Fortran code generation) had necessitated basic conceptual changes in the allocation algorithms in the compiler.

In the original version of the algorithms, temporary storage required for partial numerical results was allocated when needed during the code generation of a HAL statement, and freed-up again not later than at the end of the statement. This caused two major difficulties. Firstly, when temporary storage was required to hold the value of an argument in a procedure or user function invocation, special "unfreeable" temporary storage had to be used to prevent it from possibly being reallocated in the body of the procedure or user function. Secondly, in HAL statements containing user function invocations, (possibly nested), temporary storage allocated for partial results before the invocation code was generated had also to be masked "unfreeable" for the same reason. Other more subtle considerations finally made a complete restructuring of the algorithm essential.

The idea of providing completely dynamic storage allocation of execution time was rejected as requiring too many basic changes in the mode of operation of Phase II of the compiler. Instead a static scheme similar in some respects to the original

was adopted. In this scheme temporary storage for the program, and for each procedure or user function are allocated within mutually non-overlapping segments. A program or subprogram may have one segment or several non-contiguous segments of varying sizes dedicated for its use, depending on its requirements. The sum total of all segments constitutes a single continuous area of storage (except possibly for word boundary alignments). At the microscopic level within the bounds of a segment, storage is allocated and freed exactly as it was in the original allocation scheme.

Under this scheme, no temporary storage need be marked "unfreeable" no matter to what use it is put. Furthermore, the scheme has resulted in considerable simplification and unification of other storage allocation mechanisms in operations at code generation time.

## 4.2 Miscellaneous Improvements

### 1. Arraynesses:

- \* reorganization of the mechanism controlling the utilization of statement arraynesses, especially with regard to utilization by arrayed subscripts of arrayed variables, and by the arguments of user functions.
- \* implementation of the arrayed subscripted variable as an input or assign argument in a function or procedure call.

### 2. Cosmetics and Statistics:

- \* generation of Phase II timing information, improvement of error message format, generation of statistics on certain critical parameters of Phase II operations
- \* introduction of toggle directives to control Phase II and subsequent Fortran IV operation.

### 3. Shaping Functions:

- \* introduction of a limited range of shaping and conversion functions: INTEGER, SCALAR, MATRIX, and VECTOR (no arrayed arguments or results).

### 4. Program Calling:

- \* setting-up operating mechanisms for calling independent (i.e. separately compiled) HAL programs to "any" nest level, non-recursive

- \* creation of mechanisms for saving HAL programs in an object library.

5. I/O Routines

- \* first, implementation of full-scale HAL READ/WRITE statements fixed, uni-channel input and output, fixed record length (printer and punch only).

6. Bit Strings:

- \* fundamental bit string operators were implemented. Included were terminal and array subscripting and the AND, OR, and NOT operations. Bit strings are limited to not more than 32 bits; they have been implemented in full-word, half-word, and byte form.

## 5.0 HAL TRANSFERABILITY

### 5.1 Technical Approach

#### 5.1.1 Background

The quest for easy transfer of operational programs from one computer to another has occupied the minds of many men since the early days of computer technology. The importance of this capability has grown considerably as the computer explosion has populated our society with countless kinds and types of computers with ever decreasing and more attractive price tags, and yet soaring software costs through higher programmer salaries has made conversion more difficult due to the huge investment in operational software for existing computers. The solutions to the programming transferability problems can be categorized into one of the following types:

1. Hardware emulators - In order to maintain compatibility many modern computers have included hardware or micro-program features that permit them to simulate other (usually older) computers. Thus, existing programs can still be executed.
2. Software translators - A program is developed that will take programs that were written for machine X and translate them into equivalent programs for machine Y. This approach has been limited since the technique is seldom 100% successful, even when the two computers are almost identical.
3. Higher level languages - If programming is confined to high level languages, hopefully machine independent, and a translator or compiler is used to produce the actual machine code, then it should theoretically be possible to feed the same higher level source statements into a translator to another brand of computer and produce a program that performs functionally equivalent tasks. The difficulty here is whether the language and the interpretations given it by compiler writers are truly machine independent.

#### 5.1.2 Level of Transfer

In the design of the HAL compiler system for the 360 implementation, Fortran was adopted as the output language from the code generator. A principal reason expressed for the

somewhat unusual procedure was to promote machine transferability. Fortran IV is the most widely used programming language and ANSI Fortran IV purports to be defined in a machine independent way. Production of the HAL code generator was initiated with the avowed intent of producing ANSI standard Fortran IV. If this could have been rigidly adhered to, transferability would have been automatically produced at the lowest level. The output of Pass 2 would be suitable for submission to any Fortran compiler. As it is, there exists some 360 specific Fortran output and some assembly language subroutines, but the job required to take the Fortran output of the HAL compiler and move to another computer is a minor one. Figure 5.1 depicts the steps in the HAL compilation process. The Fortran output of Pass 2 may be physically moved (in card or tape form) to another computer facility.

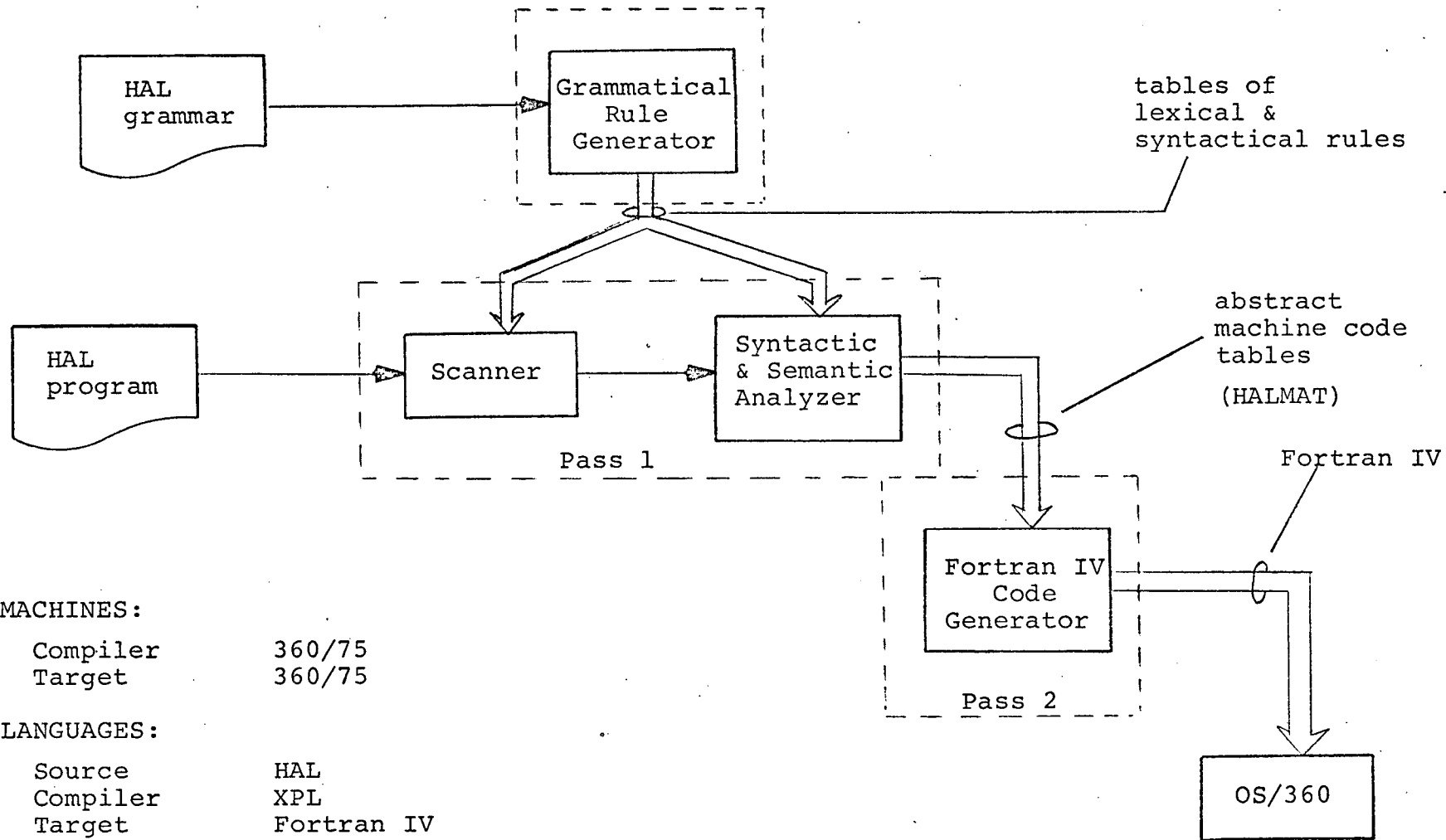
Contrast this transferability with the proposed system for construction of a HAL code generator for a flight computer. (See Figure 5.2) In this case, a new Pass 2 is required, the same output of Pass 1 (HALMAT) is used. This is the traditional approach. Every time that HAL is desired for a different target computer, another version of Pass 2 is required. This is a mid-level transfer.

However, neither of the above approaches will satisfy the needs of another general purpose computer facility. The reason is that they are only partial transfers. Although they produce code for another computer, the compiler itself still must run on the initial computer, the IBM 360 in this case. This is poor operationally. It means that a user must submit his HAL source program to the 360 for compilation and then take the object program to the other computer for execution. (This approach is perfectly adequate for a flight computer where the usual mode of operation is via simulation on his general purpose computer. Besides, the flight computer is usually of such limited size that compilation on it is not possible even if one were physically available.)

A total transfer is needed for implementation on another large commercial computer. It requires that the entire system be transferred, "lock, stock and barrel". Then the user can compile and execute on the new facility with no further need of the 360. This is a more demanding requirement since it necessitates moving the entire compiler to a different computer complex. The result is a high level transfer or complete conversion.

### 5.1.3 Method of Attack

There are three avenues of approach that might be followed to achieve a compiler transfer. They are:



**MACHINES:**

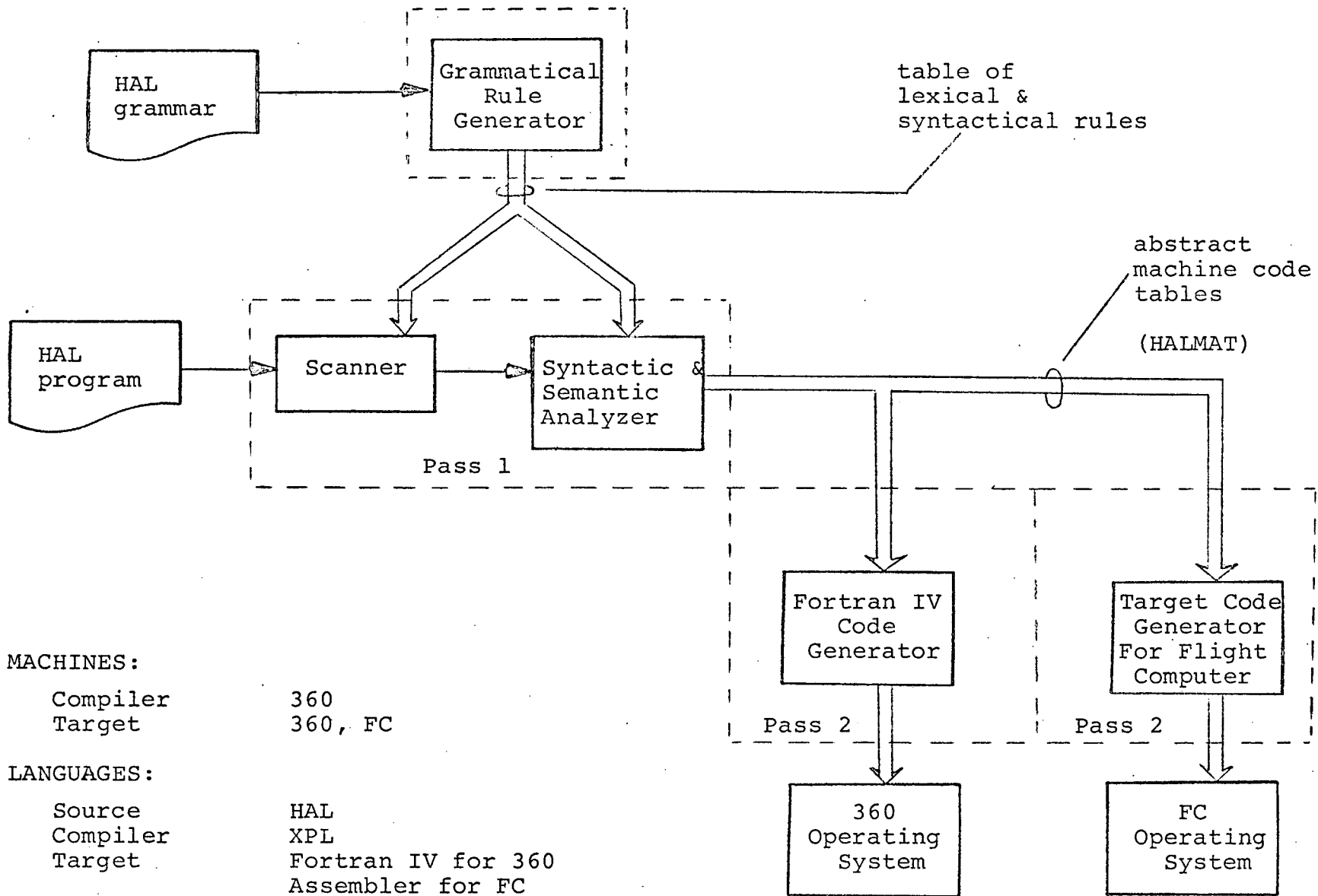
Compiler      360/75  
 Target        360/75

**LANGUAGES:**

Source        HAL  
 Compiler     XPL  
 Target        Fortran IV

Fig. 5.1 Construction of the HAL Compiler System





**MACHINES:**

Compiler	360
Target	360, FC

**LANGUAGES:**

Source	HAL
Compiler	XPL
Target	Fortran IV for 360 Assembler for FC

Fig. 5.2 Proposed Construction of HAL Compiler for Flight Computers

1. Reprogram the HAL compiler for Brand X: This technique looks at the process as a one-of-a-kind step and selects whatever seems most appropriate for machine X, be it assembly language or whatever. Then the job is done. This is the brute force approach and has no generality whatsoever.
2. Reprogram the XPL compiler for Brand X: HAL is written in XPL, a simplified subset of PL/1. Thus, it would be relatively easy to transfer to another computer that supports PL/1; however, there are few that do. But we could transfer XPL to another computer. Since XPL is itself written in XPL, the transfer could be accomplished by a mid-level transfer. (A new code generator on the 360 that produced code for Brand X would allow a version of XPL to be compiled that would execute on Brand X.) However, this approach also lacks generality; each new computer requires another code generator, itself not an easy task.
3. Reprogram HAL into a language more widely supported: If HAL could be rewritten in a language that was universally supported, than transfer problems would be minimized. The most widely used language is Fortran. And since Fortran is now produced by HAL, an interesting variation of this technique is immediately suggested. If HAL was rewritten in HAL and compiled on the current HAL compiler then the result would be Fortran source cards that would be suitable for compilation on any computer with a Fortran compiler. Thus, the transfer of HAL to almost any large scale computer could be achieved by minor changes to the Fortran output (chiefly in the area of data types and declarations) and the recoding of machine-dependent library routines. But the latter must be done anyway if HAL is to execute on Brand X; even the low-level transfer needed it. The extra task is the effort needed to rewrite HAL in HAL. But having done it once, it would not need to be done again to affect other transfers. The generality of this approach resolved the issue in its favor.

## 5.2 Translation of XPL Programs Into HAL

### 5.2.1 Introduction

This is a brief discussion of the methods used when translating a program from XPL into HAL. It is intended to provide a useful guide to a process which requires a considerable amount of analysis and judgement on the part of the individuals performing the work. This end is achieved by presenting the

essentials of language differences and by discussing examples of coding economies possible through the use of HAL. Explicit illustrations demonstrate the translation of several XPL constructs into HAL.

Translation of some form was necessary since it was intended that a copy of the HAL compiler be implemented on the Univac 1108. There were two general strategies available to Intermetrics as alternate means to affect this implementation. As one possibility, we could have rewritten or modified the XPL system to implement it on the 1108, then we would have been able to recompile the original XPL source code of HAL on the 1108. This approach lacks generality and involves the difficulties of emitting executable and efficient low level code for a machine with extant high level software. As a second alternative, we could rewrite the HAL compiler in a source language which maps via an existing processor onto a target language recognized by existing 1108 software.

This latter course was chosen, using HAL itself as the source language, and using the HAL/360 compiler as the mapping onto Fortran IV, a target language understood by the 1108 (as well as other large scale computers). This course provides a large amount of generality, and also proves to be easiest to carry out because of HAL's many high level features and the convenient degree of similarity between HAL and XPL.

The two-dimensional input scanner employed in Pass 1 of the HAL compiler was chosen as an initial goal. If a program as complex as this worked satisfactorily once debugged, we could be fairly certain that no part of the compiler would create a problem. The translation strategies and methods described in this document were devised in the process of successfully rewriting the input scanner. As an added bonus of this choice of translation strategy, the use of HAL as a source language proved to be exceptionally helpful in the process of debugging the current HAL/360 compiler. Quite a number of bugs which were invisible prior to this large scale application were exposed and repaired in the process.

## 5.2.2 Methodology

5.2.2.1 Variables. Variable declarations differ somewhat between XPL and HAL. Each individual DECLARE statement must be examined for possible changes.

BIT variables are declared identically in both XPL and HAL. The length specification is also identical in the two languages.

Ex. DECLARE X BIT(8);

FIXED variables in XPL are functionally identical to INTEGER variables in HAL. Therefore, mere substitution of the word INTEGER for FIXED is all that is necessary to make the language change. The REPLACE facility in HAL is the simplest method of substitution. Note here that any XPL variable names which correspond to HAL reserved words must be changed or augmented; i.e., the XPL identifier VECTOR could become VECTOR1 in HAL. Note also that the break characters @, #, and \$ are not legal identifier break characters in HAL and XPL identifiers using them must be replaced by legal HAL identifiers.

CHARACTER variables have somewhat different properties in XPL and HAL. In XPL, character variables are implicitly varying with a maximum length of 256 characters. However, VARYING character strings in HAL are currently limited to a length of 255. Thus, the general substitution rule for character declares is to change:

DECLARE ALPHA CHARACTER; (XPL)

to

DECLARE ALPHA CHARACTER (255) VARYING; (HAL)

In cases where a string is known to have a maximum length considerably less than 255 characters, it may be declared as such. Also, if a string is to be of fixed length (as with an initial unchanging value), the VARYING attribute should also be omitted.

Factored declarations in XPL and HAL are also implemented differently and involve a complete rewriting of the statements. For example:

DECLARE (I,J,K) FIXED, L BIT(8); (XPL)

becomes:

DECLARE INTEGER, I, J, K; (HAL)

DECLARE L BIT(8);

A word of caution is necessary at this point. XPL initializes all FIXED and BIT variables to "0" and all character strings to null strings unless otherwise specified by the INITIAL modifier. Any variable not explicitly initialized in HAL will have unpredictable contents. When in doubt as to whether the program itself initializes variables, include an INITIAL(0) specification

on the DECLARE statement (INITIAL(') for character strings).

The LITERALLY attribute in XPL is used to perform macro substitution for identifiers. The REPLACE statement in HAL performs the same function. The statement:

```
DECLARE FOREVER LITERALLY 'WHILE "1";           (XPL)
```

becomes

```
REPLACE FOREVER BY 'WHILE TRUE';             (HAL)
```

5.2.2.2 Arrays. When transferring array declarations and specifications from XPL to HAL there are a number of ground rules to follow. First, XPL subscripts start at 0 and the dimension specified is the highest allowable subscript. Therefore, an XPL array declared with an arrayness of 99 actually consists of 100 elements and must be declared as such in HAL, since all HAL subscripts start at 1 for an array. Thus,

```
DECLARE ABLE(99) FIXED;                       (XPL)
```

becomes

```
DECLARE ABLE ARRAY(100) INTEGER;             (HAL)
```

The word ARRAY must be supplied in HAL in array declarations.

Frequently in XPL the name of an array appears without a subscript. This means an implied reference to the 0th element of the array. However, in HAL, an array name without an explicit subscript implies reference to the entire array, not the first element. Therefore, for conversion, all such occurrences of non-subscripted array names must be translated with the explicit subscript of 1. Thus, the following XPL segment:

```
DECLARE ARR (9) FIXED, B CHARACTER;
```

```
B = ARR;
```

becomes in HAL:

```
DECLARE ARR ARRAY (9) INTEGER,
```

```
      B CHARACTER (255) VARYING;
```

```
B = ARR1;
```

In general, unless the 0th element of an XPL array is known not to be used, indexing expressions must be augmented by adding one to the original subscript expression, and not by changing the computation of indices in other statements. This is especially true when array references are made using Boolean values of 0 and 1 as switches for referring to one of two array elements. Thus, the XPL sequence:

```
IX = IY + IZ;  
VALUE = ABLE (IX);
```

should become in HAL:

```
IX = IY + IZ;  
VALUE = ABLEIX + 1;
```

as opposed to the HAL sequence:

```
IX = IY + IZ + 1;  
VALUE = ABLEIX;
```

as the latter form could possibly change the operational characteristics of the program.

Finally, XPL allows the specification of a subscript on a variable which is not declared as an array. This allows certain machine dependent coding "tricks" to be performed. Consider the following XPL sequence:

```
DECLARE INDEX FIXED, INDEXTAB (199) BIT (8);  
DO I = 1 TO 50;  
  INDEX (I) = 0;  
END;
```

This program in effect zero's out INDEXTAB with 50 references, rather than the 200 required to clear the individual INDEXTAB elements. This sequence is illegal in HAL and may be coded as follows in HAL:

```
INDEXTAB = 0;
```

where the non-subscripted version of the name implies setting the array to zero.

An extra step is involved when translating statements utilizing arrays of character strings or bit strings. A colon (:) must follow the array element subscript, to distinguish it from the individual character or bit subscript which is the default in HAL for these types of variables. It may be used following any array subscript, but is required in the above named instances to prevent ambiguity.

Ex:

```

DECLARE A CHARACTER(5),
        B ARRAY(10) CHARACTER(10);

B1 TO 5 = A;

B1 TO 5: = A;

```

In the first statement, characters 1 to 5 of all ten array elements of B are set to the value of the characters in A; in the second statement, the first five array elements of B are set to the value of A (padded with blanks to make the total length ten).

**5.2.2.3 Built-in Functions.** The XPL functions ADDR and INLINE are not available in HAL, and because of the machine independence of the language no corresponding functions exist. In the HAL compiler, fortunately, most such functions are used to manipulate data types not existing in XPL but which do exist in HAL.

The SUBSTR and BYTE functions in XPL are replaced with character string subscript notation in HAL. Examples of both forms of BYTE and SUBSTR substitution follow:

- a. BYTE (CHAR\_STRING) (XPL)  
     becomes  
     CHAR\_STRING<sub>1</sub> (HAL)
- b. BYTE (CHAR\_STRING, N) (XPL)  
     becomes  
     CHAR\_STRING<sub>N+1</sub> (HAL)
- c. SUBSTR (CHAR\_STRING, START) (XPL)  
     becomes  
     CHAR\_STRING<sub>START+1</sub> TO # (HAL)

d. SUBSTR (CHAR\_STRING, START, N) (XPL)

becomes

CHAR\_STRING<sub>N</sub> AT START+1 (HAL)

The functions SHL and SHR are used for doing word manipulation in XPL. For positive arguments, the SHL function may be replaced by multiplication by the appropriate power of two. The SHR function is more complex as integer division is not allowed in HAL. Since SHR is normally used to isolate a field of a packed word, the BIT shaping function can be used to achieve the same results.

For example, the XPL sequence:

```
DECLARE (ENTRY, PART) BIT (16),  
        WORD FIXED;  
PART = SHR (ENTRY, 4);  
PART = SHR (WORD, 16);
```

becomes in HAL:

```
DECLARE BIT (16), ENTRY, PART;  
DECLARE WORD INTEGER;  
PART = ENTRY1 TO 12;  
PART = BIT1 TO 16 (WORD);
```

In XPL, it is legal to assign the result of a relational expression to a BIT type variable. This is illegal in HAL. Thus, the statement

```
TEST = A > B; (XPL)
```

must become:

```
IF A > B THEN TEST = TRUE; ELSE TEST = FALSE; (HAL)
```

5.2.2.4 Constants. The following constant conversion rules apply:



"ABDF" or "(4)ABDF"	becomes	HEX'ABDF'
"(1)11010"	becomes	BIN'11010'
"(3)70346"	becomes	OCT'70346'

The quartal constant "(2)20312" must be converted to either BIN'1000110110' or HEX'236'.

The use of the BYTE function, notably BYTE('C'), to allow use of the internal representation of the character as a numeric quantity is accomplished in HAL by stating BIT('C'), or INTEGER (BIT('C')) where implicit Bit-to-Integer conversion may not take place.

5.2.2.5 Procedures. In XPL, all subroutines and functions are declared as PROCEDURE's. The RETURN statement may or may not pass back a value. If an XPL PROCEDURE which returns a value is called by the CALL statement, the returned value is ignored. In HAL, there are two classes of routines: PROCEDURE's and FUNCTION's. A PROCEDURE does not allow a value to be returned in the RETURN statement, whereas a FUNCTION demands that a value be returned. Thus, XPL PROCEDURE's that return values must be declared as FUNCTION's in HAL. Any such FUNCTIONS invoked by the CALL statement in XPL must be changed to the form:

```
DUMMY_VARIABLE = FUNCTION_NAME(X);
```

where the dummy variable is some unused name in the HAL program with the mode of the called function.

Also, in XPL, all formal parameters are call-by-value parameters. This presents a problem in HAL because, 1) PROCEDURE and FUNCTION parameters may not be assigned values within HAL programs, unlike XPL which freely allows such assignments, 2) the alternative in HAL, the ASSIGN list, is treated as a list of call-by-reference parameters, where assignments to such parameters are passed back to the calling program, whereas in XPL, parameter assignments do not reflect back to the calling program. Therefore, in all FUNCTIONS and PROCEDURES where assignments to formal parameters are made, a procedure prologue must be coded to assign the formal parameter (with an augmented name) to a local variable with the same declared properties with the original parameter name. Thus, the following XPL program segment:

```

ALPHA: PROCEDURE (BETA, GAMMA);
    DECLARE BETA FIXED, GAMMA CHARACTER;
        .
        .   BETA = BETA + 1;
        .
    END ALPHA;

```

becomes in HAL:

```

ALPHA: PROCEDURE (BETA_PRIME, GAMMA);
    DECLARE INTEGER, BETA, BETA_PRIME;
    DECLARE GAMMA CHARACTER (255) VARYING;
    BETA = BETA_PRIME;
        .
        .   BETA = BETA + 1;
        .
    CLOSE ALPHA;

```

when BETA is used as an assigned variable in the procedure, whereas GAMMA is not. The HAL compiler itself can be used to detect such occurrences, since assignments to parameters will be flagged as errors, significantly reducing the amount of program scanning necessary.

Note: Notice that the word CLOSE was used on the last line of the sample rather than the standard END. When closing a function or procedure in HAL, the word CLOSE is substituted for END. END is only used to signify the end of the DO loop or a DO case.

5.2.2.6 DO Statements. The DO case statement in HAL is similar to that in XPL, the only difference being that the first group of statements are executed when the DO case argument is equal to 0 in XPL and the first group of statements in HAL are executed when this argument is equal to 1. The following XPL sequence:

```

DO CASE I;
DO; /*CASE0*/
...
END; /*CASE0*/
...
END; /*OF DO CASE*/

```

translates to the following HAL sequence:

```

DO CASE I + 1;
DO; /*CASE1*/
...
END; /*OF CASE 1*/
...
END; /*OF DO CASE*/

```

When translating a DO case group from XPL to HAL, 1 must be added to the argument of the DO case statement rather than to change the value of the variable itself. DO case statements are translated in this manner to preclude the possibility of causing errors elsewhere in the compiler. It is not really possible to be certain that the change of the variable's value might not cause problems elsewhere. The looping statement:

```
DO IX = 1 to 10; (XPL)
```

simply becomes:

```
DO FOR IX = 1 to 10; (HAL)
```

The word FOR is required to distinguish this type of DO statement from the DO CASE or DO WHILE statements.

5.2.2.7 INPUT/OUTPUT. The primary input/output statements in XPL are the INPUT and OUTPUT pseudo-variables. To read a card image, the following statement is used:

```
CARD_IMAGE = INPUT;
```

Similarly, to write a line the following statement is used:

```
OUTPUT = NEXT_OUTPUT_LINE;
```

Both pseudo-variables are character string type and imply a new input/output record on each occurrence.

The corresponding HAL statements to read the same card image and print the same line are as follows:

```
READALL(5) CARD_IMAGE;  
  
WRITE(6) NEXT_OUTPUT_LINE;
```

Note that READALL, not READ, is used for input, as this forces reading an entire card image. READ into a character variable stops at any legal input delimiter.

5.2.2.8 Format of HAL File. The format of the input file is for the most part free of conventions. The only exception to this is that Column 1 may only be used to contain special letters. The following letters may appear in Column 1: C, D, E, M, and S.

These letters specify what type of line is contained on that current image. The letter C is to specify that the following text is to be treated as a comment and not actually compiled. D is used only for special compiler directives such as an INCLUDE file specified on this line. The letter E constitutes an exponent line which is part of the multi-line input format which HAL offers. M specifies that the following is the main line of the multi-line input, and S specifies a subscript line again which is part of the multi-line input. When using the single line format of HAL input, the M may be omitted from the line as long as text begins in Column 2 or after. The letter M is assumed on all lines which do not contain a character in Column 1. The above exception is the only one which pertains to the format of a HAL program.

Long and complicated HAL statements may be continued over as many cards as necessary just as in XPL. Certain equations which are broken up into several steps in the XPL version may be condensed into one large equation in HAL, resulting in a savings of temporary variables. (This is because XPL limits the number of expression temporaries in a statement to three registers. HAL has no such restriction.)

### 5.2.3 Debugging

The debugging procedure can be made quite simple by the use of various options which may be specified when compiling a HAL program. One may specify toggles on comment lines in HAL, which produce an identifier trace, a listing of the HALMAT code produced, and a list of the Fortran produced from Phase 2.

When errors occur, it is easy to trace the problem by consulting the three listings as mentioned above. Also, a check of the cross-reference listing produced greatly speeds debugging time since it is possible to determine in which statement a variable is either referenced, declared, or set. When a new section of code is added, a toggle can be set in a comment line at the beginning of the HAL program, which disables the call to Phase 2 of the HAL compiler. This is done to save computer time, since Phase 1 could perform a syntax check. When all syntax errors are eliminated, Phase 2 could then be called and Fortran output could be produced and subsequently compiled by the Fortran compiler.

It is, of course, much easier to debug a higher level language program than to debug assembly code, since ideas are clearly specified by the code being read, whereas in assembly language the intent is not always quite clear. In fact, when translating the in-line code it was sometimes necessary to speak to the person who had originally coded that section before a clear understanding could be gotten in order that the translation could be performed.

#### 5.2.4 Conclusions

At the time of this writing, some HAL features are still unimplemented. Because of this, certain sections of the translated code have as of now not been tried or debugged. However, that code which has been debugged and executed seems to prove that HAL is a language with which a large compiler can be easily written and debugged. The fact that HAL implements floating point arithmetic also eliminated a great deal of the complicated code necessary in the original XPL version. This fact alone made readability of the final copy much easier than the complicated in-line code which appears in the corresponding sections of the original copy.

## 5.3 Feasibility Demonstration ("HAL-in-HAL")

### 5.3.1 Objective

The HAL-in-HAL program was written as an experiment to prove whether or not the HAL language was suitable for writing translator systems, as well as aerospace applications. The program consists of a rewrite of the two-dimensional read routines originally coded in XPL for the HAL/360 compiler, utilizing the conversion techniques outlined in Section 5.2 above. These routines represent a full test of the character and bit manipulation facilities normally required for translator and system coding.

### 5.3.2 Test Program Description

The test program consists of an elementary scanning routine which utilizes the STREAM procedure for receiving its character-by-character input. STREAM converts the two-dimensional HAL input cards into the corresponding one-line format which is required by the scanner and subsequently the lexical analyzer. The test scanner repeatedly calls STREAM building-up identifier and numeric strings as tokens, as well as treating any special character as an automatic token. These are printed out as they are encountered. The test scanner is concurrently building-up an output line image which is a reflection of the input character received from STREAM. Whenever a semi-colon (;) is encountered, the current statement line, along with its corresponding over-punch markers, is printed, showing what the one-line format of the HAL statements looks like. A question mark (?) is used to indicate the end of the input stream for the purposes of this test. See Figure 5.3 for a flow chart of the test scanner. (Program listing - Statements 544-582 in Appendix B.)

### 5.3.3 Results

The HAL-in-HAL experiment has proved conclusively that HAL can be used successfully as a compiler implementation tool. Although HAL has no machine dependent features, (which frequently are designed into implementation languages), this experiment has proved that such features are not a requirement for compiler implementation, but rather merely a convenience item to circumvent known code generation inadequacies in the compiler. The HAL implementation is concise, readily followed, and understandable (even more so than the XPL version of the same program).

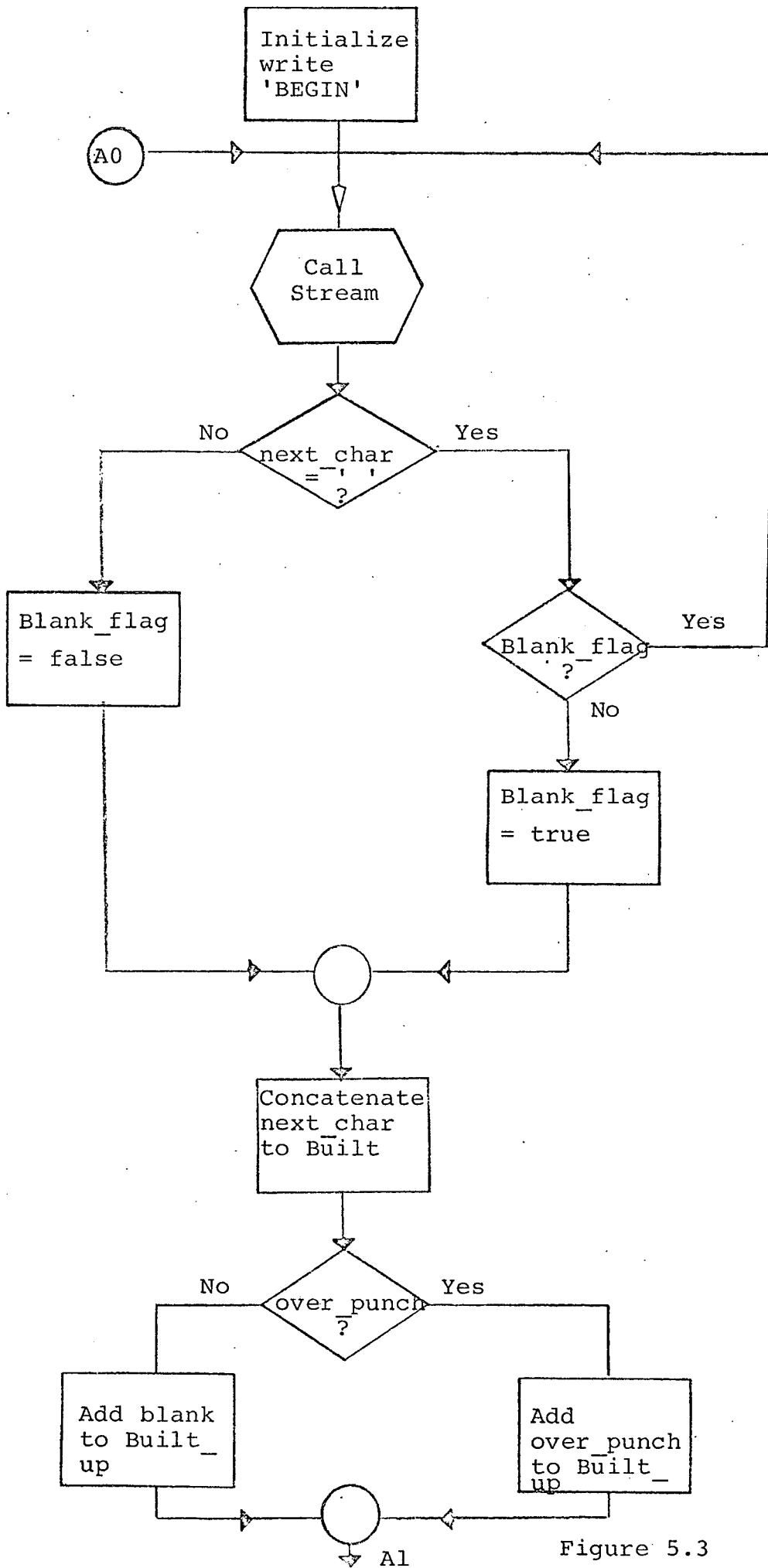


Figure 5.3

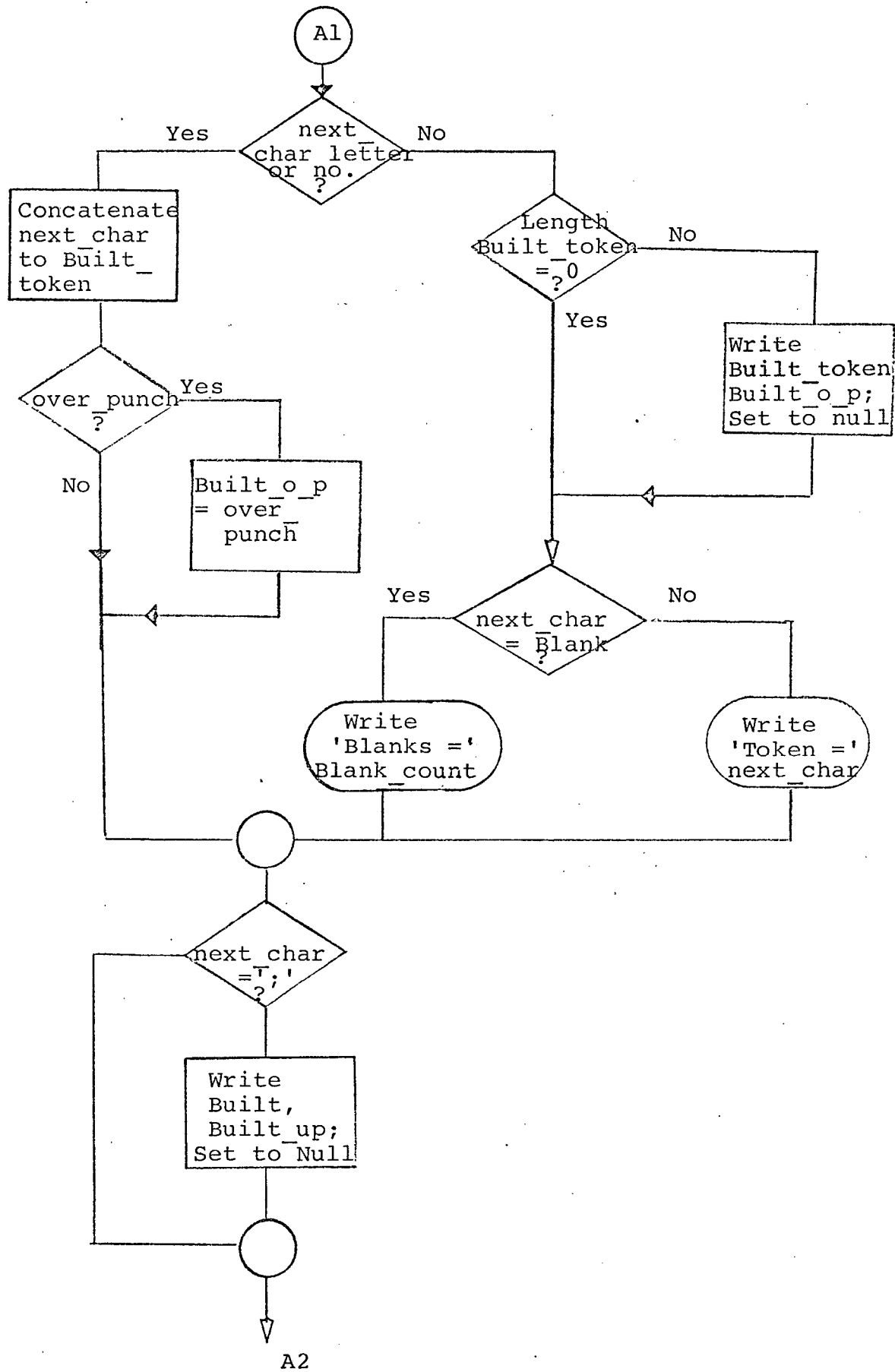


Figure 5.3 (Cont.)



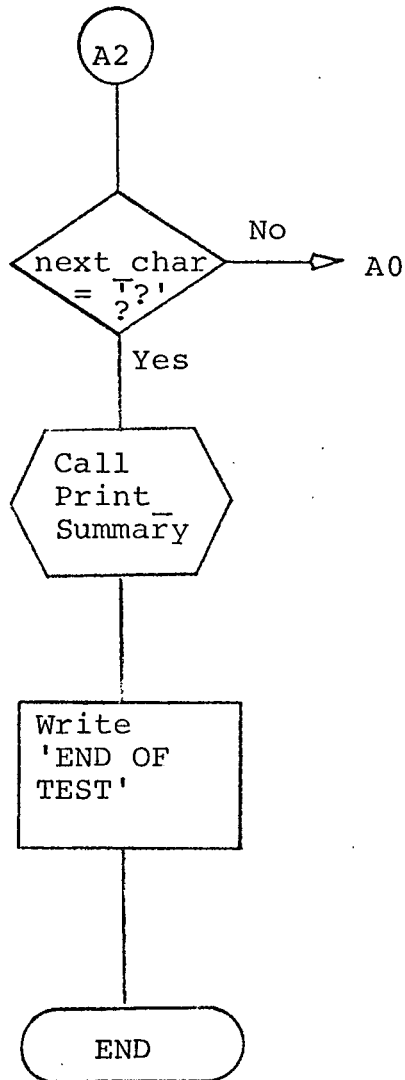


Figure 5.3 (continued)

The checking facilities of the HAL compiler can detect numerous logical errors without having to impose the strict definition rules of XPL. Uninitialized variables are easily detected, as are parameter mis-matches. The bulk of the debugging time for HAL-in-HAL was in streamlining the program to make the HAL version more readable, as well as more efficient, since the rule of adding one to all XPL subscripts as a general rule turned out to be both awkward and confusing in many instances. The final version of STREAM is much more efficient than the original translation performed utilizing the rules of Section 5.2. Programs originally coded in HAL will obviously not experience this problem.

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APPENDIX A.

HAL Course Material

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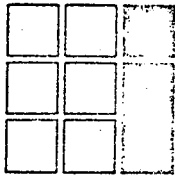
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A.1 OVERVIEW

## PREFACE

- \* HAL developed by Intermetrics, Inc.
  - \* Language design
  - \* Compiler design and implementation
  
- \* Significant Objectives
  - \* Increased readability
  - \* Increased reliability
  - \* Real time control
  
- \* Capabilities
  - \* Primarily designed for on-board computer
  - \* General enough for:

ground support and verification  
other real-time applications



SLIDE 1

**INTERMETRICS**

## SHUTTLE LANGUAGE REQUIREMENTS

### \* Software Applications

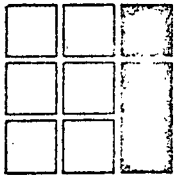
- \* Navigation, guidance, targeting
- \* Vehicle control
- \* Operating systems
- \* Data management
- \* Communications and displays
- \* Support software
- \* On-board checkout and monitor

### \* Computer Environment

- \* Wide range of computers (Flight and Ground)
- \* Fixed- and floating-point
- \* Simplex, multi-computer, multi-processor

### \* Language Characteristics

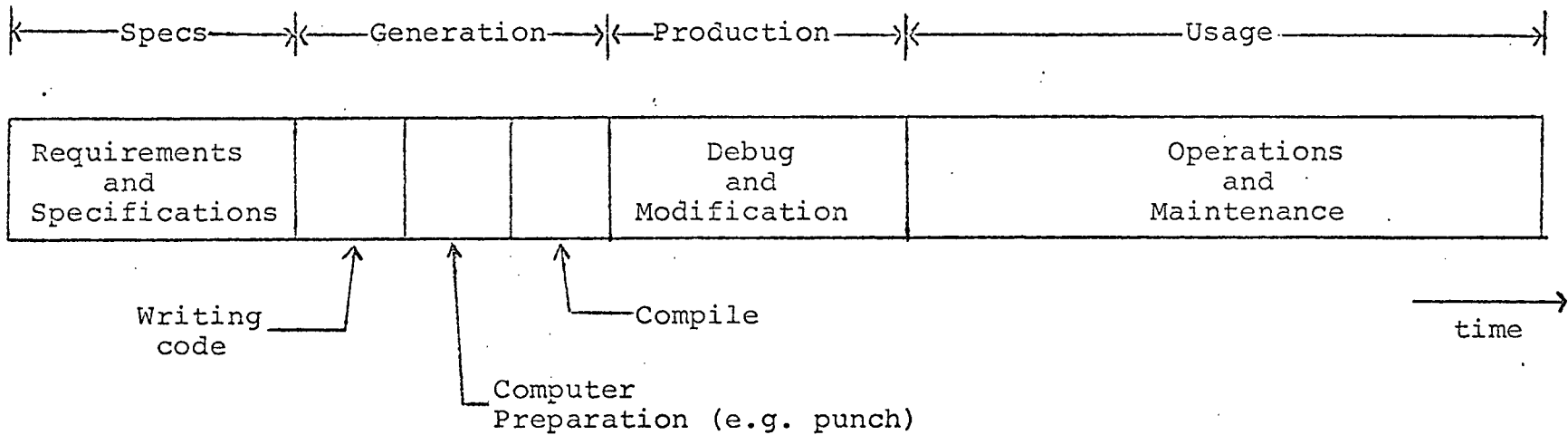
- \* Clarity and readability
- \* Enforcement of standards and conventions
- \* Extensive automatic checking (compile- and run-time)
- \* Facilitate software management
- \* Promote modularization



SLIDE 2

**INTERMETRICS**

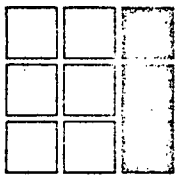
## CHRONOLOGY OF SOFTWARE DEVELOPMENT



46

### Some Observations

1. The writing of code is closely tied to the specifications.
2. The time required for computer preparation is small compared to the program life.
3. A lengthy period of debug and modification must be provided.
4. Period of program usage extends many times that of program generation.
5. Many more people will use a program than generated it.



### Conclusion

The computer language should promote understanding of the software. The listing should tend toward self-documentation.

**INTERMETRICS**

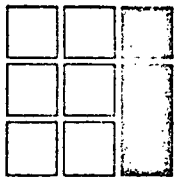
## SALIENT FEATURES OF HAL

### Capability

1. Two-dimensional Input-Output  
Annotation of variables
2. Complete vector-matrix arithmetic
3. Data array and structure handling
4. Bit and character manipulations
5. Real-time control statements
6. Data-Pool (COMPOOL), controlled sharing and name scope

### Requirement

- \* Increased readability
- \* Targeting, guidance and control
- \* Data management
- \* Systems, communications and I/O
- \* Command and control
- \* Increased reliability



**INTERMETRICS**

SLIDE 4

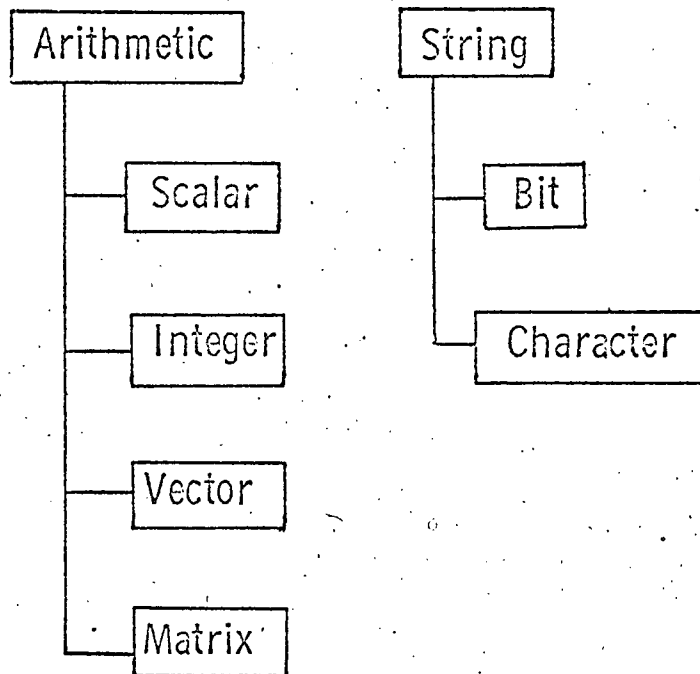


## ADVANCED FEATURES

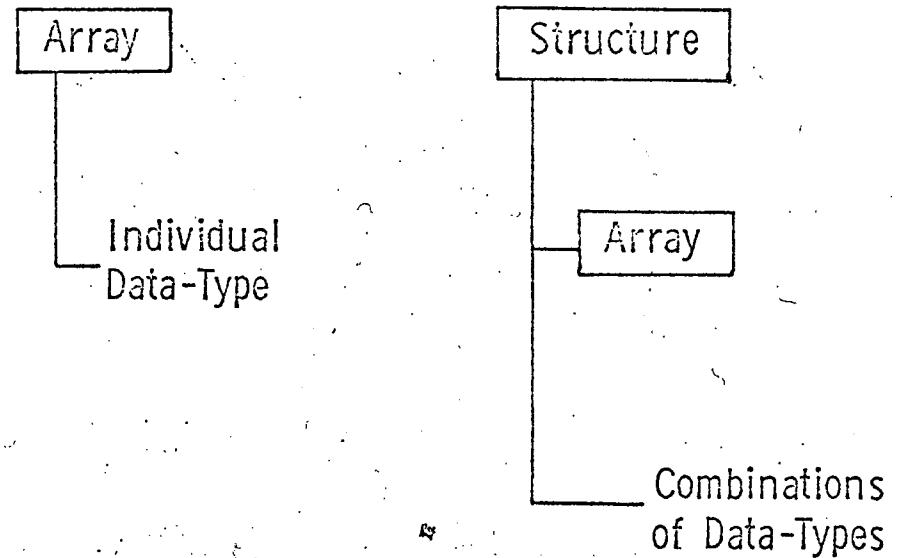
- TWO-DIMENSIONAL (MULTI-LINE) INPUT AND OUTPUT
- VECTOR AND MATRIX DATA TYPES AND OPERATORS
- TASK SCHEDULING AND SYNCHRONIZATION STATEMENTS FOR REAL-TIME CONTROL
- CONTROLLED SHARING OF DATA AMONG MULTIPLE USERS THROUGH A COMPOOL AND DATA LOCKING STATEMENTS
- STATEMENTS TO MANIPULATE DATA-GROUPS (ARRAYS AND STRUCTURES) AND POWERFUL METHODS TO PARTITION AND INDEX THEM
- OUTPUT ORIENTED. LANGUAGE IS SLANTED TOWARDS PRODUCTION OF UNDERSTANDABLE AND UNAMBIGUOUS OUTPUT LISTING RATHER THAN MINIMIZING KEYSTROKES ON INPUT
- A SIMPLE SCIENTIFIC SUBSET IS DEFINED AT THE OUTSET THAT WILL PERMIT EASY USE BY THOSE WITH A SCIENTIFIC BACKGROUND

## HAL Data Types and Organizations

### Types



### Organizations



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Unique notation: VECTOR:NAME, BIT STRING:NAME, CHARACTER STRING:NAME,  
MATRIX:NAME, ARRAY:[NAME], STRUCTURE:{NAME}

## EXAMPLES OF DATA TYPES

SCALAR: -126.04

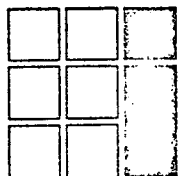
INTEGER: 126

VECTOR: 5, -26.4, 3.061

MATRIX: 5, -26.4, 3.061,  
1, -67.2, 106.1,  
0, 73.29, 0.06

BIT STRING: 1 0 1 1 0 1 0 1

CHARACTER STRING: VOLTAGE ON BATTERY B 2 VOLTS  
BELOW SPEC

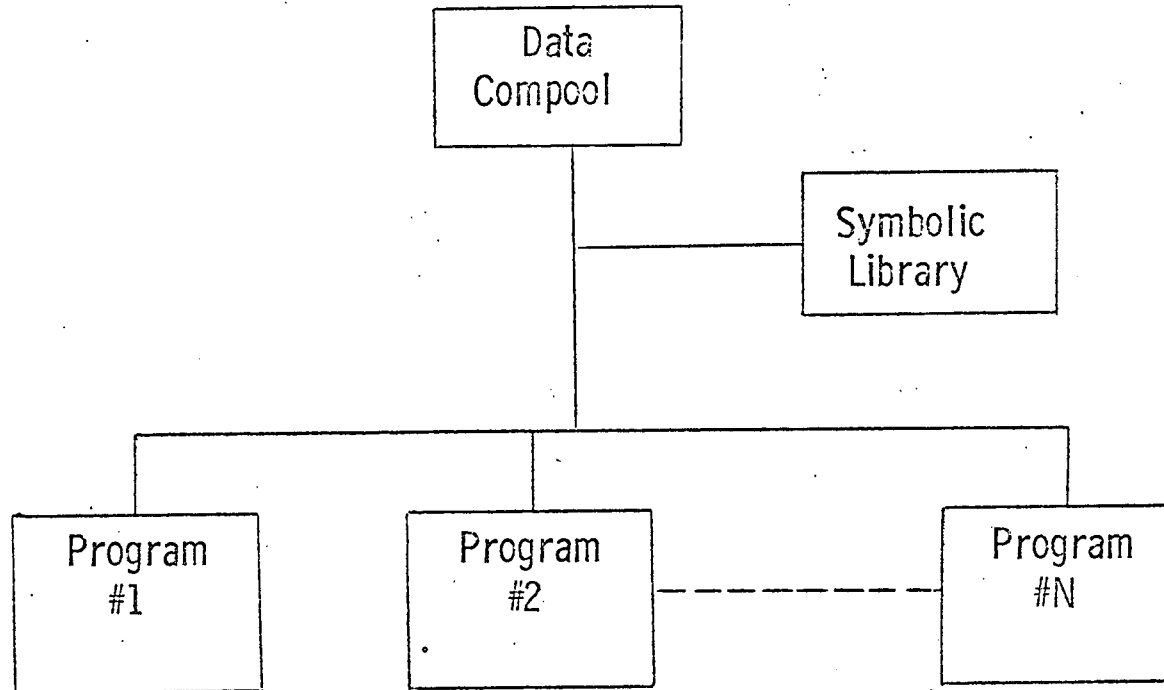


INTERMETRICS

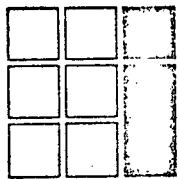
## STRUCTURE ORGANIZATION OF DATA

```
DECLARE 1 NAV_STATE (2);  
      2 STATE (2),  
        3 TIME PRECISION (8),  
        3 R VECTOR PRECISION (10),  
        3 V VECTOR PRECISION (10),  
      2 STATE_FLAGS,  
        3 BODY_FLAG BIT INITIAL (TRUE),  
        3 PHASE_FLAG BIT,  
      2 W MATRIX (9, 9) PRECISION (10);
```

# HAL PROGRAM ORGANIZATION



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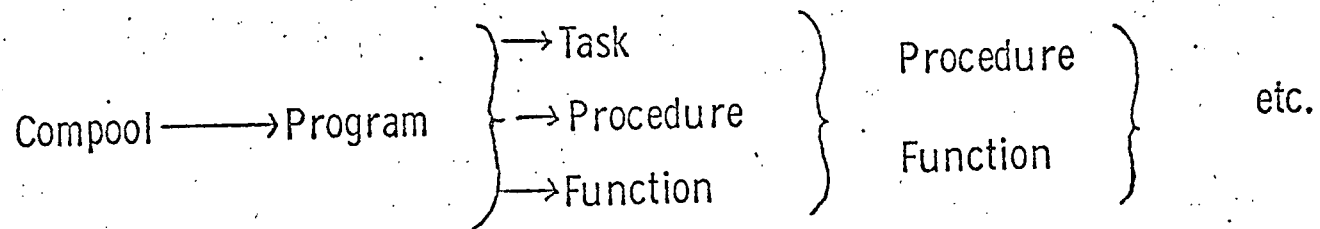


SLIDE 6

**INTERMETRICS**

## SCOPE OF NAMES

- Scope is the region in which a name is recognized.
- Scopes are defined from the outermost block toward the inner; i.e.,



- Names defined in an inner block are never recognized in an outer block. Inner blocks effectively isolate locally defined variables.

BLOCKS OF CODE (NAME SCOPE)

ABLE:

PROGRAM;

DECLARE VECTOR (5) A,B,C;

$\bar{A} = \bar{B} + \bar{C};$

⋮

BAKER:

TASK;

DECLARE A INTEGER;

CHARLIE:

PROCEDURE;

DECLARE X;

DECLARE A BIT (10);

⋮

END CHARLIE;

END BAKER;

GRAB:

PROCEDURE;

DECLARE X VECTOR (4);

⋮

END GRAB;

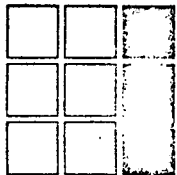
END ABLE;

→ A,B,C are vectors (5)

→ B,C are vectors (5)  
A is now an integer

→ B,C are vectors (5)  
A is now a bit string  
X is a scalar

→ A,B,C are vectors (5)  
X is a vector (4)



INTERMETRICS

## HAL Statements

### 1. Assignment

LABEL:VARIABLE = EXPRESSION;

### 2. Declare

DECLARE -----

### 3. Control

GO TO ----, IF-statements, DO-statements,

### 4. Block

Procedures, Functions, Tasks, Updates, Programs

### 5. Real-time Control

Schedules, Waits, Signals, Locks



EXAMPLES OF ARITHMETIC OPERATIONS

(From Apollo Navigation Equations)

HAL

$$\bar{Z} = \bar{W}^{*T} \bar{E};$$

$$\bar{\Omega} = \bar{Z} \bar{W}^{*T} / (ZMAG^2 + ALPHA^2);$$

$$D\bar{E}LX = \bar{\Omega} DELQ;$$

$$\bar{X} = \bar{X}' + D\bar{E}LX;$$

$$F = 1 + (ALPHA^2 / (ZMAG^2 + ALPHA^2))^{1/2};$$

$$\bar{W}^* = \bar{W}' - \bar{\Omega} \bar{Z} / F;$$

GSOP Specification

$$\underline{Z} = \underline{W}'^T \underline{b}$$

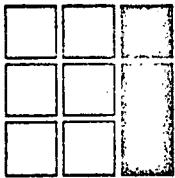
$$\underline{\omega}^T = \frac{1}{Z^2 + \bar{\alpha}^2} \underline{Z}^T \underline{W}'^T$$

$$\delta \underline{X} = \underline{\omega} \delta Q$$

$$\underline{X} = \underline{X}' + \delta \underline{X}$$

$$\underline{W} = \underline{W}' - \frac{\underline{\omega} \underline{Z}^T}{1 + \sqrt{\frac{\bar{\alpha}^2}{Z^2 + \bar{\alpha}^2}}}$$

where  $\underline{b} \equiv$  geometry vector  
 $\underline{W} \equiv$  square root of covariance  
 $\bar{\alpha}^2 \equiv$  measurement variance  
 $\underline{X} \equiv$  state vector



**INTERMETRICS**

SLIDE 9

CONTROL, LOGIC AND COMPUTATION  
(Cross product steering of Apollo vehicle)

Involves scalars, 3-d vectors, 3x3 matrices, "Booleans"

```
XSTEER:  IF TGO < 4 THEN DO;
           $\bar{\Omega}_{CNB} = 0;$ 
          SW = OFF;
          SCHEDULE ENGINE_OFF AT (TIME+TGO)
          PRIORITY (20)  E_OFF_ID;
          GO TO START;
          END;
```

$$\bar{DELM} = C \bar{B} DELT - \bar{DELV};$$

$$\bar{\Omega}_C = K(\bar{VG} * \bar{DELM}) / (ABVAL(\bar{VG}) ABVAL(\bar{DELM}));$$

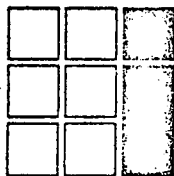
$$\bar{\Omega}_{CNB} = \bar{S}MNB \bar{R}EFSMMAT \bar{\Omega}_C;$$

GO TO START;

where TGO  $\equiv$  "time-to-go"

$\bar{VG} \equiv$  "velocity-to-be-gained"

$\bar{\Omega}_C \equiv$  rate command



**INTERMETRICS**

SLIDE 10

## EXAMPLES OF MATRIX PARTITIONING

Given: 9x9 covariance matrix E of errors in position, velocity and landmark location. That is,

$${}^*E = \begin{bmatrix} {}^*E_{p-p} & {}^*E_{p-v} & {}^*E_{p-l} \\ {}^*E_{v-p} & {}^*E_{v-v} & {}^*E_{v-l} \\ {}^*E_{l-p} & {}^*E_{l-v} & {}^*E_{l-l} \end{bmatrix}$$

### 1. RMS Errors

$$\text{RMS\_POS} = \text{SQRT}(\text{TRACE}({}^*E_{1 \text{ TO } 3, 1 \text{ TO } 3}));$$

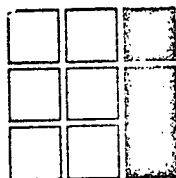
$$\text{RMS\_VEL} = \text{SQRT}(\text{TRACE}({}^*E_{4 \text{ TO } 6, 4 \text{ TO } 6}));$$

### 2. Initialize ${}^*E$ for new landmark

$${}^*E_{1 \text{ TO } 6, 7 \text{ TO } 9} = 0;$$

$${}^*E_{7 \text{ TO } 9, 1 \text{ TO } 6} = 0;$$

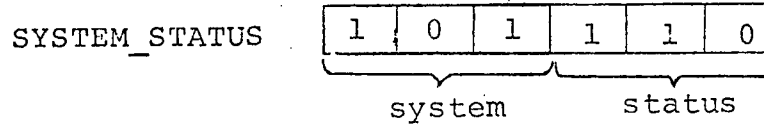
$${}^*E_{7 \text{ TO } 9, 7 \text{ TO } 9} = \text{MATRIX}_{3,3} \begin{pmatrix} A^2, & 0, & 0 \\ 0, & B^2, & 0 \\ 0, & 0, & C^2 \end{pmatrix};$$



**INTERMETRICS**

BIT AND CHARACTER MANIPULATIONS

Suppose the system-status word is made up as follows:



Example:

```

      A = 'SYSTEM STATUS: ';
DECODE: DO CASE SYSTEM_STATUS_1 TO 3;

      MESSAGE = 'ENGINE' || A;           CASE 1
      MESSAGE = 'POWER' || A;           CASE 2
      MESSAGE = 'IMU' || A;             CASE 3
      MESSAGE = 'LIFE_SUPPORT' || A;     CASE 4

      *                                     **
      *                                     **

      END;

      DO CASE SYSTEM_STATUS_4 TO 6;

      MESSAGE = MESSAGE || 'O.K.';       CASE 1
      MESSAGE = MESSAGE || 'RECONFIGURED'; CASE 2
      MESSAGE = MESSAGE || 'IN SELF-CHECK'; CASE 3

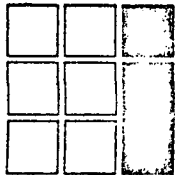
      *                                     **
      *                                     **

      END;

END_DECODE: WRITE(DISPLAY) MESSAGE;

```

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**INTERMETRICS**

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EXAMPLE OF IF-STATEMENTS  
(Flag-checking in Apollo Rendezvous Data Processing)

A: WAIT FOR SYNCH\_SIGNAL;  
IF REFSMMAT\_FLAG THEN  
    IF R\_60\_OP THEN GO TO A;  
    ELSE IF UPDATE\_FLAG THEN DO;  
        IF VHF\_RANGE THEN  
            IF TIME>60-TIME\_VHF THEN GO TO VHFREAD;  
            GO TO D;  
        END;  
    ELSE IF TRACKFLAG THEN GO TO D;  
GO TO EXIT;

Note: ELSE always refers to immediately preceding IF (except when IF is within a DO group)

EXAMPLE OF A PROCEDURE

(The Apollo Time-Radius Routine  
from GSOP)

CALLER:

```
CALL TIME_RADIUS ( $\bar{RT}2$ ,  $\bar{VT}2$ , (ABVAL ( $\bar{RT}2$ ) - 30480), MU_EARTH,  
T_R_FLAG); ASSIGN (TIME_32,  $\bar{RT}3$ ,  $\bar{VT}3$ );
```

SUBROUTINE:

```
TIME_RADIUS: PROCEDURE ( $\bar{A}$ ,  $\bar{B}$ , C, D,  $\bar{E}$ ) ASSIGN (F,  $\bar{G}$ ,  $\bar{H}$ );
```

```
[Statements  
|  
|  
|  
|  
|
```

```
RETURN;
```

```
END TIME_RADIUS;
```

NOTE: "Call-By-Name", "Call-By-Value"

EXAMPLE OF A FUNCTION

```
ABLE:  N = TRACER(A+B);  
        GO TO BAKER;  
TRACER: FUNCTION(Q);  
        DECLARE Q MATRIX(A,*);  
        IF TRACE(Q) > 100 THEN  
            RETURN (Q Q-1 + Q + QQ + QQQ)  
        ELSE RETURN(0);  
END TRACER;
```

NOTE: "CALL BY VALUE", "run-time" dimensions

# PROGRAMMING REQUIREMENTS FOR REAL TIME SPACE APPLICATIONS

## Scheduling and Tasking

Software performs time critical functions and responds to interrupts in a complex environment requiring the capability to schedule, control and synchronize tasks.

## Recovery From Error Conditions

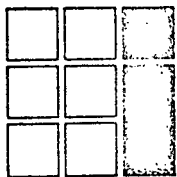
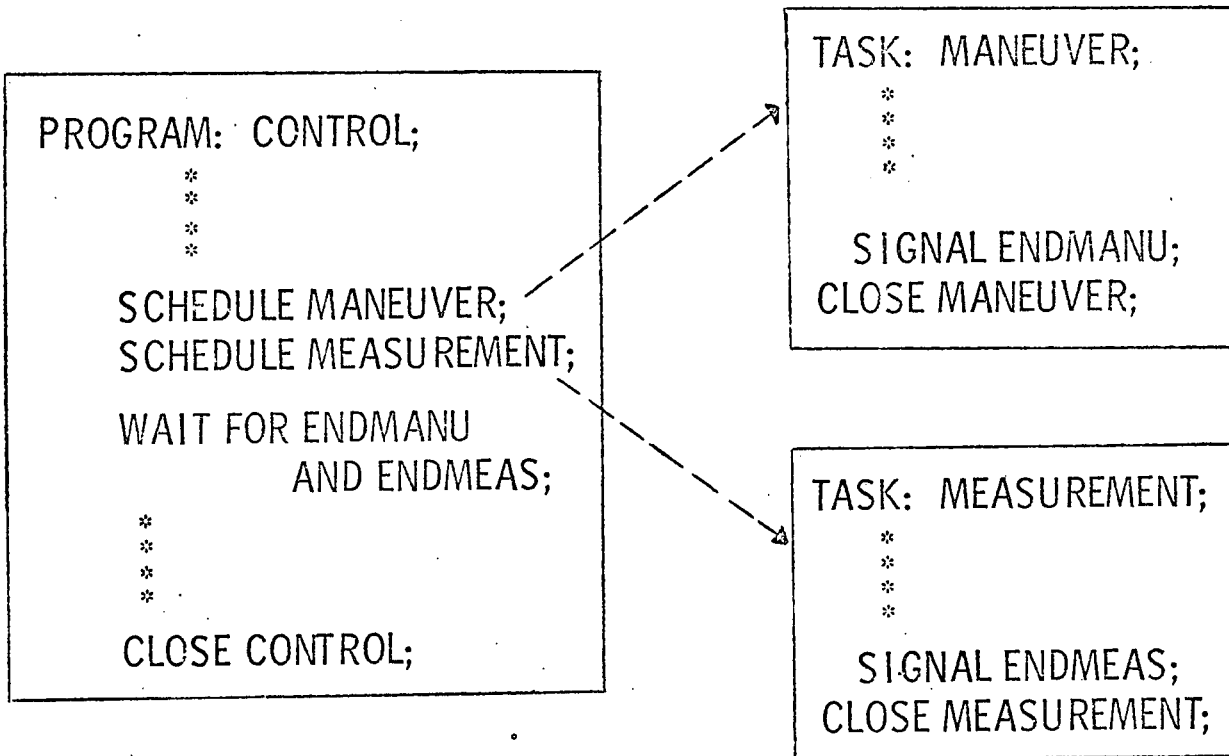
Techniques are required to protect and enable system to "continue" after detection of unexpected error condition.

## Common Memory Sharing and Control

Techniques are required to dynamically control the use of common data elements among tasks in the environment.



EXAMPLE OF HAL REAL-TIME CONTROL



ENDMANU and ENDMEAS are programmer-defined events

**INTERMETRICS**

## REAL TIME STATEMENT EXAMPLES

SCHEDULE TARGETING PRIORITY(3);

SCHEDULE RADAR ON R\_RUPT PRIORITY(PRIO + 2) RADAR\_PROG;

IF TRACKFLAG = ON THEN SCHEDULE AUTOMANEUVER IN 5;

ELSE WAIT UNTIL (TIME + 5);

SCHEDULE STEERING AT(IGNITION + 3)PRIORITY(10) INDEPENDENT;

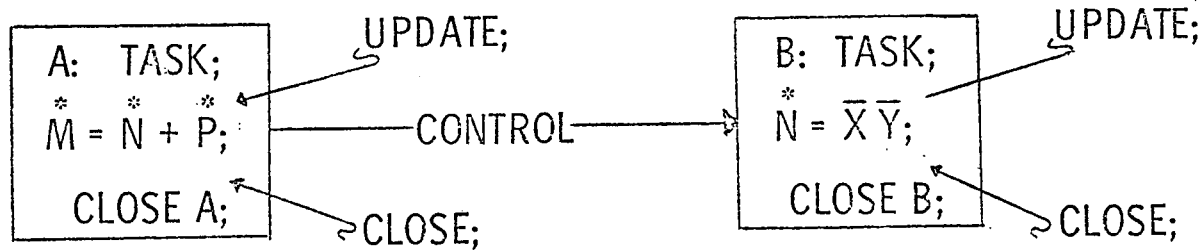
TERMINATE RADAR\_PROG;

WAIT FOR OK;

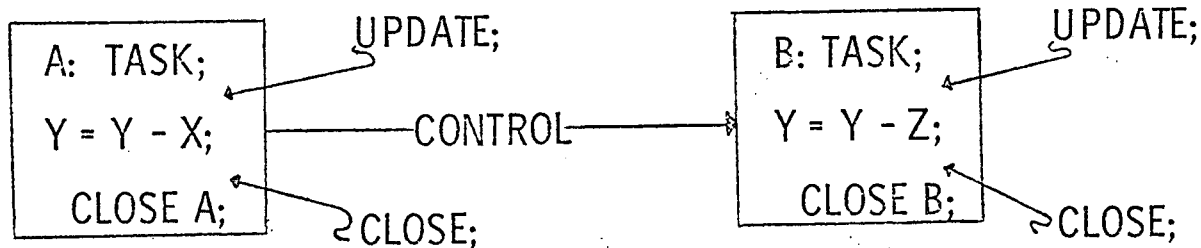
SIGNAL OK;

## CONTROL OF SHARED DATA

### EXAMPLE 1: READ AND WRITE CONFLICTS

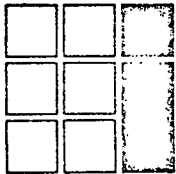


### EXAMPLE 2: UPDATE CONFLICTS



### NOTES:

1. B "INTERRUPTS" A IN BOTH CASES
2. #1 TASK A RESUMES USING OLD AND NEW VALUES FOR N\*
3. #2 TASK A RESUMES "CLOBBERING" THE VALUE FOR Y SET BY TASK B



ERROR CONDITION STATEMENTS EXAMPLES :

ON ERROR<sub>12</sub> GO TO ABLE;

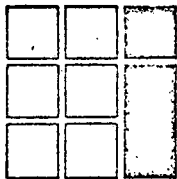
ON ERROR<sub>1 TO 5</sub> GO TO BAKER;

ON ERROR<sub>4</sub> SYSTEM;

E\_RUPT ERROR<sub>6</sub>;

## SUMMARY

- \* HAL emphasizes reliability
  - \* Readability
  - \* Data protection
  
- \* HAL is a full-capability language
  - \* Includes all data types
  - \* Real-time control statements
  - \* Supports on-board computer software
  - \* Floating- or fixed-point syntax
  - \* Supports ground, checkout, simulation software
  
- \* Schedule of Events
  - \* First version delivery to MSC in June, 1971
  - \* Development to continue compatible with Shuttle schedule
  - \* Object-code-module required for selected on-board computer



SLIDE 14

INTERMETRICS

A.2 LONGER HAL COURSE

## HAL Data Operations - I

### Arithmetic

All common operations including:

- \* Vector dot, cross and outer - products
- \* Matrix multiplication, inverse transpose
- \* Integer mathematics
- \* Combined integer-scalar operation

### Bit String

Logical AND, OR, NOT, of:

- \* Long bit strings, bit-by-bit
- \* Single bit "booleans"

Concatenation

### Character String

Concatenation of characters and data into messages

### Arrays

Most valid element-by-element operations apply to arrays

## HAL Data Operations - II

### Comparisons

A comparison of data always results in a single TRUE or FALSE answer.

#### \* Absolute comparison

VECTOR, MATRIX, ARRAY, STRUCTURE

=  
NOT=

#### \* Absolute and Relative Comparisons

scalar

integer

bit

character

=  
NOT=  
<=  
>=  
NOT <  
NOT >  
>  
<



EXAMPLES OF DATA OPERATIONS - I

Arithmetic Operations

$$I = (-J)^3 + K(3+J); \quad \text{---- all integers}$$

$$A = P/R + R^P - \bar{V} \cdot \bar{M}^* \bar{V}; \quad \text{---- scalars, vectors, matrices}$$

$$\bar{B} = -\bar{P}^* \bar{V} + 5 \bar{M}^* \bar{V} + A(\bar{V} \cdot \bar{P})^2 \bar{F} + \bar{F}/(\bar{V} \cdot \bar{V} \bar{M}^*); \quad \text{--- scalars, vectors, matrices}$$

$$\bar{C} = -\bar{M}^* \bar{N}^* + \bar{M}^{*-1} + \bar{V} \bar{V}/A - (\bar{M}^* \bar{N}^*)^T; \quad \text{---- scalars, vectors, matrices}$$

$$A = I/J + B J + J^{-6.4}; \quad \text{----- integers, scalars}$$

## EXAMPLES OF STRING OPERATIONS

### Bit String

$\dot{D} = \dot{B} \text{ AND } \dot{C};$

$\dot{D} = \dot{A} \text{ OR } (\dot{B} \text{ AND } \dot{C});$

$\dot{A} = \dot{D} \text{ || NOT } \dot{B} \text{ || } (\dot{B} \text{ AND } \dot{C});$

### Character String

$C = \text{'PLEASE'} \text{ || 'HELP'};$

$D = \text{'THE ANS. IS'} \text{ || X || 'N.M.'};$

EXAMPLES OF ARRAY OPERATIONS

"Two-array" operations:

$$-[P]/[A], [\bar{P}] * [\bar{V}], [R]^{[P]}, [\dot{A}] \text{ OR } [\dot{B}]$$

One-array operations:

$$-P/[A], \bar{P} * [\bar{V}], R^{[P]}, [\dot{A}] \text{ OR } \dot{B}$$

## EXAMPLES OF COMPARISON OPERATIONS

IF  $I > J$  THEN -----

IF  $A \text{ NOT } < (X^2 - 5 \bar{V} \cdot \bar{V})$  THEN -----

DO WHILE  $S = (A \text{ AND } B)$  -----

IF  $M^* = N^*$  THEN -----

DO WHILE  $[A] \text{ NOT } = [B]$  -----

## HAL Explicit Declarations

- \* In general, all data must be declared by declare statements (with the exception of those permitted by implicit declarations).
- \* The declare statement specifies the name, organization, type and attributes of data quantities.

Keywords used in declare statements:

<u>Organization</u>	<u>Type</u>	<u>Attributes</u>
ARRAY	INTEGER SCALAR (optional) VECTOR MATRIX BIT CHARACTER	PRECISION INITIAL CONSTANT STATIC AUTOMATIC LOCKTYPE DENSE ALIGNED VARYING

## EXAMPLES OF EXPLICIT DECLARATIONS

1. DECLARE I INTEGER INITIAL (65);

I is an integer with an initial value = 65.

2. DECLARE X PRECISION (8) AUTOMATIC INITIAL (6.061);

X is a floating point scalar with at least 8 significant decimal digits.

3. DECLARE A ARRAY (5, 3, 4) VECTOR (6) PRECISION (10);

A is a 5x3x4 array. Each element is a 6 dimensional vector with components represented to 10 significant decimal digits.

4. DECLARE MATRIX (3, 4) INITIAL (0) AUTOMATIC

A, B, C PRECISION (10);

A, B, and C are all (3, 4) matrices with automatic storage. All components are set to zero.

5. DECLARE A PRECISION (10, 15)

A is a fixed point scalar with 10 integer bits and at least 15 fractional bits (i. e. maximum value  $< 2^{10}$ , granularity  $\leq 2^{-15}$ ).

## INDEX OPERATORS

1. The "TO" operator

Selects a subset of elements from element-i  
"TO" element-j.

For example:

$A_1$  TO 10

2. The "AT" operator

Selects a subset of N-elements starting at  
element-i.

For example:

$A_{10}$  AT 1

3. The number of elements in any "partition"  
must be known at compile-time.

EXAMPLES OF INDEXING - I

1. Vectors and Matrices (given  $\bar{V}$ ,  $\bar{M}$ )

$V_2$  ----- scalar element,

$\bar{V}_1$  TO 5 ----- sub-vector,

$M_{2,3}$  ----- scalar element,

$\bar{M}_{*,1}$  ----- vector element,

$\bar{M}_{3 \text{ AT } P, 3 \text{ AT } Q}^*$  ----- sub-matrix

2. Bit and Character Strings (given  $\dot{S}$ ,  $C$ )

$\dot{S}_3$  ----- single bit,

$\dot{S}_2$  TO 10 ----- sub-string,

$\dot{S}_6$  AT P ----- sub-string

$C'_3$  TO # ----- sub-string,



## EXAMPLES OF INDEXING - II

3. Arrays (given  $[A]$ , a two-dimensional array of matrices)

$[A]_{1 \text{ TO } 4, 3 \text{ TO } 6}$  ----- sub-array

$[A]_{P, Q}$  ----- an array of scalar elements

$[\bar{A}]_{*, *, *, 2} \equiv [\bar{A}]_{*, 2}$  ----- sub-array of vector elements

4. Array of Bit Strings

$[A]_{3 \text{ TO } 5, *: 1 \text{ TO } 6}$  ----- sub-array of sub-strings

$A_{6, 4:3}$  ----- one particular bit

## DO - STATEMENTS

- DO - statements block out a set of statements which are to be treated as a single unit.
- There are four types of DO - statements
  1. Simple DO-END
  2. Iterative DO-FOR
  3. Iterative DO-WHILE
  4. Selective DO-CASE

EXAMPLES OF DO-STATEMENTS - ISimple DO-END

IF X>5 THEN BAKER: DO;

    B = A = B;

    C = D;

    GO TO ABLE:

END BAKER;

ELSE CHARLIE: DO;

$\bar{Z} = \bar{M} * \bar{V}$ ;

$\bar{F} = \bar{V} * \bar{Z}$ ;

    IF Y = 0 THEN GO TO OUT;

END CHARLIE;

EXAMPLES OF DO-STATEMENTS - II

1. Iterative DO-FOR

ABLE: DO FOR I = P TO (N/S) BY L WHILE N>0;

limits and increment  
are computed once.

$X = Y^2 + A_I;$

$N = N - .006 X;$

$P = 1; S = 2; L = 3;$

END ABLE;

2. Iterative DO-WHILE

ABLE: DO WHILE (X > Y AND GO\_FLAG = ON);

X, Y and GO\_FLAG  
are recomputed.

$X = Y^2 + P \cdot \text{LOG}(Z);$

GO\_FLAG = TRAKFLAG OR NAV\_FLAG;

[ Statements  
: :  
: :  
: : ]

END ABLE;

EXAMPLE: SEARCHING AN ARRAY OF DATA

The final phase reference for Apollo reentry:

```
DECLARE ARRAY (13) VREF CONSTANT (994, 2103, 3922,...);  
DECLARE ARRAY (13) RDOTREF CONSTANT (-690.0, -719,...);  
DECLARE ARRAY (13) DREFR CONSTANT (41.15, 60, 81.5,...);  
:  
:  
:  
etc.
```

```
INTERPOLATE:  I = 0;  
              DO WHILE (VREFI NOT< V) AND (V NOT< VREFI+1);  
              I = I + 1;  
              END;  
              GRAD = (V - VREFI) / (VREFI+1 - VREFI);  
              RDOTREFV = RDOTREFI + GRAD (RDOTREFI+1 - RDOTREFI);  
              GO TO CONTINUE;
```

EXAMPLES OF DO-STATEMENTS - III

Selective DO-CASE (Computed DO-Statement)

ABLE: DO CASE N;

$X = Y^2;$  CASE 1

BAKER: DO CASE P; CASE 2

$F = A + B;$  CASE 1

$\bar{G} = \bar{M} * \bar{V};$  CASE 2

END;

GO TO CHARLIE; CASE 3

$\bar{Z} = \bar{W} + \bar{B};$  CASE 4

END ABLE;

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EXAMPLES OF IF-STATEMENTS

1. Simple:

IF X = 5 AND Y > 6 THEN ABLE: GO TO PLACE;  
ELSE GO TO TRY\_AGAIN;

2. Complex:

IF X = 5 THEN IF Y > 6 THEN IF B OR C THEN  $\bar{Z} = M^* \bar{V}$ ;  
ELSE CHOICE:  $\bar{Z} = M^{*-1} \bar{V}$ ;

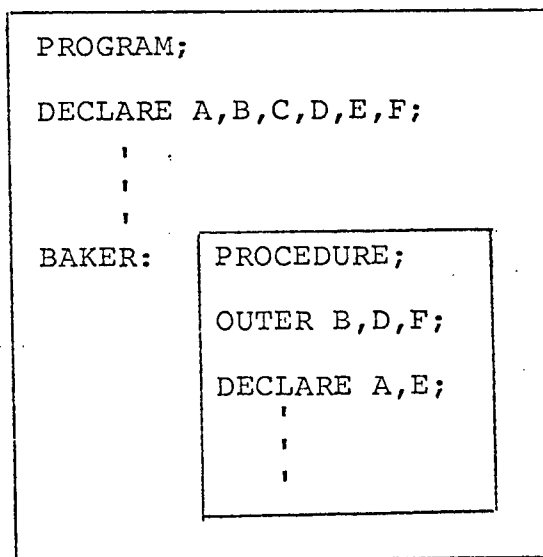
3. More Complex:

IF S = (A OR B) THEN IF X > 5 AND Y > 6 THEN GO TO OUT;  
ELSE IF [A] NOT= [B] THEN [A] = [C];  
ELSE IF -----THEN -----; ELSE IF ----- THEN -----;  
ELSE GO TO TRY\_AGAIN;

note: ELSE always refers to immediately preceding IF.

SELECTIVE INCLUSION OF OUTER-NAMES

ABLE;



- Only B,D,F are recognized outer names. (A,C,E, are "rejected").
- A,E, are defined locally.

Note: COMPOOL variable-names may be accepted, rejected and/or locally defined by combinations of DECLARE and OUTER statements. In order to use implicit declarations within a block (except for PROGRAM-level) an OUTER-statement must be present.



APPENDIX B.

HAL-in-HAL Detailed Description and Listing

B.1 INTRODUCTION

The output from the HAL-in-HAL experiment consists of the following:

1. The HAL program listing, symbol table, and cross reference
2. The output of the HAL program execution
3. A listing of the source data read by HAL-in-HAL (a sample HAL program)

B.2 HAL-in-HAL PROGRAM DESCRIPTION

This section will describe the general function of each of the routines in HAL-in-HAL. Refer to the program listing for specific details.

B.2.1 STREAM (Statements 100-542)

The overall functions of the STREAM procedure are as follows:

1. Convert the multi-line format of the input cards to a one-line format, which is required by the scanning and syntactic analysis routines.
2. Process Comment and Heading cards to aid readability of the source program, and also to enable certain compiler toggles for assisting the person who is debugging the compiler.
3. Eliminate HAL in-line comments (strings contained between /\* and \*/).
4. Perform substitutions for replace type strings (not demonstrated).

To convert input to one-line format requires the following:

1. Enclosure of each level of subscripting in parentheses, preceded by the dollar-sign character (\$).

Ex.

$$\begin{array}{c} M \\ S \\ S \end{array} S1_I = S2_{IX_I};$$

becomes

$$S1\$(I) = S2\$(IX\$(I));$$

2. Enclosure of each level of exponentiation in parentheses, preceded by two asterisks (\*\*).

Ex.

$$\begin{array}{c} E \\ E \\ M \end{array} S2^2 = X^{I^2};$$

becomes

$$S2^{**}(2) = X^{**}(I^{**}(2));$$

STREAM contains ten local subroutines which assist it in performing its function. They are local because they are of no value outside of STREAM. The non-local procedures are general interest routines, which are useful at levels other than within the STREAM procedure.

## B2.2 CARD\_TYPE (Statements 26-33)

CARD\_TYPE is a function which receives as input the first character from an input card and returns an integer typifying the card in one of five classes.

## B2.3 CHAR\_INDEX (Statements 34-44)

CHAR\_INDEX locates one string within another, returning the relative position of the desired substring if found, and 0 otherwise.

## B2.4 PAD (Statements 45-53)

The PAD function forces a varying character string to a minimum specified fixed length by appending trailing blanks.

Its primary function is for WRITE list items.

B2.5 I\_FORMAT (Statements 54-61)

The I\_FORMAT function first converts a number to a character string, and then adds high order blanks to force a right justified integer string of a specified fixed length. This routine is also primarily for WRITE list items.

B2.6 ERRORS (Statements 62-82)

ERRORS both prints error messages when reported, and saves a record of their occurrence for later reporting. It also will terminate the compilation if either too many or too severe errors occur during compilation.

B2.7 PROCESS\_COMMENT (Statements 116-133)

This routine processes heading cards, as well as looking for special debugging directives on comment cards.

B2.8 STACK\_RETURN\_CHAR (Statements 134-144)

This routine locates an available position in the return stack and records both a count and the character to be added to the output stream to formulate one-line output out of multi-line input.

B2.9 READ\_CARD (Statements 145-157)

This routine reads the next input card and prints the card previously read. (This is because a group is only defined by the next non-group card.) It also counts the input cards and checks for an end of input condition.

B2.10 ORDER\_OK (Statements 158-189)

This routine verifies that cards are in the proper sequence to formulate a proper HAL group. It also signals when a group is completed.

B2.11 COMMENT (Statements 190-201)

This routine removes '/'\* \*/' type comments from E and S lines, when they exist.

B2.12 SCAN\_CARD (Statements 202-220)

This routine scans E and S cards for non-blank characters, compressing multiple lines into one line, with an indicator recording which level the character appeared on. If an overlap occurs, the one closest to the M line is retained, and a diagnostic is issued.

B2.13 COMP (Statements 221-244)

This routine is called when either an E line or S line is first encountered in a group. It keeps reading cards and calling SCAN until an entire E or S group is compressed into a single line as described in B2.12

B2.14 GET\_GROUP (Statements 245-297)

This routine is called to assemble a complete group, which consists of an M line and one or more E and/or S lines, formulating as output a single E line, M line, and S line, with corresponding indicators.

B2.15 CHOP (Statements 298-306)

This routine advances the M line character index by 1, forcing a new group to be read when the M line termination is reached. The information concerning the last character on the previous card is retained.

B2.16 STACK (Statements 307-325)

This routine builds on Exponent or Subscript stack corresponding to a single blank field on an M line; i.e., those subscripts and/or exponents related to a specified variable or function. This is in preparation for outputting from STREAM within stacked Return characters (see B2.8).

## B2.17 BUILD\_XSCRIPTS (Statements 326-362)

This procedure invokes STACK for building both the E and S stacks for a single blank field, including M line comments as blank fields. Residue blanks for which no E or S line characters exist are treated as blank fields by the scanner (blank fields by the scanner (blank is a legal delimiter in HAL)).

### B.3 DESCRIPTION OF HAL-IN-HAL OUTPUT

The output of the HAL program is interpreted thusly. The program reads in a group of data cards. A group can consist of:

1. a single M line
2. one or more E lines followed by a single M line
3. a single M line followed by one or more S lines
4. one or more E lines followed by a single M line followed by one or more S lines

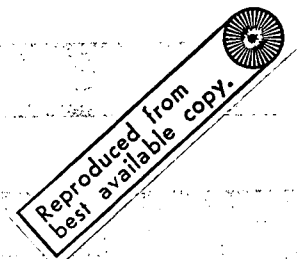
Each group is then converted into one line of output to the scanning routine. A group of type 1 is transmitted directly except for elimination of redundant blanks. A group of the other three types involves processing of S line and/or E line stacks and the addition of the appropriate subscript and/or exponent enclosures (subscripting is performed first if the M line identifier has both a subscript and an exponent attached). As the scanning routine continuously calls STREAM, identifier tokens are formed. All identifier and special character tokens are printed as they are encountered, including identifier over punches. When the token is a blank, a count of the blanks scanned is printed.

The full output consists of the following:

1. a printout of the entire card group just read, complete with card numbering
2. a printout of the individual tokens encountered within the group just read (signalled by TOKEN= or BLANKS=)
3. a printout of the combined one-line format and any possible over-punch characters whenever a semi-colon (;) is returned by STREAM. This represents a complete HAL statement as seen by the scanning routine (signalled by OVER and MAIN in succession).

This sequence is repeated until the input is exhausted (signalled by a "?" for this test). Note in the TOGGLE sequences where the subscript and exponent enclosures are added to the characters on the actual input cards.

STMT	SOURCE	LINE	CURRENT SCOPE
1 M	HALINHAL: PROGRAM:	1	
CI	A B C D E F G	2	
CI	THE FOLLOWING GROUP OF DECLARES ARE ONLY TEMPORARY	3	
2 M	REPLACE R BY 'REPLACE';	4	HALINHAL
3 M	R FALSE BY '0'; R TRUE BY '1'; R BIT_1 BY 'INTEGER';	5	HALINHAL
6 M	R IF_NCT BY 'IF 0 ='; R IFF BY 'IF 1 =';	6	HALINHAL
8 M	R CUT BY 'WRITE(6)';	7	HALINHAL
9 M	R FOREVER BY 'WHILE 1 = 1';	8	HALINHAL
10 M	DEFAULT INTEGER;	9	HALINHAL
11 M	DECLARE X70 CHARACTER(70) INITIAL (' ');	10	HALINHAL
12 M	DECLARE CHARACTER(255) VARYING, BUILT, BUILT_UP, BUILT_TOKEN, BUILT_O_P;	11	HALINHAL
13 M	DECLARE BLANK_FLAG BIT_1 INITIAL(FALSE);	12	HALINHAL
14 M	DECLARE INTEGER INITIAL(0),	13	HALINHAL
14 M	ERRCR_COUNT, MACRO_POINT, MACRO_LIMIT,	14	HALINHAL
14 M	OLD_LEVEL, NEW_LEVEL,	15	HALINHAL
14 M	MAX_SEVERITY, STATEMENT_SEVERITY,	16	HALINHAL
14 M	SAVE_SEVERITY ARRAY(100),	17	HALINHAL
14 M	SAVE_LINE ARRAY(100), PREVIOUS_ERROR, 1;	18	HALINHAL
15 M	DECLARE COMPILING BIT_1 INITIAL(TRUE);	19	HALINHAL
16 M	DECLARE INTEGER, TIME1, TIME2;	20	HALINHAL
17 M	DECLARE MACRO_STREAM CHARACTER(255) VARYING;	21	HALINHAL
18 M	DECLARE MACRO_FOUND BIT_1 INITIAL(FALSE);	22	HALINHAL
19 M	DECLARE CARD_COUNT INTEGER INITIAL(-1);	23	HALINHAL
20 M	DECLARE DISASTER LABEL;	24	HALINHAL
21 M	DECLARE NEXT_CHAR CHARACTER(1);	25	HALINHAL



96

3"

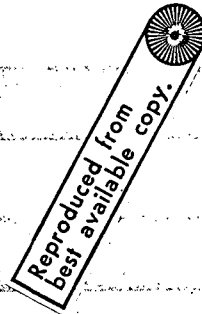
2"

1"

STMT	SOURCE	LINE	CURRENT SCOPE
22 MI	DECLARE OVER_PUNCH CHARACTER(1);	26	HALINHAL
23 MI	DECLARE BLANK_COUNT INTEGER INITIAL(0);	27	HALINHAL
24 MI	DECLARE CTRLCL ARRAY(10) BIT(32);	28	HALINHAL
25 MI	DECLARE TCGGLES CHARACTER(10) INITIAL('1234567890');	29	HALINHAL
CI		30	HALINHAL
CI		31	HALINHAL
CI	THE ARRAY CARD_TYPE IN XPL IS REPLACED BY THIS FUNCTION.	32	HALINHAL
CI		33	HALINHAL
26 MI	CARD_TYPE:	34	HALINHAL
26 MI	FUNCTION(SELECT) INTEGER;	35	CARD_TYPE
27 MI	DECLARE SELECT CHARACTER(1);	36	CARD_TYPE
28 MI	IF SELECT = 'M' OR SELECT = 'A' THEN RETURN 2;	37	CARD_TYPE
29 MI	IF SELECT = 'S' THEN RETURN 3;	38	CARD_TYPE
30 MI	IF SELECT = 'B' THEN RETURN 1;	39	CARD_TYPE
31 MI	IF SELECT = 'C' OR SELECT = 'D' OR SELECT = 'H'	40	CARD_TYPE
31 MI	THEN RETURN 4;	41	CARD_TYPE
32 MI	RETURN 0;	42	CARD_TYPE
33 MI	CLOSE CARD_TYPE;	43	HALINHAL
CI		44	HALINHAL
CI		45	HALINHAL
CI	THE CHAR_INDEX FUNCTION IS THE SAME AS THE INDEX BUILT IN FUNCTION	46	HALINHAL
CI		47	HALINHAL
34 MI	CHAR_INDEX: FUNCTION(String,PATTERN) INTEGER;	48	CHAR_INDEX
35 MI	DECLARE CHARACTER(*), STRING, PATTERN;	49	CHAR_INDEX
36 MI	DECLARE INTEGER, I, J, K;	50	CHAR_INDEX
37 MI	J = LENGTH(STRING);	51	CHAR_INDEX
38 MI	K = LENGTH(PATTERN);	52	CHAR_INDEX
39 MI	IF K > J THEN RETURN 0;	53	CHAR_INDEX
40 MI	ELSE DO FOR I = 1 TO J-K+1;	54	CHAR_INDEX
41 MI	IF STRING\$(K AT I) = PATTERN THEN RETURN I;	55	CHAR_INDEX
42 MI	END;	56	CHAR_INDEX



STMT	SOURCE	LINE	CURRENT SCOPE
43 M	RETURN 0;	57	CHAR_INDEX
44 M	CLOSE CHAR_INDEX;	58	HALINHAL
CI		59	HALINHAL
CI		60	HALINHAL
CI	HERE IS PAD A FUNCTION FROM HALPASS1	61	HALINHAL
CI		62	HALINHAL
45 M	PAD: FUNCTION (STRING,WIDTH) CHARACTER(255) VARYING;	63	PAD
46 M	DECLARE CHARACTER(*), STRING;	64	PAD
47 M	DECLARE CHARACTER(255) VARYING, TEMP_STRING;	65	PAD
48 M	DECLARE INTEGER, L, WIDTH;	66	PAD
49 M	L = LENGTH(STRING);	67	PAD
50 M	IF L < WIDTH THEN	68	PAD
50 M	TEMP_STRING = STRING    X70\$(1 TO WIDTH-L);	69	PAD
51 M	ELSE TEMP_STRING = STRING;	70	PAD
52 M	RETURN TEMP_STRING;	71	PAD
53 M	CLOSE PAD;	72	HALINHAL
CI		73	HALINHAL
CI		74	HALINHAL
CI	HERE IS THE FUNCTION I_FORMAT FROM HALPASS1	75	HALINHAL
CI		76	HALINHAL
54 M	I_FORMAT: FUNCTION(NUMBER,WIDTH) CHARACTER(*);	77	I_FORMAT
55 M	DECLARE CHARACTER(255) VARYING, STRING;	78	I_FORMAT
56 M	DECLARE INTEGER, L, NUMBER, WIDTH;	79	I_FORMAT
57 M	STRING = NUMBER;	80	I_FORMAT
58 M	L = LENGTH(STRING);	81	I_FORMAT
59 M	IF L < WIDTH THEN	82	I_FORMAT
59 M	STRING = X70\$(1 TO WIDTH-L)    STRING;	83	I_FORMAT
60 M	RETURN STRING;	84	I_FORMAT
61 M	CLOSE I_FORMAT;	85	HALINHAL
CI		86	HALINHAL
CI		87	HALINHAL
CI	THE FOLLOWING IS THE ERROR PROCEDURE FROM PASS1.	88	HALINHAL

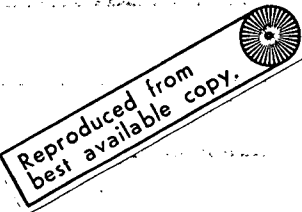


STMT	SOURCE	LINE	CURRENT SCOPE
	CI	80	HALINHAL
62 M	ERRORS: PROCEDURE (MESSAGE,SEVERITY):	90	ERRORS
	CI PRINTS AND ACCOUNTS FOR ALL ERRORS	91	ERRORS
63 M	DECLARE MSG CHARACTER(255) VARYING;	92	ERRORS
64 M	DECLARE MESSAGE CHARACTER(*);	93	ERRORS
65 M	DECLARE SEVERITY INTEGER;	94	ERRORS
66 M	ERROR_COUNT = ERROR_COUNT + 1;	95	ERRORS
67 M	MSG = '***** ERROR # '    ERROR_COUNT    ' OF SEVERITY '	96	ERRORS
67 M	SEVERITY    ': '    MESSAGE;	97	ERRORS
68 M	IF ERROR_COUNT > 1 THEN	98	ERRORS
68 M	MSG = MSG    '. LAST ERROR ON LINE '    PREVIOUS_ERROR;	99	ERRORS
69 M	WRITE(6)MSG    '. *****';	100	ERRORS
70 M	SAVE_SEVERITY\$(ERROR_COUNT) = SEVERITY;	101	ERRORS
71 M	SAVE_LINE\$(ERROR_COUNT) = CARD_COUNT;	102	ERRORS
72 M	PREVIOUS_ERROR = CARD_COUNT;	103	ERRORS
73 M	IF ERROR_COUNT >= 99 THEN GO;	104	ERRORS
74 M	IF ERROR_COUNT = 100 THEN	105	ERRORS
74 M	WRITE(6) 'TOO MANY ERRORS: COMPILATION ABORTED.';	106	ERRORS
75 M	IFF COMPILING THEN COMPILING = FALSE;	107	ERRORS
76 M	ELSE GO TO DISASTER; /* THEN END IT THERE */	108	ERRORS
77 M	END;	109	ERRORS
78 M	IF SEVERITY > MAX_SEVERITY THEN MAX_SEVERITY = SEVERITY;	110	ERRORS
79 M	IF SEVERITY > STATEMENT_SEVERITY THEN	111	ERRORS
79 M	STATEMENT_SEVERITY = SEVERITY;	112	ERRORS
80 M	IF SEVERITY > 3 THEN IFF COMPILING THEN	113	ERRORS
80 M	COMPILING = FALSE;	114	ERRORS
81 M	ELSE GO TO DISASTER;	115	ERRORS
	CI END OF PASS1 ERROR	116	ERRORS

STMT	SOURCE	LINE	CURRENT SCOPE
82 M	CLOSE ERRORS; /* CLOSING ERRORS AND RETURNING. */	117	HALINHAL
C1		118	HALINHAL
C1		119	HALINHAL
C1	THIS NEXT PROCEDURE IS PRINT_SUMMARY	120	HALINHAL
C1		121	HALINHAL
83 M	PRINT_SUMMARY: PROCEDURE;	122	PRINT_SUMMARY
84 M	WRITE(6)CARD_COUNT    ' CARDS WERE PROCESSED';	123	PRINT_SUMMARY
85 M	IF ERRCR_COUNT = 0 THEN WRITE(6) ' NO ERRORS WERE DETECTED.';	124	PRINT_SUMMARY
86 M	ELSE IF ERROR_COUNT > 1 THEN DO;	125	PRINT_SUMMARY
87 M	WRITE(6) ERRCR_COUNT    ' ERRORS WERE DETECTED,'	126	PRINT_SUMMARY
87 M	' THE LAST ERROR WAS ON LINE '    PREVIOUS_ERROR;	127	PRINT_SUMMARY
88 M	WRITE(6) '*****SUMMARY OF DETECTED ERRORS.*****';	128	PRINT_SUMMARY
89 M	DO FOR I = 1 TO ERROR_COUNT;	129	PRINT_SUMMARY
90 M	WRITE(6) 'ERROR # '    I    ' ON LINE '	130	PRINT_SUMMARY
90 M	SAVE_LINE\$I    ' OF SEVERITY '	131	PRINT_SUMMARY
90 M	SAVE_SEVERITY\$I    ' .';	132	PRINT_SUMMARY
91 M	END;	133	PRINT_SUMMARY
92 M	END;	134	PRINT_SUMMARY
93 M	ELSE WRITE(6) 'ONE ERROR WAS DETECTED WHICH OCCURRED ON LINE '	135	PRINT_SUMMARY
93 M	PREVIOUS_ERROR    ' OF SEVERITY '    SAVE_SEVERITY\$1    ' .';	136	PRINT_SUMMARY
94 M	TIME2 = TIME;	137	PRINT_SUMMARY
95 M	T = TIME2 - TIME1;	138	PRINT_SUMMARY
96 M	WRITE(6) 'CARD-PROCESSING RATE: '    6000 * CARD_COUNT / T	139	PRINT_SUMMARY
96 M	' CARDS PER MINUTE.';	140	PRINT_SUMMARY
97 M	WRITE(6) ' CLOCK TIME IS '    TIME2;	141	PRINT_SUMMARY
98 M	RETURN;	142	PRINT_SUMMARY
99 M	CLOSE PRINT_SUMMARY;	143	HALINHAL
100 M	STREAM: PROCEDURE;	144	STREAM
C1	THIS PROCEDURE FILLS THE VARIABLES NEXT_CHAR,	145	STREAM
C1	ARROW, AND OVER_FUNCH.	146	STREAM

STMT	SOURCE	LINE	CURRENT SCOPE
CI	NEXT_CHAR IS A ONE BYTE VARIABLE THAT CONTAINS THE	147	STREAM
CI	NEXT CHARACTER IN THE INPUT STREAM.	148	STREAM
CI	ARROW IS A HALF WORD VARIABLE THAT CONTAINS THE	149	STREAM
CI	INTEGER WHICH REPRESENTS THE RELATIVE	150	STREAM
CI	DISPLACEMENT OF THE CHARACTER IN NEXT_CHAR	151	STREAM
CI	WITH RESPECT TO THE LAST CHARACTER.	152	STREAM
CI	A POSITIVE VALUE INDICATES A MOVE UP.	153	STREAM
CI	OVER_PUNCH IS A ONE BYTE VARIABLE THAT IS FILLED WITH	154	STREAM
CI	A NON ZERO VALUE WHEN A CHARACTER OTHER	155	STREAM
CI	THAN A BLANK APPEARS DIRECTLY OVER AN	156	STREAM
CI	M-LIN CHARACTER--THE VALUE IS THE BYTE VALUE	157	STREAM
CI	OF THE OVER_PUNCH.	158	STREAM
CI		159	STREAM
101 MI	DECLARE CHARACTER(255) VARYING, E_LINE, S_LINE, M_LINE;	160	STREAM
102 MI	DECLARE CHARACTER(90), SAVE_CARD, CURRENT_CARD;	161	STREAM
103 MI	DECLARE BLANKS CHARACTER(128) INITIAL(' ');	162	STREAM
104 MI	DECLARE E_LINE_ERR CHARACTER(50) VARYING	163	STREAM
104 MI	INITIAL('E_LINE CHARACTER MORE THAN 1 LINE ABOVE PRECEEDING CHARACTER');	164	STREAM
105 MI	DECLARE ORDER_ERR CHARACTER(27) INITIAL('SOURCE PROGRAM OUT OF ORDER');	165	STREAM
106 MI	DECLARE CHARACTER(1), PREV_CARD, SAVE_NEXT_CHAR, SAVE_OVER_PUNCH;	166	STREAM
107 MI	DECLARE INTEGER INITIAL(0), LAST_E_IND, LAST_S_IND, E_BLANKS, M_BLANKS,	167	STREAM
107 MI	S_BLANKS, EP, SP, TEXT_LIMIT, E_COUNT, LAST_E_COUNT, S_COUNT,	168	STREAM
107 MI	LAST_S_COUNT, INDX, CP, POINTER, ARROW, II;	169	STREAM
108 MI	DECLARE ARRAY(256) INTEGER INITIAL(0), E_IND, S_IND, E_INDICATOR,	170	STREAM
109 MI	S_INDICATOR;	171	STREAM
109 MI	DECLARE BIT_1 INITIAL(FALSE), RETURNING_E, RETURNING_M, RETURNING_S,	172	STREAM
109 MI	END_GROUP, M_COMMENT, ARROW_FLAG;	173	STREAM
110 MI	DECLARE BIT_1 INITIAL(TRUE), FIRST_CALL_TO_STREAM;	174	STREAM
111 MI	DECLARE ARRAY(3) CHARACTER(1), TYPE_CHAR;	175	STREAM
112 MI	DECLARE APRAY(3) INTEGER INITIAL(0), RETURN_CHAR;	176	STREAM
113 MI	DECLARE CHARACTER(255) VARYING, E_STACK, S_STACK;	177	STREAM
114 MI	DECLARE INPUT_PAD CHARACTER(23) INITIAL	178	STREAM
114 MI	('M /* */ EOF' EOF EOF');	179	STREAM
CI		180	STREAM
CI		181	STREAM

STMT	SOURCE	LINE	CURRENT SCOPE
CI	TWO DIMENSIONAL PROCEDURES.	182	STREAM
CI		183	STREAM
115 MI	GO TO STREAM_START;	184	STREAM
116 MI	PROCESS_COMMENT: PROCEDURE;	185	PROCESS_COMMENT
117 MI	DECLARE K CHARACTER(1);	186	PROCESS_COMMENT
118 MI	DECLARE J INTEGER;	187	PROCESS_COMMENT
119 MI	IF CURRENT_CARD\$1 = 'F' THEN	188	PROCESS_COMMENT
119 MI	WRITE(6) LINE(1), CURRENT_CARD\$(2 TO 4) /* ISSUE-A NEW HEADER*/	189	PROCESS_COMMENT
119 MI	;	190	PROCESS_COMMENT
120 MI	ELSE	191	PROCESS_COMMENT
120 MI	IF CURRENT_CARD\$1 = 'C' THEN	192	PROCESS_COMMENT
120 MI	DO FOR I=1 TO TEXT_LIMIT-1;	193	PROCESS_COMMENT
121 MI	IF CURRENT_CARD\$I = ' ' THEN DO;	194	PROCESS_COMMENT
122 MI	K = CURRENT_CARD\$(I+1);	195	PROCESS_COMMENT
123 MI	J = CHAR_INDEX(TOGGLES, K);	196	PROCESS_COMMENT
124 MI	IF J /= 0 THEN DO;	197	PROCESS_COMMENT
125 MI	IF I < TEXT_LIMIT-1 THEN	198	PROCESS_COMMENT
125 MI	K = CURRENT_CARD\$(I+2);	199	PROCESS_COMMENT
126 MI	ELSE GO TO COMPLEMENT;	200	PROCESS_COMMENT
127 MI	IF K = '+' THEN CONTROL = TRUE;	201	PROCESS_COMMENT
127 SI	J:	202	PROCESS_COMMENT
129 MI	ELSE IF K = '-' THEN CONTROL = FALSE;	203	PROCESS_COMMENT
129 SI	J:	204	PROCESS_COMMENT
129 MI	ELSE	205	PROCESS_COMMENT
129 MI	COMPLEMENT: CONTROL = NOT CONTROL ;	206	PROCESS_COMMENT
129 SI	J: J:	207	PROCESS_COMMENT
130 MI	END;	208	PROCESS_COMMENT
131 MI	END;	209	PROCESS_COMMENT
132 MI	END;	210	PROCESS_COMMENT
133 MI	CLOSE PROCESS_COMMENT;	211	STREAM



STMT	SOURCE	LINE	CURRENT SCOPE
CI		212	STREAM
CI		213	STREAM
134 M	STACK_RETURN_CHAR:	214	STREAM
134 M	PROCEDURE(NUMBER, CHARR):	215	STACK_RETURN_CHAR
135 M	DECLARE INTEGER, NUMBER, I:	216	STACK_RETURN_CHAR
136 M	DECLARE CHARR CHARACTER(1):	217	STACK_RETURN_CHAR
137 M	DO FOR I = 1 TO 3:	218	STACK_RETURN_CHAR
138 M	IF RETURN_CHAR = 0 THEN DO;	219	STACK_RETURN_CHAR
138 S	I:	220	STACK_RETURN_CHAR
139 M	RETURN_CHAR = NUMBER;	221	STACK_RETURN_CHAR
139 S	I:	222	STACK_RETURN_CHAR
140 M	TYPE_CHAR = CHARR;	223	STACK_RETURN_CHAR
140 S	I:	224	STACK_RETURN_CHAR
CI	THE FOLLOWING STATEMENT IS FOR DEBUGGING PURPOSES ONLY.	225	STACK_RETURN_CHAR
CI	OUT 'STACKED "'    CHARR    "' ("    NUMBER    ") - INDEX('    I    ')';	226	STACK_RETURN_CHAR
141 M	RETURN;	227	STACK_RETURN_CHAR
142 M	END;	228	STACK_RETURN_CHAR
143 M	END;	229	STACK_RETURN_CHAR
144 M	CLOSE STACK_RETURN_CHAR;	230	STREAM
CI		231	STREAM
CI		232	STREAM
145 M	READ_CARD:	233	STREAM
145 M	PROCEDURE :	234	READ_CARD
146 M	DECLARE END_OF_INPUT BIT_1 INITIAL(FALSE);	235	READ_CARD
147 M	IFF END_OF_INPUT THEN DO;	236	READ_CARD
148 M	CURRENT_CARD = INPUT_PAC;	237	READ_CARD
149 M	RETURN;	238	READ_CARD
150 M	END;	239	READ_CARD
151 M	READ(5) CURRENT_CARD;	240	READ_CARD
CI	END_OF_INPUT = TRUE;	241	READ_CARD
152 M	CARD_CCUNT = CARD_CCUNT + 1;	242	READ_CARD

STMT	SOURCE	LINE	CURRENT SCOPE
153 M	IF CARD_COUNT NOT = 0 THEN DO;	243	READ_CARD
154 M	WRITE(6) I_FORMAT(CARD_COUNT, 4)	244	READ_CARD
154 M	' '    SAVE_CARD\$    ' '    SAVE_CARD\$(2 TO #)    ' '    CARD_COUNT;	245	READ_CARD
155 M	END;	246	READ_CARD
156 M	SAVE_CARD = CURRENT_CARD;	247	READ_CARD
157 M	CLOSE READ_CARD;	248	STREAM
		249	STREAM
		250	STREAM
158 M	ORDER_OK:	251	STREAM
159 M	FUNCTION(TYPE) BIT_1;	252	ORDER_OK
159 M	DECLARE TYPE CHARACTER(1);	253	ORDER_OK
160 M	CC CASE CARD_TYPE(CURRENT_CARD\$1) + 1;	254	ORDER_OK
161 M	DO; /*CASE 1--ILLEGAL CARD TYPE*/	255	ORDER_OK CASE 1
162 M	END_GROUP = FALSE;	256	ORDER_OK
163 M	RETURN FALSE;	257	ORDER_OK
164 M	END; /*OF CASE 1*/	258	ORDER_OK
165 M	/* CASE 2--E CARD*/	259	ORDER_OK
165 M	E_CARD:	260	ORDER_OK CASE 2
165 M	DO;	261	ORDER_OK
166 M	IF CARD_TYPE(TYPE) = 2 OR	262	ORDER_OK
166 M	CARD_TYPE(TYPE) = 3 THEN END_GROUP = TRUE;	263	ORDER_OK
167 M	ELSE END_GROUP = FALSE;	264	ORDER_OK
168 M	RETURN TRUE;	265	ORDER_OK
169 M	END; /*CASE 2*/	266	ORDER_OK
170 M	/*CASE 3--M CARD*/ GO TO E_CARD;	267	ORDER_OK CASE 3
171 M	DO; /* CASE 4--S CARD*/	268	ORDER_OK CASE 4
172 M	END_GROUP = FALSE;	269	ORDER_OK
173 M	IF CARD_TYPE(TYPE) = 2 OR	270	ORDER_OK

STMT	SOURCE	LINE	CURRENT SCOPE
173 MI	CARD_TYPE(TYPE) = 3 THEN RETURN TRUE;	271	ORDER_OK
174 MI	ELSE RETURN FALSE;	272	ORDER_OK
175 MI	END; /* CASE 4*/	273	ORDER_OK
176 MI	DC; /*CASE 5--A COMMENT*/	274	ORDER_OK CASE 5
177 MI	IF CARD_TYPE(TYPE) = 2 OR	275	ORDER_OK
177 MI	CARD_TYPE(TYPE) = 3 THEN END_GROUP = TRUE;	276	ORDER_OK
178 MI	ELSE END_GROUP = FALSE;	277	ORDER_OK
179 MI	IF CURRENT_CARDS1 = 'P' THEN DO;	278	ORDER_OK
180 MI	IF TYPE = 'C' THEN RETURN TRUE;	279	ORDER_OK
181 MI	ELSE RETURN FALSE;	280	ORDER_OK
182 MI	END;	281	ORDER_OK
183 MI	ELSE DC;	282	ORDER_OK
184 MI	IF CARD_TYPE(TYPE) = 1 THEN RETURN FALSE;	283	ORDER_OK
185 MI	ELSE RETURN TRUE;	284	ORDER_OK
186 MI	END;	285	ORDER_OK
187 MI	END; /* CASE 5 */	286	ORDER_OK
188 MI	END; /*CF DO CASE*/	287	ORDER_OK CASE 6
189 MI	CLOSE ORDER_OK;	288	STREAM
CI		289	STREAM
CI		290	STREAM
190 MI	COMMENT:	291	STREAM
190 MI	FUNCTION BIT_1;	292	COMMENT
191 MI	IF CURRENT_CARDSCP = '/' THEN	293	COMMENT
191 MI	IF CP < TEXT_LIMIT THEN	294	COMMENT
191 MI	IF CURRENT_CARDS(CP + 1) = '*' THEN DO;	295	COMMENT
CI	LOOK FOR END OF COMMENT.	296	COMMENT
192 MI	DC FOR CP = CP+2 TO TEXT_LIMIT-1;	297	COMMENT
193 MI	IF CURRENT_CARDSCP = '*' THEN	298	COMMENT



STMT	SOURCE	LINE	CURRENT	SCOPE
192 M	IF CURRENT_CARD\$(CP+1) = '/' THEN	299	COMMENT	
193 M	RETURN TRUE;	300	COMMENT	
192 M	DO;	301	COMMENT	
194 M	CP = CP + 1;	302	COMMENT	
195 M	RETURN TRUE;	303	COMMENT	
196 M	END;	304	COMMENT	
197 M	END;	305	COMMENT	
198 M	GO TO RETURN_TRUE;	306	COMMENT	
199 M	END;	307	COMMENT	
200 M	RETURN FALSE;	308	COMMENT	
201 M	CLOSE COMMENT;	309	STREAM	
CI		310	STREAM	
CI		311	STREAM	
202 M	SCAN_CARD:	312	STREAM	
202 M	PROCEDURE(TYPE,CCOUNT) ASSIGN(LIN,INDICATOR);	313	SCAN_CARD	
203 M	DECLARE INDICATOR ARRAY(256) INTEGER;	314	SCAN_CARD	
204 M	DECLARE INTEGER, TYPE, CCOUNT;	315	SCAN_CARD	
205 M	DECLARE LIN CHARACTER(*);	316	SCAN_CARD	
206 M	DO FOR CP = 2 TO TEXT_LIMIT;	317	SCAN_CARD	
207 M	IF CURRENT_CARD\$CP NOT = ' ' THEN DO;	318	SCAN_CARD	
209 M	IF COMMENT THEN GO TO CONTINUE;	319	SCAN_CARD	
209 M	IF LIN\$CP NOT = ' ' THEN	320	SCAN_CARD	
209 M	DO CASE TYPE + 1; /*HAL DO CASES START WITH 1*/	321	SCAN_CARD	
210 M	/* CASE 1*/ CALL ERRORS('OVERLAPPING E-LINE CHARACTERS',1);	322	SCAN_CARD CASE 1	
211 M	DO; /* CASE 2*/	323	SCAN_CARD CASE 2	
212 M	CALL ERRORS('OVERLAPPING S-LINE CHARACTERS',1);	324	SCAN_CARD	
213 M	GO TO CONTINUE;	325	SCAN_CARD	
214 M	END; /*CP CASE 2*/	326	SCAN_CARD	

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STMT	SOURCE	LINE	CURRENT SCOPE
215 M	END; /*CF DO CASE TYPE + 1*/	327	SCAN_CARD CASE 3
216 M	INDICATOR\$CP = CCUNT;	328	SCAN_CARD
217 M	LIN\$CP = CUPRENT_CARD\$CP;	329	SCAN_CARD
218 M	END;	330	SCAN_CARD
219 M	CONTINUE;	331	SCAN_CARD
219 M	END;	332	SCAN_CARD
220 M	CLOSE SCAN_CARD;	333	STREAM
	CI	334	STREAM
	CI	335	STREAM
221 M	COMP:	336	STREAM
221 M	PROCEDURE(TYPE) ASSIGN(LIN,INDICATOR,CCUNT);	337	COMP
222 M	DECLARE PCINT CHARACTER(1);	338	COMP
223 M	DECLARE INTEGER, CCUNT, TYPE;	339	COMP
224 M	DECLARE CHARACTER(*), LIN;	340	COMP
225 M	DECLARE ARRAY(256) INTEGER, INDICATOR;	341	COMP
226 M	IF TYPE = 1 THEN PCINT = 'S';	342	COMP
227 M	ELSE PCINT = 'E';	343	COMP
228 M	CCUNT = 1;	344	COMP
229 M	DO FOREVER;	345	COMP
230 M	CALL SCAN_CARD(TYPE,CCUNT) ASSIGN	346	COMP
230 M	(LIN,INDICATOR);	347	COMP
3" 231 M	CALL READ_CARD;	348	COMP
232 M	IF CURRENT_CARD\$1 NOT = PCINT THEN DO;	349	COMP
	CI NO MORE OF THIS TYPE CARD.	350	COMP
2" 233 M	IF_NOT ORDER_CHK(PCINT) THEN	351	COMP
233 M	CALL ERRORS(ORDER_ERR,1);	352	COMP
234 M	DO FOR CP = 2 TO TEXT_LIMIT;	353	COMP
1" 235 M	IF LIN\$CP = ' ' THEN	354	COMP

STMT	SOURCE	LINE	CURRENT SCOPE
235 MI	INDICATOR\$CP = 0;	355	COMP
236 MI	ELSE IF TYPE = 0 THEN DO;	356	COMP
237 MI	INDICATOR\$CP = CCOUNT-INDICATOR\$CP +1;	357	COMP
238 MI	END;	358	COMP
239 MI	END;	359	COMP
240 MI	RETURN;	360	COMP
241 MI	END;	361	COMP
242 MI	CCOUNT = CCOUNT + 1;	362	COMP
243 MI	END;	363	COMP
244 MI	CLOSE COMP;	364	STREAM
CI		365	STREAM
CI		366	STREAM
245 MI	GFT_GROUP:	367	STREAM
245 MI	PROCEDURE;	368	GET_GROUP
246 MI	E_LINE = E_LINE\$(INDX TO #)    BLANKS;	369	GET_GROUP
247 MI	S_LINE = S_LINE\$(INDX TO #)    BLANKS;	370	GET_GROUP
248 MI	LAST_E_COUNT = E_COUNT;	371	GET_GROUP
249 MI	LAST_S_COUNT = S_COUNT;	372	GET_GROUP
250 MI	E_COUNT, S_COUNT = 0;	373	GET_GROUP
251 MI	GO TO LOOP;	374	GET_GROUP
252 MI	READ_IT:	375	GET_GROUP
252 MI	CALL READ_CARD;	376	GET_GROUP
253 MI	IF NOT ORDER_CHK(PREV_CARD) THEN DO;	377	GET_GROUP
254 MI	CALL ERRORS(CRDER_ERR,1);	378	GET_GROUP
255 MI	GO TO READ_IT;	379	GET_GROUP
256 MI	END;	380	GET_GROUP
257 MI	LOOP:	381	GET_GROUP
257 MI	IF END_GROUP THEN GO TO FCUND_GROUP;	382	GET_GROUP

3"

2"

1"

STMT	SOURCE	LINE	CURRENT SCOPE
259 M	DO CASE CARD_TYPE(CURRENT_CARD\$1) + 1;	382	GET_GROUP
259 M	DO; /*CASE 1, A DUMMY*/ END;	384	GET_GROUP CASE 1
261 M	DO; /*CASE 2 E-LIN*/	385	GET_GROUP CASE 2
262 M	CALL CCMP(0) ASSIGN	386	GET_GROUP
262 M	(E_LINE,F_INDICATOR,E_CCUNT);	387	GET_GROUP
263 M	GO TO LOOP;	388	GET_GROUP
264 M	END; /*OF CASE 2*/	389	GET_GROUP
265 M	DO; /*CASE 3 M-LIN*/	390	GET_GROUP CASE 3
266 M	M_LINE = M_LINE\$(INDX TO #)    CURRENT_CARD\$(2 TO #);	391	GET_GROUP
267 M	SAVE_CARD\$1 = 'M';	392	GET_GROUP
268 M	PREV_CARD = CURRENT_CARD\$1;	393	GET_GROUP
269 M	GO TO READ_IT;	394	GET_GROUP
270 M	END; /*OF CASE 3*/	395	GET_GROUP
271 M	DO; /*CASE 4 S-LIN*/	396	GET_GROUP CASE 4
272 M	CALL CCMP(1) ASSIGN	397	GET_GROUP
272 M	(S_LINE,S_INDICATOR,S_CCUNT);	398	GET_GROUP
273 M	GO TO LOOP;	399	GET_GROUP
274 M	END; /*OF CASE 4*/	400	GET_GROUP
275 M	DO; /*CASE 5 COMMENT*/	401	GET_GROUP CASE 5
276 M	PREV_CARD = CURRENT_CARD\$1;	402	GET_GROUP
277 M	CALL PROCESS_COMMENT;	403	GET_GROUP
278 M	GO TO READ_IT;	404	GET_GROUP
279 M	END; /*OF CASE 5*/	405	GET_GROUP
280 M	END; /*OF DO CASE*/	406	GET_GROUP CASE 6
281 M	FOUND_GROUP:	407	GET_GROUP
281 M	WRITE(6) SKIP(1);	408	GET_GROUP
282 M	END_GROUP = FALSE;	409	GET_GROUP
283 M	E_LINE = E_LINE\$(1 TO LENGTH(M_LINE));	410	GET_GROUP

STMT	SOURCE	LINE	CURRENT SCOPE
284 MI	IF E_COUNT NOT > 0 THEN DO;	411	GET_GROUP
285 MI	DO FOR CP = 2 TO TEXT_LIMIT;	412	GET_GROUP
286 MI	E_INDICATOR\$CP = 0;	413	GET_GROUP
287 MI	END;	414	GET_GROUP
288 MI	E_COUNT = LAST_E_COUNT;	415	GET_GROUP
289 MI	END;	416	GET_GROUP
290 MI	S_LINE = S_LINES(1 TO LENGTH(M_LINE));	417	GET_GROUP
291 MI	IF S_COUNT NOT > 0 THEN DO;	418	GET_GROUP
292 MI	DO FOR CP = 2 TO TEXT_LIMIT;	419	GET_GROUP
293 MI	S_INDICATOR\$CP = 0;	420	GET_GROUP
294 MI	END;	421	GET_GROUP
295 MI	S_COUNT = LAST_S_COUNT;	422	GET_GROUP
296 MI	END;	423	GET_GROUP
CI	THE FOLLOWING STATEMENTS ARE FOR DEBUGGING PURPOSES ONLY	424	GET_GROUP
CI	OUT 'E_LINE='    E_LINE    ' ';	425	GET_GROUP
CI	OUT 'M_LINE='    M_LINE    ' ';	426	GET_GROUP
CI	OUT 'S_LINE='    S_LINE    ' ';	427	GET_GROUP
297 MI	CLOSE GET_GROUP;	428	STREAM
CI		429	STREAM
CI		430	STREAM
298 MI	CHOP:	431	STREAM
298 MI	PROCEDURE;	432	CHOP
299 MI	INDX = INDX + 1;	433	CHOP
300 MI	IF INDX = TEXT_LIMIT THEN DO;	434	CHOP
CI	OUT OF DATA, GET MORE.	435	CHOP
301 MI	E_INDICATOR\$1 = E_INDICATOR\$TEXT_LIMIT;	436	CHOP
302 MI	S_INDICATOR\$1 = S_INDICATOR\$TEXT_LIMIT;	437	CHOP
303 MI	CALL GET_GROUP;	438	CHOP
304 MI	INDX = 1;	439	CHOP
305 MI	END;	440	CHOP

STMT	SOURCE	LINE	CURRENT SCOPE
306 M	CLOSE CHOP;	441	STREAM
	CI	442	STREAM
	CI	443	STREAM
307 M	STACK:	444	STREAM
307 M	PROCEDURE(TYPE, INDICATOR, LIN) ASSIGN	445	STACK
307 M	(IND, STACK, PP);	446	STACK
308 M	DECLARE INTEGER, TYPE, PP;	447	STACK
309 M	DECLARE CHARACTER(*), LIN, STACK;	448	STACK
310 M	DECLARE ARRAY(256) INTEGER, INDICATOR, IND;	449	STACK
311 M	IF PP < 1 THEN GO TO NOT_MULTIPLE;	450	STACK
312 M	IF LIN\$INDX = ' ' THEN DO;	451	STACK
313 M	IF STACK\$PP = ' ' THEN	452	STACK
313 M	IND\$PP = IND\$PP + 1;	453	STACK
314 M	ELSE GO TO NOT_MULTIPLE;	454	STACK
315 M	END;	455	STACK
316 M	ELSE DO;	456	STACK
317 M	NOT_MULTIPLE:	457	STACK
317 M	PP = PP + 1;	458	STACK
318 M	IF PP > 256 THEN DO CASE TYPE+1;	459	STACK
319 M	CALL ERRORS('EXPCNENT STRING OVER FLOW', 3);	460	STACK CASE 1
320 M	CALL ERRORS('SUBSCRIPT STRING OVER FLOW', 3);	461	STACK CASE 2
321 M	END; /*CF DO CASE*/	462	STACK CASE 3
322 M	STACK = STACK    LIN\$INDX;	463	STACK
323 M	IND\$PP = INDICATOR\$INDX;	464	STACK
324 M	END;	465	STACK
325 M	CLOSE STACK;	466	STREAM
	CI	467	STREAM
	CI	468	STREAM
326 M	BUILD_XSCRIPTS:	469	STREAM

STMT	SOURCE	LINE	CURRENT SCOPE
326 M	PROCEDURE;	470	BUILD_XSCRIPTS
327 M	E_STACK, S_STACK = '';	471	BUILD_XSCRIPTS
328 M	F_BLANKS, S_BLANKS = -1;	472	BUILD_XSCRIPTS
329 M	EP, SP = 0;	473	BUILD_XSCRIPTS
330 M	CHECK_M:	474	BUILD_XSCRIPTS
330 M	IF M_LINE\$INDX = ' ' THEN	475	BUILD_XSCRIPTS
330 M	PROCESS:	476	BUILD_XSCRIPTS
330 M	DO;	477	BUILD_XSCRIPTS
331 M	CALL STACK(0,E_INDICATOR,E_LINE) ASSIGN	478	BUILD_XSCRIPTS
331 M	(E_IND,E_STACK,EP);	479	BUILD_XSCRIPTS
332 M	CALL STACK(1,S_INDICATOR,S_LINE) ASSIGN	480	BUILD_XSCRIPTS
332 M	(S_IND,S_STACK,SP);	481	BUILD_XSCRIPTS
333 M	CALL CHOP;	482	BUILD_XSCRIPTS
334 M	GO TO CHECK_M;	483	BUILD_XSCRIPTS
335 M	END;	484	BUILD_XSCRIPTS
336 M	ELSE /* A NON BLANK ON S-LIN*/	485	BUILD_XSCRIPTS
336 M	DC;	486	BUILD_XSCRIPTS
337 M	IFF M_COMMENT THEN DC;	487	BUILD_XSCRIPTS
339 M	IF M_LINE\$INDX = '*' THEN	488	BUILD_XSCRIPTS
339 M	IF M_LINE\$(INDX +1) = '/' THEN DC;	489	BUILD_XSCRIPTS
339 M	M_COMMENT = FALSE;	490	BUILD_XSCRIPTS
340 M	M_LINE\$(INDX +1) = ' ';	491	BUILD_XSCRIPTS
341 M	END;	492	BUILD_XSCRIPTS
342 M	GO TO PROCESS;	493	BUILD_XSCRIPTS
343 M	END;	494	BUILD_XSCRIPTS
344 M	IF M_LINE\$INDX = '/' THEN	495	BUILD_XSCRIPTS
344 M	IF M_LINE\$(INDX +1) = '*' THEN DO;	496	BUILD_XSCRIPTS
345 M	M_COMMENT = TRUE;	497	BUILD_XSCRIPTS

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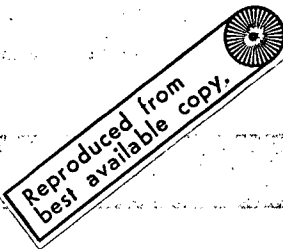
STMT	SOURCE	LINE	CURRENT SCOPE
346 M	M_LINE\$(INDX +1) = ' ';	498	BUILD_XSCRIPTS
347 M	GO TO PROCESS;	499	BUILD_XSCRIPTS
348 M	END;	500	BUILD_XSCRIPTS
349 M	IF S_STACK\$SP = ' ' THEN DO;	501	BUILD_XSCRIPTS
350 M	IF SP > 1 THEN S_STACK = S_STACK\$(1 TO SP-1);	502	BUILD_XSCRIPTS
351 M	ELSE S_STACK = ' ';	503	BUILD_XSCRIPTS
352 M	S_BLANKS = S_IND\$SP;	504	BUILD_XSCRIPTS
353 M	END;	505	BUILD_XSCRIPTS
354 M	IF E_STACK\$EP = ' ' THEN DO;	506	BUILD_XSCRIPTS
355 M	IF EP > 1 THEN E_STACK = E_STACK\$(1 TO EP-1);	507	BUILD_XSCRIPTS
356 M	ELSE E_STACK = ' ';	508	BUILD_XSCRIPTS
357 M	E_BLANKS = E_IND\$EP;	509	BUILD_XSCRIPTS
358 M	END;	510	BUILD_XSCRIPTS
359 M	IF E_BLANKS >= S_BLANKS THEN	511	BUILD_XSCRIPTS
359 M	M_BLANKS = S_BLANKS;	512	BUILD_XSCRIPTS
360 M	ELSE M_BLANKS = E_BLANKS;	513	BUILD_XSCRIPTS
361 M	END;	514	BUILD_XSCRIPTS
C	THE FOLLOWING STATEMENTS ARE FOR DEBUGGING PURPOSES ONLY	515	BUILD_XSCRIPTS
C	OUT 'E_STACK='    E_STACK    ' ';	516	BUILD_XSCRIPTS
C	OUT 'S_STACK='    S_STACK    ' ';	517	BUILD_XSCRIPTS
362 M	CLOSE BUILD_XSCRIPTS;	518	STREAM
C		519	STREAM
C		520	STREAM
363 M	STREAM_START;	521	STREAM
363 M	IFF FIRST_CALL_TO_STREAM THEN DO;	522	STREAM
364 M	TIME1 = TIME;	523	STREAM
365 M	FIRST_CALL_TO_STREAM = FALSE;	524	STREAM
366 M	AGAIN;	525	STREAM
366 M	CALL READ_CARD;	526	STREAM



STMT	SOURCE	LINE	CURRENT	SCOPE
367 M	IF_NCT CRDER_CK('C') THEN DO;	527	STREAM	
368 M	CALL ERRORS(CRDER_ERR,1);	528	STREAM	
369 M	GO TO AGAIN;	529	STREAM	
370 M	END;	530	STREAM	
371 M	TEXT_LIMIT = LENGTH(CURRENT_CARD) ;	531	STREAM	
372 M	RETURNING_M = TRUE;	532	STREAM	
373 M	M_LINE, E_LINE, S_LINE = ' ';	533	STREAM	
374 M	INDX = 1;	534	STREAM	
375 M	CALL GET_GROUP;	535	STREAM	
376 M	M_PLANKS = -1;	536	STREAM	
377 M	INDX = 2;	537	STREAM	
378 M	II = 1;	538	STREAM	
379 M	END;	539	STREAM	
380 M	IFF MACRO_FOUND THEN DO;	540	STREAM	
381 M	IF MACRC_POINT < MACRO_LIMIT THEN DO;	541	STREAM	
382 M	NEXT_CHAR =	542	STREAM	
382 M	MACRC_STREAM\$MACRC_POINT;	543	STREAM	
383 M	MACRO_POINT = MACRC_POINT +1;	544	STREAM	
384 M	RETURN;	545	STREAM	
385 M	END;	546	STREAM	
386 M	IF SAVE_BLANK_COUNT >= 0 THEN DO;	547	STREAM	
387 M	NEXT_CHAR = ' ';	548	STREAM	
388 M	BLANK_COUNT = SAVE_BLANK_COUNT;	549	STREAM	
389 M	SAVE_BLANK_COUNT = -1;	550	STREAM	
390 M	RETURN;	551	STREAM	
391 M	END;	552	STREAM	
392 M	MACRO_FOUND = FALSE;	553	STREAM	
393 M	NEXT_CHAR = SAVE_NEXT_CHAR;	554	STREAM	

STMT	SOURCE	LINE	CURRENT	SCOPE
394 M	OVER_PUNCH = SAVE_OVER_PUNCH;	555	STREAM	
395 M	MACRC_STREAM = '';	556	STREAM	
396 M	RETURN;	557	STREAM	
397 M	END;	558	STREAM	
398 M	BLANK_CCUNT = -1;	559	STREAM	
399 M	STACK_CHECK;	560	STREAM	
399 M	DO FOR II = II TO 3; /*CHECK LIMIT*/	561	STREAM	
400 M	IF RETURN_CHAR NOT = 0 THEN DO;	562	STREAM	
400 S	II:	563	STREAM	
401 M	ARROW_FLAG = TRUE;	564	STREAM	
402 M	RETURN_CHAR = RETURN_CHAR -1;	565	STREAM	
402 S	II:        II:	566	STREAM	
403 M	NEXT_CHAR = TYPE_CHAR ;	567	STREAM	
403 S	II:	568	STREAM	
404 M	OVER_PUNCH = 0;	569	STREAM	
405 M	RETURN;	570	STREAM	
406 M	END;	571	STREAM	
407 M	END;	572	STREAM	
408 M	II = 1;	573	STREAM	
409 M	IFF ARROW_FLAG THEN DO;	574	STREAM	
410 M	ARROW_FLAG = FALSE;	575	STREAM	
411 M	NEXT_CHAR = SAVE_NEXT_CHAR;	576	STREAM	
412 M	OVER_PUNCH = SAVE_OVER_PUNCH;	577	STREAM	
413 M	BLANK_CCUNT = SAVE_BLANK_COUNT;	578	STREAM	
414 M	RETURN;	579	STREAM	
415 M	END;	580	STREAM	
416 M	BEGINING;	581	STREAM	
416 M	IFF RETURNING_M THEN DO;	582	STREAM	
417 M	IF M_BLANKS >= 0 THEN DO;	583	STREAM	

STMT	SOURCE	LINE	CURRENT SCOPE
418 M	NEXT_CHAR = ' ';	584	STREAM
419 M	ARROW = - LAST_E_IND;	585	STREAM
420 M	LAST_E_IND = C;	586	STREAM
421 M	FLANK_CCUNT = M_BLANKS;	587	STREAM
422 M	M_BLANKS = -1;	588	STREAM
423 M	GO TO FLUND_CHAR;	589	STREAM
424 M	END;	590	STREAM
425 M	IF M_LINE\$INDX NOT = ' ' THEN DO;	591	STREAM
426 M	IF M_LINE\$INDX = '/' THEN	592	STREAM
426 M	IF M_LINE\$(INDX +1) = '*' THEN DO;	593	STREAM
C	HERE IS A START OF A COMMENT	594	STREAM
427 M	M_COMMENT = TRUE;	595	STREAM
428 M	M_LINE\$INDX = ' ';	596	STREAM
429 M	M_LINE\$(INDX +1) = ' ';	597	STREAM
430 M	GO TO BLANK;	598	STREAM
431 M	END;	599	STREAM
432 M	IF E_CCUNT > 0 THEN DO;	600	STREAM
433 M	IF E_LINE\$INDX NOT = ' ' THEN DO;	601	STREAM
434 M	IF E_INDICATOR\$INDX NOT = 1 THEN	602	STREAM
434 M	CALL ERRORS('E-LINE OVERLAPS M-LINE',1);	603	STREAM
435 M	ELSE OVER_PUNCH = E_LINE\$INDX;	604	STREAM
436 M	END;	605	STREAM
437 M	ELSE OVER_PUNCH = 0;	606	STREAM
438 M	END;	607	STREAM
439 M	ELSE OVER_PUNCH = 0;	608	STREAM
440 M	IF S_CCUNT > 0 THEN	609	STREAM
440 M	IF S_LINE\$INDX NOT = ' ' THEN	610	STREAM
440 M	CALL ERRORS('S-LINE CVFRLAPS M-LINE',1);	611	STREAM



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STMT	SOURCE	LINE	CURRENT	SCOPE
		612	STREAM	
441 MI	APROW = - LAST_E_IND;	613	STREAM	
442 MI	LAST_E_IND = 0;	614	STREAM	
443 MI	NEXT_CHAR = M_LINE\$INX;	615	STREAM	
444 MI	CALL CHCP;	616	STREAM	
445 MI	GO TO FCUND_CHAR;	617	STREAM	
446 MI	END;	618	STREAM	
447 MI	ELSE	619	STREAM	
447 MI	PLANK:	620	STREAM	
447 MI	DO;	621	STREAM	
448 MI	CALL BUILD_XSCRIPTS;	622	STREAM	
449 MI	CVER_PUNCH = 0;	623	STREAM	
450 MI	RETURNING_M = FALSE;	624	STREAM	
451 MI	LAST_S_IND = 0;	625	STREAM	
452 MI	RETURNING_S = TRUE;	626	STREAM	
453 MI	PCINTER = 1; /*THIS ONE MAY NOT BE NEEDED*/	627	STREAM	
454 MI	END;	628	STREAM	
455 MI	END;	629	STREAM	
456 MI	IFF RETURNING_S THEN DO;	630	STREAM	
457 MI	IF LENGTH(S_STACK) > 0 &	631	STREAM	
457 MI	PCINTER <= LENGTH(S_STACK) THEN DO;	632	STREAM	
458 MI	IF S_STACK\$PCINTER = ' ' THEN DO;	633	STREAM	
459 MI	IF S_IND\$PCINTER >= 0 THEN	634	STREAM	
	CI MORE LEFT	635	STREAM	
459 MI	DO;	636	STREAM	
460 MI	NEXT_CHAR = ' ';	637	STREAM	
461 MI	BLANK_COUNT = S_IND\$PCINTER;	638	STREAM	
462 MI	PCINTER = PCINTER + 1;	639	STREAM	
463 MI	APROW = LAST_S_IND - S_IND\$PCINTER;			

STMT	SOURCE	LINE	CURRENT SCOPE
464 M	LAST_S_IND = S_IND\$POINTER;	640	STREAM
465 M	END;	641	STREAM
466 M	END;	642	STREAM
467 M	ELSE DC; /*A NON BLANK*/	643	STREAM
468 M	NEXT_CHAR = S_STACK\$POINTER;	644	STREAM
469 M	ARROW = LAST_S_IND - S_IND\$POINTER;	645	STREAM
470 M	LAST_S_IND = S_IND\$POINTER;	646	STREAM
471 M	POINTER = POINTER + 1;	647	STREAM
472 M	END;	648	STREAM
473 M	GO TO FOUND_CHAR;	649	STREAM
474 M	END;	650	STREAM
475 M	ELSE DC; /* CAN NOT RETURN*/	651	STREAM
476 M	RETURNING_S = FALSE;	652	STREAM
477 M	RETURNING_E = TRUE;	653	STREAM
478 M	LAST_E_IND = -LAST_S_IND;	654	STREAM
479 M	POINTER = 1;	655	STREAM
480 M	END;	656	STREAM
481 M	END;	657	STREAM
482 M	IF RETURNING_E THEN DC;	658	STREAM
483 M	IF LENGTH(E_STACK) > 0 &	659	STREAM
484 M	POINTER <= LENGTH(E_STACK) THEN DC;	660	STREAM
485 M	IF E_STACK\$POINTER = ' ' THEN DO;	661	STREAM
486 M	IF E_IND\$POINTER >= 0 THEN	662	STREAM
2"	C  MORE TO GO.	663	STREAM
488 M	DC;	664	STREAM
489 M	NEXT_CHAR = ' ';	665	STREAM
1"	490 M  BLANK_COUNT = E_IND\$POINTER;	666	STREAM
491 M	POINTER = POINTER + 1;	667	STREAM

STMT	SOURCE	LINE	CURRENT	SCOPE
489 MI	ARROW = E_IND\$PCINTER - LAST_E_IND;	668	STREAM	
490 MI	LAST_E_IND = E_IND\$PCINTER;	669	STREAM	
491 MI	END;	670	STREAM	
492 MI	END;	671	STREAM	
493 MI	ELSE	672	STREAM	
CI	A NON BLANK	673	STREAM	
497 MI	DO;	674	STREAM	
494 MI	NEXT_CHAR = F_STACK\$PCINTER;	675	STREAM	
495 MI	ARROW = E_IND\$PCINTER - LAST_E_IND;	676	STREAM	
496 MI	LAST_E_IND = E_IND\$PCINTER;	677	STREAM	
497 MI	PCINTER = PCINTER + 1;	678	STREAM	
498 MI	END;	679	STREAM	
499 MI	GO TO FOUND_CHAR;	680	STREAM	
500 MI	END;	681	STREAM	
501 MI	ELSE	682	STREAM	
CI	CAN NOT RETURN	683	STREAM	
501 MI	DO;	684	STREAM	
502 MI	RETURNING_E = FALSE;	685	STREAM	
503 MI	RETURNING_M = TRUE;	686	STREAM	
504 MI	END;	687	STREAM	
505 MI	END;	688	STREAM	
506 MI	GO TO BEGINING;	689	STREAM	
507 MI	FOUND_CHAR:	690	STREAM	
507 MI	IF ARROW NOT = 0 THEN DO;	691	STREAM	
508 MI	OLD_LEVEL = NEW_LEVEL;	692	STREAM	
509 MI	NEW_LEVEL = NEW_LEVEL + ARROW;	693	STREAM	
510 MI	SAVE_OVER_PUNCH = OVER_PUNCH;	694	STREAM	
511 MI	SAVE_NEXT_CHAR = NEXT_CHAR;	695	STREAM	

STMT	SOURCE	LINE	CURRENT SCOPE
512 M	SAVE_BLANK_COUNT = BLANK_COUNT;	696	STREAM
513 M	IF OLD_LEVEL > 0 THEN DO;	697	STREAM
514 M	IF ARRCW < 0 THEN	698	STREAM
514 M	CALL STACK_RETURN_CHAR(-ARROW, '(');	699	STREAM
515 M	ELSE	700	STREAM
515 M	EXONENT:	701	STREAM
515 M	DO;	702	STREAM
516 M	IF ARRCW > 1 THEN	703	STREAM
516 M	CALL ERRORS(E_LINE_ERR, 1);	704	STREAM
517 M	CALL STACK_RETURN_CHAR(2, '*');	705	STREAM
518 M	CALL STACK_RETURN_CHAR(ARROW, '(');	706	STREAM
519 M	END;	707	STREAM
520 M	END;	708	STREAM
521 M	ELSE IF OLD_LEVEL = 0 THEN DO;	709	STREAM
522 M	IF ARRCW < 0 THEN	710	STREAM
522 M	SUBS:	711	STREAM
522 M	DO;	712	STREAM
523 M	IF ARRCW < -1 THEN CALL ERRORS(	713	STREAM
523 M	'S-LINE CHARACTER MORE THAN 1 LINE BELOW PRECEDING CHARACTER', 1);	714	STREAM
524 M	CALL STACK_RETURN_CHAR(1, 'S');	715	STREAM
525 M	CALL STACK_RETURN_CHAR(-ARROW, '(');	716	STREAM
526 M	END;	717	STREAM
527 M	ELSE GO TO EXONENT;	718	STREAM
528 M	END;	719	STREAM
529 M	ELSE /*OLD < 0*/ DO;	720	STREAM
530 M	IF ARRCW < 0 THEN GO TO SUBS;	721	STREAM
531 M	IF NEW_LEVEL <= 0 THEN	722	STREAM
531 M	CALL STACK_RETURN_CHAR(ARROW, '(');	723	STREAM

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STMT	SOURCE	LINE	CURRENT	SCOPE
532 M	ELSE DC;	724	STREAM	
533 M	CALL STACK_RETURN_CHAR(-CLC_LEVEL, '*');	725	STREAM	
534 M	IF NEW_LEVEL > 1 THEN CALL ERRORS(E_LINE_ERR, 1);	726	STREAM	
535 M	CALL STACK_RETURN_CHAR(2, '*');	727	STREAM	
536 M	CALL STACK_RETURN_CHAR(NEW_LEVEL, '(');	728	STREAM	
537 M	END;	729	STREAM	
538 M	END;	730	STREAM	
539 M	ARROW = 0;	731	STREAM	
540 M	GO TO STACK_CHECK;	732	STREAM	
541 M	END;	733	STREAM	
542 M	RETURN;	734	STREAM	
543 M	CLOSE STREAM;	735	HALINHAL	
CI		736	HALINHAL	
CI		737	HALINHAL	
CI	PROGRAM TO TEST THE STREAM PROCEDURE	738	HALINHAL	
CI		739	HALINHAL	
544 M	MAIN_PROGRAM:	740	HALINHAL	
544 M	WRITE(6) 'BEGIN TEST OF HAL IN HAL';	741	HALINHAL	
545 M	BUILT, BUILT_UP, BUILT_TOKEN, BUILT_C_P='';	742	HALINHAL	
546 M	MAIN_LOOP: DO FOREVER;	743	HALINHAL	
547 M	CALL STREAM;	744	HALINHAL	
548 M	IF NEXT_CHAR = ' ' THEN DC;	745	HALINHAL	
549 M	IF BLANK_FLAG THEN GO TO MAIN_LOOP;	746	HALINHAL	
550 M	ELSE BLANK_FLAG = TRUE;	747	HALINHAL	
551 M	END;	748	HALINHAL	
552 M	ELSE BLANK_FLAG = FALSE;	749	HALINHAL	
553 M	BUILT = BUILT    NEXT_CHAR;	750	HALINHAL	
554 M	IF CVER_PUNCH = '0' THEN DO;	751	HALINHAL	
CI	OUT 'STREAM RETURNS "'    NEXT_CHAR    "' - BLANKS='	752	HALINHAL	
CI	BLANK_COUNT;	753	HALINHAL	



STMT	SOURCE	LINE	CURRENT SCOPE
555 M	BUILT_UP = BUILT_UP    ! ';	754	HALINHAL
556 M	END;	755	HALINHAL
557 M	ELSE DO;	756	HALINHAL
C	OUT 'STREAM RETURNS "'    NEXT_CHAR    "' - OVER PUNCH "'	757	HALINHAL
C	OVER_PUNCH    "'    ' - BLANKS='    BLANK_COUNT;	758	HALINHAL
558 M	BUILT_UP = BUILT_UP    OVER_PUNCH;	759	HALINHAL
559 M	END;	760	HALINHAL
560 M	IF NEXT_CHAR NOT < 'A' THEN DO;	761	HALINHAL
561 M	BUILT_TOKEN = BUILT_TOKEN    NEXT_CHAR;	762	HALINHAL
562 M	IF OVER_PUNCH NOT = '0' THEN BUILT_O_P = OVER_PUNCH;	763	HALINHAL
563 M	END;	764	HALINHAL
564 M	ELSE DO;	765	HALINHAL
565 M	IF LENGTH(BUILT_TOKEN) NOT = 0 THEN DO;	766	HALINHAL
566 M	IF BUILT_O_P = '' THEN WRITE(6) '***TOKEN='    BUILT_TOKEN;	767	HALINHAL
567 M	ELSE WRITE(6) '***TOKEN='    BUILT_TOKEN    ',MARKER='    BUILT_O_P;	768	HALINHAL
568 M	BUILT_TOKEN, BUILT_C_P = '';	769	HALINHAL
569 M	END;	770	HALINHAL
570 M	IF NEXT_CHAR NOT = ' ' THEN WRITE(6) '***TOKEN='    NEXT_CHAR;	771	HALINHAL
571 M	ELSE WRITE(6) '***BLANKS='    BLANK_COUNT + 1;	772	HALINHAL
572 M	END;	773	HALINHAL
573 M	IF NEXT_CHAR = ';' THEN DO;	774	HALINHAL
574 M	WRITE(6) SKIP(2), '***CVP***'    BUILT_UP;	775	HALINHAL
575 M	WRITE(6) '***MAIN***'    BUILT, SKIP(2);	776	HALINHAL
576 M	BUILT, BUILT_UP = '';	777	HALINHAL
577 M	END;	778	HALINHAL
578 M	IF NEXT_CHAR = '?' THEN GO TO DISASTER;	779	HALINHAL
579 M	END; /* DO FOREVER */	780	HALINHAL
580 M	DISASTER:	781	HALINHAL

STMT	SOURCE	LINE	CURRENT SCOPE
580 M	CALL PRINT_SUMMARY;	782	HALINHAL
581 M	WRITE(6) * THIS TEST IS NOW COMPLETE.;	783	HALINHAL
582 M	CLOSE HALINHAL;	784	
C1		785	
C1	A B C D E	786	

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2"  
1"

SYMBOL TABLE LISTING:

LCC	NAME	TYPE	CLASS	LENGTH	PRECISION	NEST	OUTER_LEVEL	FLAGS	SYT_PTR	ARRAY
1	HALNHAL	PROG	LABEL	0	0	0	0	40	2	
2	R	0	REPL	0	0	1	0	8000	1	
3	FALSE	0	REPL	0	0	1	0	8000	2	
4	TRUE	0	REPL	0	0	1	0	9000	3	
5	BIT_1	0	REPL	0	0	1	0	8000	4	
6	IF_NOT	0	REPL	0	0	1	0	8000	5	
7	IF	0	REPL	0	0	1	0	8000	6	
8	CUT	0	REPL	0	0	1	0	8000	7	
9	FOREVER	0	REPL	0	0	1	0	8000	8	
10	XTO	CHAR	VAR	70	FIXED	1	0	9A0B	0	
11	BUILT	CHAR	VAR	255	VARYING	1	0	820B	0	
12	BUILT_UP	CHAR	VAR	255	VARYING	1	0	820B	0	
12	BUILT_TOKEN	CHAR	VAR	255	VARYING	1	0	820B	0	
14	BUILT_C_P	CHAR	VAR	255	VARYING	1	0	920B	0	
15	PLANK_FLAG	INT	VAR	0	0	1	0	8A0B	0	
16	PREFCR_COUNT	INT	VAR	0	0	1	0	8A0B	0	
17	MACRO_PUNIT	INT	VAR	0	0	1	0	8A0B	0	
18	MACRO_LIMIT	INT	VAR	0	0	1	0	8A0B	0	
19	CLR_LEVEL	INT	VAR	0	0	1	0	8A0B	0	
20	NEW_LEVEL	INT	VAR	0	0	1	0	9A0B	0	
21	MAX_SEVERITY	INT	VAR	0	0	1	0	8A0B	0	
22	STATEMENT_SEVERITY	INT	VAR	0	0	1	0	9A0B	0	
23	SAVE_SEVERITY	INT	ARRAY	0	0	1	0	8A0B	0	100
24	SAVE_LINE	INT	ARRAY	0	0	1	0	8A0B	0	100
25	PREVIOUS_ERROR	INT	VAR	0	0	1	0	8A0B	0	
26	I	INT	VAR	0	0	1	0	8A0B	0	
27	COMPILING	INT	VAR	0	0	1	0	8A0B	0	
28	TIME1	INT	VAR	0	0	1	0	820B	0	
29	TIME2	INT	VAR	0	0	1	0	820B	0	
30	MACRO_STREAM	CHAR	VAR	255	VARYING	1	0	820B	0	
31	MACRO_COUNT	INT	VAR	0	0	1	0	8A0B	0	
32	CARD_COUNT	INT	VAR	0	0	1	0	8A0B	0	
33	DISASTER	STMT	LABEL	0	0	1	0	8040	0	
34	NEXT_CHAR	CHAR	VAR	1	FIXED	1	0	820B	0	
35	OVER_PUNCH	CHAR	VAR	1	FIXED	1	0	820B	0	
36	PLANK_COUNT	INT	VAR	0	0	1	0	8A0B	0	
37	CONTROL	BIT	ARRAY	32	0	1	0	820B	0	10
38	TOGGLES	CHAR	VAR	10	FIXED	1	0	8A0B	0	
39	CARD_TYPE	INT	FUNC	0	0	1	0	8040	40	
40	SELECT	CHAR	VAR	1	FIXED	2	0	860B	0	
41	CHAR_INDEX	INT	FUNC	0	0	1	0	8040	42	
42	STRING	CHAR	VAR	**	FIXED	2	0	860B	0	
43	PATTERN	CHAR	VAR	**	FIXED	2	0	860B	0	
44	I	INT	VAR	0	0	2	0	820B	0	
45	J	INT	VAR	0	0	2	0	820B	0	
46	K	INT	VAR	0	0	2	0	820B	0	
47	PAD	CHAR	FUNC	255	VARYING	1	0	8040	48	
48	STRING	CHAR	VAR	**	FIXED	2	0	860B	0	
49	WIDTH	INT	VAR	0	0	2	0	860B	0	
50	TEMP_STRING	CHAR	VAR	255	VARYING	2	0	820B	0	
51	L	INT	VAR	0	0	2	0	820B	0	
52	I_FORMAT	CHAR	FUNC	**	FIXED	1	0	8040	53	

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LCC	NAME	TYPE	CLASS	LENGTH	PRECISION	NEST	OUTER_LEVEL	FLAGS	SYT_PTR	ARRAY
53	NUMBER	INT	VAR	0	0	2	0	860B	0	
54	WIDTH	INT	VAR	0	0	2	0	860B	0	
55	STRING	CHAR	VAR	255	VARYING	2	0	820B	0	
56	L	INT	VAR	0	0	2	0	C20B	0	
57	ERRORS	PROC	LABEL	0	0	1	0	8040	58	
58	MESSAGE	CHAR	VAR	**	FIXED	2	0	840B	0	
59	SEVERITY	INT	VAR	0	0	2	0	840B	0	
60	MSG	CHAR	VAR	255	VARYING	2	0	820B	0	
61	DISASTER	I-ST	LABEL	0	0	2	0	C040	33	
62	PRINT_SUMMARY	PRCC	LABEL	0	0	1	0	8040	0	
63	STREAM	PROC	LABEL	0	0	1	0	8040	64	
64	F_LINE	CHAR	VAR	255	VARYING	2	0	820B	0	
65	S_LINE	CHAR	VAR	255	VARYING	2	0	820B	0	
66	M_LINE	CHAR	VAR	255	VARYING	2	0	820B	0	
67	SAVE_CARD	CHAR	VAR	80	FIXED	2	0	820B	0	
68	CURRENT_CARD	CHAR	VAR	80	FIXED	2	0	820B	0	
69	BLANKS	CHAR	VAR	128	FIXED	2	0	8A0B	0	
70	F_LINE_FRR	CHAR	VAR	50	VARYING	2	0	8A0B	0	
71	ORDER_FRR	CHAR	VAR	27	FIXED	2	0	8A0B	0	
72	PREV_CARD	CHAR	VAR	1	FIXED	2	0	820B	0	
73	SAVE_NEXT_CHAR	CHAR	VAR	1	FIXED	2	0	820B	0	
74	SAVE_OVER_PUNCH	CHAR	VAR	1	FIXED	2	0	820B	0	
75	LAST_E_IND	INT	VAR	0	C	2	0	8A0B	0	
76	LAST_S_IND	INT	VAR	0	C	2	0	8A0B	0	
77	E_BLANKS	INT	VAR	0	0	2	0	8A0B	0	
78	M_BLANKS	INT	VAR	0	0	2	0	8A0B	0	
79	S_BLANKS	INT	VAR	0	0	2	0	8A0B	0	
80	FP	INT	VAR	0	0	2	0	8A0B	0	
81	SP	INT	VAR	0	0	2	0	8A0B	0	
82	TEXT_LIMIT	INT	VAR	0	0	2	0	8A0B	0	
83	F_COUNT	INT	VAR	0	0	2	0	8A0B	0	
84	LAST_E_CCUNT	INT	VAR	0	0	2	0	8A0B	0	
85	S_CCUNT	INT	VAR	0	0	2	0	8A0B	0	
86	LAST_S_CCUNT	INT	VAR	0	0	2	0	8A0B	0	
87	INDX	INT	VAR	0	0	2	0	8A0B	0	
88	CP	INT	VAR	0	0	2	0	8A0B	0	
89	POINTER	INT	VAR	0	C	2	0	8A0B	0	
90	ARROW	INT	VAR	0	0	2	0	8A0B	0	
91	II	INT	VAR	0	0	2	0	8A0B	0	
92	F_IND	INT	ARRAY	0	0	2	0	8A0B	0	256
93	S_IND	INT	ARRAY	0	0	2	0	8A0B	0	256
94	E_INDICATOR	INT	ARRAY	0	0	2	0	8A0B	0	256
95	S_INDICATOR	INT	ARRAY	0	0	2	0	8A0B	0	256
96	RETURNING_E	INT	VAR	0	0	2	0	8A0B	0	
97	RETURNING_M	INT	VAR	0	0	2	0	8A0B	0	
98	RETURNING_S	INT	VAR	0	0	2	0	8A0B	0	
99	END_GROUP	INT	VAR	0	0	2	0	8A0B	0	
100	M_COMMENT	INT	VAR	0	0	2	0	8A0B	0	
101	ARROW_FLAG	INT	VAR	0	0	2	0	8A0B	0	
102	FIRST_CALL_TC_STREAM	INT	VAR	0	0	2	0	8A0B	0	
103	TYPE_CHAR	CHAR	ARRAY	1	FIXED	2	0	820B	0	3
104	RETURN_CHAR	INT	ARRAY	0	0	2	0	8A0B	0	3
105	E_STACK	CHAR	VAR	255	VARYING	2	0	820B	0	
106	S_STACK	CHAR	VAR	255	VARYING	2	0	820B	0	
107	INPUT_PAC	CHAR	VAR	23	FIXED	2	0	8A0B	0	
108	STREAM_START	STMT	LABEL	0	0	2	0	8040	0	

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LCC	NAME	TYPE	CLASS	LENGTH	PRECISION	NEST	OUTER_LEVEL	FLAGS	SYT_PTR	ARRAY
109	PROCESS_COMMENT	PROC	LABEL	0	0	2	0	8040	110	
110	K	CHAR	VAR	1	FIXED	3	0	820B	0	
111	J	INT	VAR	0	0	3	0	820B	0	
112	COMPLEMENT	STMT	LABEL	0	0	3	0	C040	0	
113	STACK_RETURN_CHAR	PROC	LABEL	0	0	2	0	8040	114	
114	NUMBER	INT	VAR	0	0	3	0	860B	0	
115	CHAR	CHAR	VAR	1	FIXED	3	0	860B	0	
116	I	INT	VAR	0	0	3	0	C20B	0	
117	READ_CARD	PROC	LABEL	0	0	2	0	8040	118	
118	END_OF_INPUT	INT	VAR	0	0	3	0	CA0B	0	
119	ERRR_CHK	INT	FUNC	0	0	2	0	8040	120	
120	TYPE	CHAR	VAR	1	FIXED	3	0	860B	0	
121	F_CARD	STMT	LABEL	0	0	3	0	C040	0	
122	COMMENT	INT	FUNC	0	0	2	0	8040	123	
123	RETURN_TRUE	STMT	LABEL	0	0	3	0	C040	0	
124	SCAN_CARD	PROC	LABEL	0	0	2	0	8040	125	
125	TYPE	INT	VAR	0	0	3	0	860B	0	
126	COUNT	INT	VAR	0	0	3	0	860B	0	
127	LIN	CHAR	VAR	**	FIXED	3	0	822B	0	
128	INDICATOR	INT	ARRAY	0	0	3	0	822B	0	256
129	CONTINUE	STMT	LABEL	0	0	3	0	8040	0	
130	ERRORS	I-CL	LABEL	0	0	3	0	C040	57	
131	COMP	PROC	LABEL	0	0	2	0	8040	132	
132	TYPE	INT	VAR	0	0	3	0	860B	0	
133	LIN	CHAR	VAR	**	FIXED	3	0	822B	0	
134	INDICATOR	INT	ARRAY	0	0	3	0	822B	0	256
135	CCOUNT	INT	VAR	0	0	3	0	822B	0	
136	PCINT	CHAR	VAR	1	FIXED	3	0	820B	0	
137	SCAN_CARD	I-CL	LABEL	0	0	3	0	8040	124	
138	READ_CARD	I-CL	LABEL	0	0	3	0	8040	117	
139	ERRORS	I-CL	LABEL	0	0	3	0	C040	57	
140	GET_GROUP	PROC	LABEL	0	0	2	0	8040	141	
141	LOCP	STMT	LABEL	0	0	3	0	8040	0	
142	READ_IT	STMT	LABEL	0	0	3	0	8040	0	
143	READ_CARD	I-CL	LABEL	0	0	3	0	8040	117	
144	ERRORS	I-CL	LABEL	0	0	3	0	8040	57	
145	FOUND_GROUP	STMT	LABEL	0	0	3	0	8040	0	
146	COMP	I-CL	LABEL	0	0	3	0	8040	131	
147	PROCESS_COMMENT	I-CL	LABEL	0	0	3	0	C040	109	
148	CHOP	PROC	LABEL	0	0	2	0	8040	140	
149	GET_GROUP	I-CL	LABEL	0	0	3	0	C040	140	
150	STACK	PROC	LABEL	0	0	2	0	8040	151	
151	TYPE	INT	VAR	0	0	3	0	860B	0	
152	INDICATOR	INT	ARRAY	0	0	3	0	860B	0	256
153	LIN	CHAR	VAR	**	FIXED	3	0	860B	0	
154	IND	INT	ARRAY	0	0	3	0	822B	0	256
155	STACK	CHAR	VAR	**	FIXED	3	0	822B	0	
156	PP	INT	VAR	0	0	3	0	822B	0	
157	NOT_MULTIPLE	STMT	LABEL	0	0	3	0	8040	0	
158	ERRORS	I-CL	LABEL	0	0	3	0	C040	57	
159	BUILD_XSCRIPTS	PROC	LABEL	0	0	2	0	8040	160	
160	CHECK_M	STMT	LABEL	0	0	3	0	8040	0	
161	PROCESS	STMT	LABEL	0	0	3	0	8040	0	
162	STACK	I-CL	LABEL	0	0	3	0	8040	150	
163	CHOP	I-CL	LABEL	0	0	3	0	C040	148	
164	AGAIN	STMT	LABEL	0	0	2	0	8040	0	

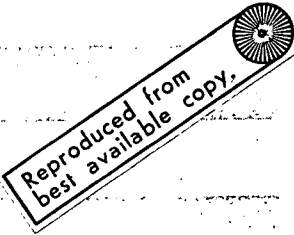
LOC	NAME	TYPE	CLASS	LENGTH	PRECISION	NEST	OUTER_LEVEL	FLAGS	SYT_PTR	ARRAY
165	ERRORS	I-CL	LABEL	0	0	2	0	8040	57	
166	*SAVE_BLANK_COUNT	INT	VAR	0	0	2	0	8218	0	
167	STACK_CHECK	STMT	LABEL	0	0	2	0	8040	0	
168	BEGINING	STMT	LABEL	0	0	2	0	8040	0	
169	FOUND_CHAR	STMT	LABEL	0	0	2	0	8040	0	
170	BLANK	STMT	LABEL	0	0	2	0	8040	0	
171	EXPCNENT	STMT	LABEL	0	0	2	0	8040	0	
172	SUBS	STMT	LABEL	0	0	2	0	C040	0	
173	MAIN_PROGRAM	STMT	LABEL	0	0	1	0	8040	0	
174	MAIN_LOOP	STMT	LABEL	0	0	1	0	C040	0	







102	C	0110	2	0263	4	0365														
103	0	C111	4	0140	2	C403														
104	0	0112	2	0139	4	0139	2	0400	6	0402										
105	0	0113	4	0327	4	0331	2	0354	6	0355	4	0356	2	0483	2	0494	2	0494		
106	0	0113	4	0327	4	0332	2	0349	6	0350	4	0351	2	0457	2	0458	2	0468		
107	0	0114	2	0149																
108	2	0115	0	0353																
109	0	0114	2	0277																
110	0	0117	4	0122	2	0123	4	0125	2	0127	2	0128								
111	0	0118	4	0123	2	C124	1	C127	1	0128	1	0129								
112	2	0126	0	0129																
113	C	0134	2	0514	2	0517	2	0518	2	0524	2	0525	2	0531	2	0533	2	0535	2	0536
114	0	0134	C	0135	2	C139														
115	0	0134	C	0136	2	C140														
116	0	0135	4	0137	1	0138	1	0139	1	0140										
117	0	C145	2	C231	2	C252	2	C366												
119	0	0146	2	0147																
119	0	0152	2	0233	2	0253	2	0367												
120	0	0153	C	0159	2	0166	2	0173	2	0177	2	0180	2	0184						
121	0	0165	2	0170																
122	0	0190	2	0208																
123	0	0193	2	0198																
124	0	0202	2	C230																
125	C	0202	C	0204	2	0209														
126	0	0202	C	0204	2	0216														
127	0	0202	C	0205	2	0209	4	0217												
128	0	0202	C	0203	4	C216														
129	2	0208	2	0213	0	C219														
131	0	0221	2	0262	2	0272														
132	C	0221	0	0223	2	0226	2	0230	2	0236										
133	0	0221	0	0224	4	0230	2	C235												
134	0	0221	0	0225	4	0230	4	0235	6	0237										
135	0	0221	0	0223	4	0228	2	0230	2	0237	6	0242								
136	0	0222	4	0226	4	0227	2	C232	2	0233										
140	0	0245	2	0303	2	0375														
141	2	0251	C	0257	2	0263	2	0273												
142	C	0252	2	0255	2	C269	2	0278												
145	2	0257	0	0281																
148	0	0298	2	0333	2	0444														
150	C	0307	2	0331	2	0332														
151	C	0307	C	0308	2	C318														
152	0	0307	C	0310	2	C323														
153	C	0307	C	0309	2	0312	2	C322												
154	C	0307	0	0310	6	C313	4	C323												
155	C	0307	C	0309	2	0313	6	C322												
156	0	0307	C	0308	2	0311	1	0312	6	0317	2	0318	1	0323						
157	2	0311	2	0314	0	0317														
159	C	0326	2	0448																
160	0	0320	2	0334																
161	0	0330	2	0342	2	0347														
164	0	0366	2	0369																
166	2	0396	2	0398	4	0399	2	0413	4	0512										
167	0	0399	2	0540																
169	0	0416	2	0506																
169	2	0423	2	0445	2	C473	2	0469	0	0507										
170	2	0430	0	0447																



171 0 0515 2 0527  
172 0 0522 2 0530  
173 0 0544 NOT REFERENCED  
174 0 0546 2 0549

## MACRO TEXT LISTING:

## LCC TEXT

1 REPLACE  
2 0  
3 1  
4 INTEGER  
5 IF 0 =  
6 IF 1 =  
7 WRITE(6)  
9 WHILE 1 = 1

STACKING DECISIONS = 13662  
CALLS TO SCAN = 3926  
CALLS TO IDENTIFY = 989  
NUMBER OF REDUCTIONS = 9727  
MAX STACK SIZE = 25  
MAX IND. STACK SIZE = 32  
END IND. STACK SIZE = 17  
MAX EXT\_ARRAY INDEX = 2  
XREF LIST ENTRIES = 870  
STATEMENT COUNT = 593  
MAX OUTER\_LIST INDEX = 0  
MAX NESTING DEPTH = 3  
FREE STRING AREA = 55355

786 CARDS WERE PROCESSED.  
NO ERRORS WERE DETECTED DURING PHASE 1 .

TOTAL ELAPSED TIME IN COMPILER 0:0:20.52.  
ELAPSED SET UP TIME 0:0:0.04.  
ACTUAL ELAPSED COMPILING TIME 0:0:19.13.  
ELAPSED CLEAN-UP TIME AT END 0:0:1.35.  
PROCESSING RATE: 2465 CARDS PER MINUTE.

BEGIN TEST OF HAL IN HAL

1 C| 2 5 6

2 M| TPARAM:PROGRAM;

\*\*\*BLANKS=1  
 \*\*\*TOKEN=TPARM  
 \*\*\*TOKEN=  
 \*\*\*TOKEN=PROGRAM  
 \*\*\*TOKEN=;

\*\*\*COVER\*\*\*

\*\*\*MAIN\*\*\* TPARAM:PROGRAM; \*

3 E| \*  
 4 M| DECLARE M1  
 5 S| 3,3

\*\*\*BLANKS=65  
 \*\*\*TOKEN=DECLARE  
 \*\*\*BLANKS=1  
 \*\*\*TOKEN=M1,MARKER=#  
 \*\*\*TOKEN=#  
 \*\*\*TOKEN=(  
 \*\*\*TOKEN=3  
 \*\*\*TOKEN=,  
 \*\*\*TOKEN=3  
 \*\*\*TOKEN=)  
 \*\*\*TOKEN=;

6 E| \*  
 7 M| M2 ;  
 8 S| 8,8

\*\*\*BLANKS=66  
 \*\*\*TOKEN=M2,MARKER=#  
 \*\*\*TOKEN=#  
 \*\*\*TOKEN=(  
 \*\*\*TOKEN=8  
 \*\*\*TOKEN=,  
 \*\*\*TOKEN=9  
 \*\*\*TOKEN=)  
 \*\*\*TOKEN=;

\*\*\*COVER\*\*\*

\*\*\*MAIN\*\*\* DECLARE M1\$(3,3), M2\$(8,8); \*

9 M| DECLARE I1 ARRAY(5) INTEGER INITIAL(3,2,5,4,1);

\*\*\*BLANKS=72  
 \*\*\*TOKEN=DECLARE  
 \*\*\*BLANKS=1  
 \*\*\*TOKEN=I1  
 \*\*\*BLANKS=1  
 \*\*\*TOKEN=ARRAY  
 \*\*\*TOKEN=(  
 \*\*\*TOKEN=5  
 \*\*\*TOKEN=)

```
**BLANKS=1
***TOKEN=INTEGER
**BLANKS=1
***TOKEN=INITIAL
***TOKEN=(
***TOKEN=3
***TOKEN=,
***TOKEN=2
***TOKEN=,
***TOKEN=5
***TOKEN=,
***TOKEN=4
***TOKEN=,
***TOKEN=1
***TOKEN=)
***TOKEN=;
```

```
***OVER***
***MAIN*** DECLARE I1 ARRAY(5) INTEGER INITIAL(3,2,5,4,1);
```

```
10 M1 DECLARE ARRAY(5) SCALAR, S1 , S2;
```

10

```
**BLANKS=32
***TOKEN=DECLARE
**BLANKS=1
***TOKEN=ARRAY
***TOKEN=(
***TOKEN=5
***TOKEN=)
**BLANKS=1
***TOKEN=SCALAR
***TOKEN=,
**BLANKS=1
***TOKEN=S1
**BLANKS=1
***TOKEN=,
**BLANKS=1
***TOKEN=S2
***TOKEN=;
```

```
***OVER***
***MAIN*** DECLARE ARRAY(5) SCALAR, S1 , S2;
```

3"

2"

1"

FUNCTION DECLARATIONS

```

11 H) FUNCTION DECLARATIONS
12 E) *
13 M) PROC: FUNCTION (P ) MATRIX (*,*)
14 S) **

```

11  
12  
13  
14

```

**BLANKS=47
***TOKEN=PROC
***TOKEN=:
**BLANKS=1
***TOKEN=FUNCTION
**BLANKS=2
***TOKEN=(
***TOKEN=P,MARKER=*
***TOKEN=$
***TOKEN=(
***TOKEN=*
***TOKEN=,
***TOKEN=*
***TOKEN=)
***TOKEN=)
**BLANKS=1
***TOKEN=MATRIX
**BLANKS=1
***TOKEN=(
***TOKEN=*
***TOKEN=,
***TOKEN=*
***TOKEN=)
***TOKEN=:

```

```

***OVER*** *
***MAIN*** PROC: FUNCTION (P$(*,*)) MATRIX (*,*)

```

```

15 E) *-1
16 M) RETURN P ;

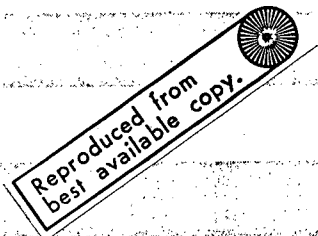
```

15  
16

```

**BLANKS=45
***TOKEN=RETURN
**BLANKS=1
***TOKEN=P,MARKER=*
***TOKEN=*
***TOKEN=*
***TOKEN=(
***TOKEN=-
***TOKEN=1
***TOKEN=)
***TOKEN=:

```



```

***OVER*** *
***MAIN*** RETURN P**(-1);

```

```

17 M) CLCSE PROC;

```

17

```

**BLANKS=67
***TOKEN=CLCSE

```

1"

```
**BLANKS=1
***TOKEN=PRO1
***TOKEN=;

***OVER***
***MAIN*** CLOSE PRO1;
```

```
18 MI PROC: FUNCTION (B) MATRIX (*,*); | 18
```

```
**BLANKS=67
***TOKEN=PRO
***TOKEN=;
**BLANKS=1
***TOKEN=FUNCTION
**BLANKS=1
***TOKEN=(
***TOKEN=B
***TOKEN=)
**BLANKS=1
***TOKEN=MATRIX
**BLANKS=1
***TOKEN=(
***TOKEN=*
***TOKEN=,
***TOKEN=*
***TOKEN=)
***TOKEN=;
```

```
***OVER***
***MAIN*** PROC: FUNCTION (B) MATRIX (*,*);
```

```
19 MI DECLARE MATRIX (*,*), P; | 19
```

```
**BLANKS=49
***TOKEN=DECLARE
**BLANKS=1
***TOKEN=MATRIX
**BLANKS=1
***TOKEN=(
***TOKEN=*
***TOKEN=,
***TOKEN=*
***TOKEN=)
***TOKEN=,
**BLANKS=1
***TOKEN=B
***TOKEN=;
```

```
***OVER***
***MAIN*** DECLARE MATRIX (*,*), P;
```

```
20 MI RETURN PRO1(B); | 20
21 SI 2 TO 4,1 TC 3 | 21
```

```
1" **BLANKS=56
```

```

***TOKEN=RETURN
**BLANKS=1
***TOKEN=PRG1
***TOKEN=(
***TOKEN=P
***TKFA=;
***TKEN=(
***TKEN=2
**BLANKS=1
***TKEN=TC
**BLANKS=1
***TKFA=4
***TKEN=,
***TKEN=1
**BLANKS=1
***TKEN=TO
**BLANKS=1
***TKEN=B
***TKEN=)
***TKEN=)
***TKEN=;

```

```

***OVER***
***MAIN*** RETURN PROJ(B$(2 TO 4,1 TO 3));

```

22 MI CLOSE PROC;

```

**BLANKS=50
***TKEN=CLOSE
**BLANKS=1
***TKEN=PRO
***TKEN=;

```

```

***OVER***
***MAIN*** CLGE PRO;

```

MAIN PROGRAM

23 MI MAIN PROGRAM

24 MI DO FOR I=1 TO 5;

| 23

| 24

```

**BLANKS=67
***TOKEN=DO
**BLANKS=1
***TOKEN=FOR
**BLANKS=1
***TOKEN=I
***TOKEN==
***TOKEN=1
**BLANKS=1
***TOKEN=TO
**BLANKS=1
***TOKEN=5
***TOKEN=;

```

```

***OVER***
***MAIN*** DO FOR I=1 TO 5;

```

25 MI DO FOR J=1 TO 5;

| 25

```

**BLANKS=63
***TOKEN=DO
**BLANKS=1
***TOKEN=FOR
**BLANKS=1
***TOKEN=J
***TOKEN==
***TOKEN=1
**BLANKS=1
***TOKEN=TO
**BLANKS=1
***TOKEN=5
***TOKEN=;

```

```

***OVER***
***MAIN*** DO FOR J=1 TO 5;

```

```

26 EI          J
27 EI          I I 5 I
28 MI M2 = S1 ;
29 SI I, J I I
30 SI          I

```

| 26

| 27

| 28

| 29

| 30

```

**BLANKS=65
***TOKEN=M2
***TOKEN=(
***TOKEN=I
***TOKEN=,
***TOKEN=J
***TOKEN=)
***TOKEN==
**BLANKS=1

```

2"

1"



```

***TOKEN=S1
***TOKEN=$
***TOKEN=(
***TOKEN=I1
***TOKEN=$
***TOKEN=(
***TOKEN=I
***TOKEN=)
***TOKEN=)
***TOKEN=#
***TOKEN=#
***TOKEN=(
***TOKEN=I1
***TOKEN=$
***TOKEN=I
***TOKEN=#
***TOKEN=#
***TOKEN=(
***TOKEN=J
***TOKEN=)
***TOKEN=)
***TOKEN=#;

```

```

***OVER***
***MAIN*** M2$(I,J)=S1$(I1$(I))**((I1$I**)(J));

```

```

31 MI END;

```

```

| 31

```

```

**BLANKS=62
***TOKEN=END
***TOKEN=:

```

```

***OVER***
***MAIN*** END;

```

```

32 C) THE NEXT GROUP HAS AN OVERLAPPING E LINE

```

```

| 32

```

```

33 E) I

```

```

| 33

```

```

***** ERROR # 1 OF SEVERITY 1: OVERLAPPING E-LINE CHARACTERS. *****

```

```

34 E) I1$I

```

```

| 34

```

```

35 MI S2 =S1 ; /* COMMENTS ARE BLANKS */

```

```

| 35

```

```

36 SI I I

```

```

| 36

```

```

3" ***BLANKS=75
***TOKEN=S2
***TOKEN=$
***TOKEN=(
***TOKEN=I
***TOKEN=)
***TOKEN=#
2" ***TOKEN=S1
***TOKEN=$
***TOKEN=(
***TOKEN=I
***TOKEN=)
1" ***TOKEN=#
***TOKEN=#

```



```
***TCKEN=(
***TCKEN=I1
***TCKEN=$
***TCKEN=I
***TCKEN=)
***TCKEN=;
```

```
***OVER***
```

```
***MAIN*** S2$(I)=S1$(I)**(I1$I);
```

37 E	2		37
****	ERROR # 2	OF SEVERITY 1: SCUPCE PROGRAM CUT OF ORDER. LAST ERROR ON LINE 33. ****	
38 C		THIS SHOULD DEMONSTRATE CARDS CUT OF ORDER	38
39 M	S2 = S2 ;		39
40 S	I		40

```
**BLANKS=67
```

```
**** ERROR # 3 OF SEVERITY 1: S-LINE OVERLAPS M-LINE. LAST ERROR ON LINE 37. ****
```

```
***TCKEN=S2
***BLANKS=1
***TCKEN=
***BLANKS=1
***TCKEN=S2
***TCKEN=$
***TCKEN=(
***TCKEN=I
***TCKEN=)
***TCKEN=*
***TCKEN=*
***TCKEN=(
***TCKEN=2
***TCKEN=)
***TCKEN=;
```

```
***OVER***
```

```
***MAIN*** S2 = S2$(I)**(2);
```

41 C	AN S LINE OVERLAP ABOVE	41
42 M	END;	42

```
**BLANKS=71
***TCKEN=END
***TCKEN=;
```

```
***OVER***
```

```
***MAIN*** END;
```

43 E	*	*	43
44 M	M1=PRC(M2 /* IGNORED */ );		44
45 S	4 AT 3,4 AT 5		45

```
**BLANKS=75
***TCKEN=M1,MARKER=*
***TCKEN=*
***TCKEN=PRC
```

```

***TOKEN=(
***TOKEN=M2,MARKER=
***TOKEN=
***TOKEN=(
***TOKEN=4
**BLANKS=1
***TOKEN=AT
**BLANKS=1
***TOKEN=3
***TOKEN=,
***TOKEN=4
**BLANKS=1
***TOKEN=AT
**BLANKS=1
***TOKEN=5
***TOKEN=)
**BLANKS=2
***TOKEN=)
***TOKEN=;

```

```

***COVER*** *
***MAIN*** M1=PRO(M2$(4 AT 3,4 AT 5) );

```

46 M | CLOSE TPARM;

| 46

```

**BLANKS=53
***TOKEN=CLOSEF
**BLANKS=1
***TOKEN=TPARM
***TOKEN=;

```

```

***COVER***
***MAIN*** CLOSE TPARM;

```

47 M | ?

| 47

```

**BLANKS=67
***TOKEN=?
47 CARDS WERE PROCESSED
3 ERRORS WERE DETECTED, THE LAST ERROR WAS ON LINE 40
*****SUMMARY OF DETECTED ERRORS*****
ERROR # 1 ON LINE 33 OF SEVERITY 1.
ERROR # 2 ON LINE 37 OF SEVERITY 1.
ERROR # 3 ON LINE 40 OF SEVERITY 1.
CARD-PROCESSING RATE: 4.4269995E+02 CARDS PER MINUTE.
CLOCK TIME IS 4361443
THIS TEST IS NOW COMPLETE.

```

3"

2"

1"

```

C 02 05 06
M TPARM: PROGRAM;
E *
M DECLARE M1
S 3,3
E *
M M2 ;
S 9,8
M DECLARE I1 ARRAY(5) INTEGER INITIAL(3,2,5,4,1);
M DECLARE ARRAY(5) SCALAR, S1 , S2;
H FUNCTION DECLARATIONS
E *
M PRO1: FUNCTION (P ) MATRIX (*,*);
S *,*
E *-1
M RETURN P ;
M CLOSE PRO1;
M PRO: FUNCTION (B) MATRIX (*,*);
M DECLARE MATRIX (*,*), B;
M RETURN PRO1(B
S 2 TO 4,1 TO 3
M CLOSE PRO;
H MAIN PROGRAM
M DO FOR I=1 TO 5;
M DO FOR J=1 TO 5;
E J
E I1$1
M M2 = S1 ;
S I,J I1
S I
M END;
C THE NEXT GROUP HAS AN OVERLAPPING E LINE
E I
E I1$1
M S2 =S1 ; /* COMMENTS ARE BLANKS */
S I I
E 2
C THIS SHOULD DEMONSTRATE CARDS OUT OF ORDER
MS2 = S2 ;
S I I
C AN S LINE OVERLAP ABOVE
M END;
E *
M M1=PRO(M2 /* IGNORED */ );
S 4 AT 3,4 AT 5
M CLOSE TPARM;
M ?

```